

**HP 3000 Computer Systems**

# **INTERPROCESS COMMUNICATION**



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# PREFACE

This book is a tutorial of Interprocess Communication internals and externals. It encompasses the information found in the MPE Intrinsic Manual, MPE File System Reference Manual and IPC code and reference material. The externals part of this document is meant for Hewlett Packard System Engineers or customers with an MPE and SPL background. IPC internals assumes the reader has an internal MPE and internal File System background.

Section 1 is a brief introduction to IPC, how it has evolved, where we use an IPC application. Section 2 is external IPC; how do we program IPC applications. There are two program examples at the end of this unit. One of these illustrates using soft interrupts.

The remainder of this manual is devoted to IPC internal data structures and their dynamics. Section 3 explains how IPC interacts with the MPE File System and Section 4 describes the queueing mechanisms for IPC and soft interrupts. Soft interrupt operation and soft interrupt procedures at the end of section 4 will help you to assimilate the material into a working model. For debugging information use the appendices A, B, C, and G. All of the material in this document is based on code written in MPE IV.

Larry Zeitman  
MPE Technical Support  
October 24, 1983

Interprocess communication (IPC) is a facility of the file system which permits multiple user processes to communicate with one another in an easy and efficient manner. It is the nonprivileged user's access to `nowait I/O`.<sup>1</sup> IPC uses message files to hold information between user processes. These message files act as first-in-first-out queues of records, with entries made by `FWRITES` and deletions made by `FREADS`.

Interprocess Communication consists of two parts. Basic IPC implements the functions as described above, and Soft Interrupts is an addition to basic IPC. The Soft Interrupt facility allows a process to trap to its own coded interrupt handler upon completing I/O.

Classically operating systems are interrupt driven, meaning that scheduling decisions, I/O, CPU or peripheral errors are all detected by the cpu by means of "interrupts". We talk of hard interrupts when either peripheral devices or hardware interrupt processing. Soft interrupts are then interrupts generated by software events. In either case the cpu is interrupted from processing and either microcode or software will cause the operating system to switch execution to an interrupt procedure. We refer to this sequence of events as a trap (also you may see trap referring to the interrupt code that executes). Soft interrupts within MPE means that when I/O completes the cpu will be interrupted, scheduling information for the process that initiated the I/O will be updated and the next time the process runs it will trap to a predetermined interrupt procedure. IPC with and without the soft interrupt feature will be explained continually throughout this manual.

## History

Prior to IPC various other system features (still available today) were used to implement process to process communication. These system features include:

### File System

- o Requires users to coordinate access, using `RINs` or `FLOCK & FUNLOCK`.
- o No mechanism to queue read or write requests to a file.
- o No mechanism to implement soft interrupts.

### Extra Data Segments

- o Processes must create and identify the DST to be used.
- o The DST is private to a session.
- o The user has to implement a locking scheme using `RINs`.
- o No mechanism to implement soft interrupts.

---

<sup>1</sup>See no-wait I/O in the MPE Intrinsic Manual.

## Introduction

### Mail

- o Transfers greater than one word require the creation and deletion of an extra data segment for each transfer.
- o Communication outside of the process tree is not supported.
- o Mail does not have the ability to queue messages.
- o Mail does not support soft interrupts.

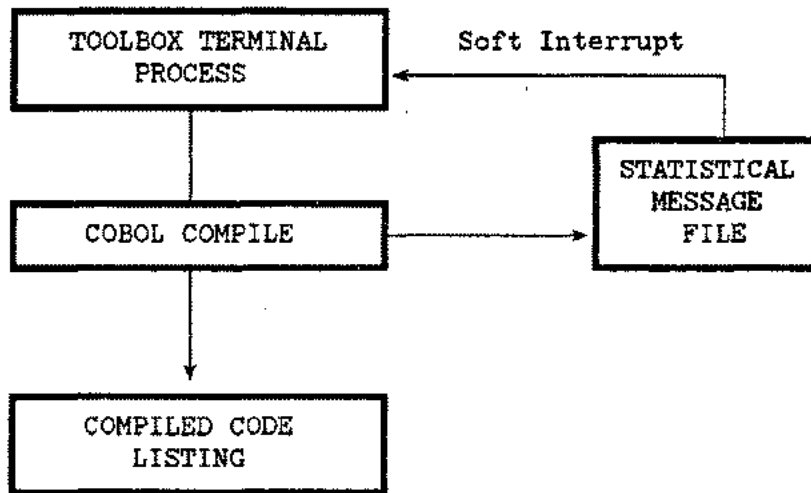
### IPC features include

- o The ability to send and receive variable-length messages or records quickly and easily.
- o The ability to support communication among multiple senders and receivers.
- o The ability to communicate with process trees.
- o The ability to queue reader and writer communication.
- o The ability to perform (non-privileged) user nowait send and receive requests.
- o Support of a trap facility when a send or receive completes (software interrupts).
- o Support of timeouts (timed limits to I/O inquiries).
- o The ability to coordinate user communication without the need for locking and un-locking overhead.



## Where Should IPC Be Used

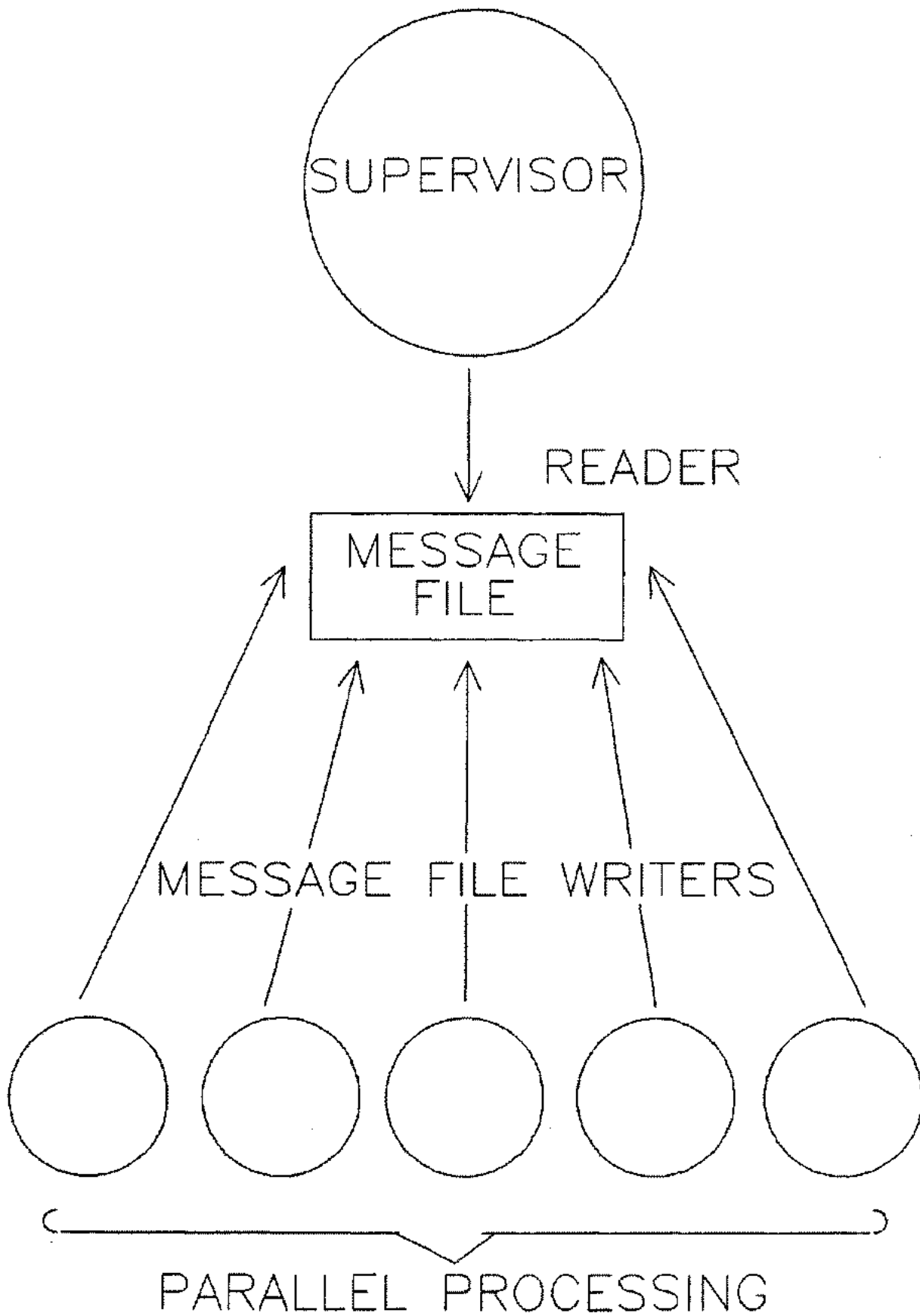
Part of IPC's power is its ability to implement parallel processing. As an example TOOLBOX using IPC.



The TOOLBOX terminal process edits COBOL source code. Then, a compile of that source code is initiated. Before the compile is started, the terminal process initiates a soft interrupt read to a statistical message file. The COBOL compile will actually occur in the background (as a son process). Once the compile has started, the terminal process is available to perform other tasks (e.g. write documentation, JCL). When the background compile is complete, the compiled source code is written to a special list file and the compile errors, warnings, and statistics are written to the statistical message file. Once the first record is written to the statistical message file, the pending soft interrupt read is satisfied, which interrupts the terminal process's activity to display the results of the COBOL compile.

Another useful implementation of Soft Interrupts would be in a supervisor/slave design where one/many slave processes are being managed by a supervisor process that performs other activities in parallel. The supervisor may request the status on the amount of transactions processed, whether the process is busy, or notification that the process needs to be terminated (e.g. this supervisor process could be servicing a terminal). Using IPC intrinsics the supervisor will post a read to a message file and wait for messages from slave processes. As messages are sent the supervisor can examine each one with an interrupt handler.

*(See the figure on the following page)*



## IPC Features

### Message Queuing

Messages are queued when a reader tries to read from an empty message file or a writer attempts to write to a full message file.

### Software Interrupts

A program may designate a trap procedure for interrupts. When I/O completes and an interrupt is generated, program control will be directed to the trap procedure. After the trap routine finishes execution, control is returned to mainline code.

### Writer ID's

When a writer process opens a message file, the file system assigns a unique 16-bit identification number to the writer. Each record the process writes to the message file is associated with this number. When the writer closes the file, the ID number no longer applies to the process and may be reused. The ID is accessible via FCONTROL 46. Note that only through programming would a reader be able to receive messages from a certain writer ID.

### Timeouts

A reader or a writer process may limit the length of time it will wait to be serviced. By issuing an FCONTROL 4, a process may specify the maximum number of seconds it will wait for its I/O request to be satisfied.

## External IPC

### Copy Access

When records are read from a message file, FREAD logically deletes them as it reads. In order to copy a message file without destroying it, the file must be opened with the file copy option specified in the FOPEN Aoptions, or the *COPY* keyword must be specified in a :FILE command. When this option is selected the message file is treated as a standard sequential file rather than as a message file, and it may be copied safely (ie, without the logical deletion of each record).

#### NOTE

*In order to access a message file in copy mode, a process must have exclusive access to the file.*

### Nondestructive Read

By issuing an FCONTROL 47, a reader may avoid deleting the record it reads, and the record will remain at the head of the message queue. This feature differs from the *copy access* feature in that it is a temporary condition. The second Fread following the Fcontrol 47 will reread the record and delete it in the usual manner.

### Global Multiaccess

Global Multiaccess allows processes in the same or different process trees to open the same message file. Each file will have the same record pointer but is allowed different "f" and "a" options. The global multiaccess option may be requested in the *aoptions* of the FOPEN to the file, or by using the *GMULTI* keyword in a :FILE command to define the file.

#### NOTE

*Global multiaccess is unavailable to message files when they have been opened with exclusive access in copy mode.*

### Appending to Variable Length Files

Variable-length files may be opened with append access. It is not necessary to have fixed-length records of the maximum possible size, so space is conserved.

## Fundamental Operation

Internal IPC procedures are called by the MPE file system and the file system is invoked by the file system intrinsics. Listed below is a brief description of those intrinsics that implement an IPC operation. The intrinsics supply the options to use the basic ipc facility and soft interrupts. The default is interprocess communication without traps.

**FOPEN** FOPEN establishes a connection to a message file. With FOPEN, a user process identifies itself as either a reader or a writer; readers access the head of the message file and writers access the tail. Incompatible parameters that are specified with FOPEN are adjusted. For example, since messages are read or written to the file one record at a time, a multirecord parameter is not allowed and the file system will correct this parameter when the file is opened. If FOPEN is used to access a new file, a new message file is created.

**FREAD** Fread reads one record from the head of a message file. The record is copied to the reader's TARGET area and is logically deleted from the message queue; the next record is now at the head of the file. If a process tries to read from an empty message file that writers have opened, the file system causes it to wait until a writer process enters a record to the file. If there are no writers associated with the message file, an end-of-file indication, CCG, is returned.

**NOTE**

*If the message file is empty and there are no writers, a read will be queued if there is either an FCONTROL 45 in effect or this is the first Fread after the reader's Fopen. See "Intrinsics - A Detailed Description page 2-11.*

**FWRITE** FWRITE appends one record to the tail of a message file. If a process tries to write to a full message file which readers are accessing, the file system causes it to wait until a reader process has read a block of records from the file. If there are no readers associated with the message file, an end-of-file indication, CCG, is returned.

**NOTE**

*If the message file is full and there are no readers, a write will be queued if there is either an FCONTROL 45 in effect, or this is the first Fwrite after the writer's FOPEN. See page 2-11.*

**FCONTROL** Supplies various control functions to a process that is using a message file. These control functions permit a process to take advantage of the additional features of IPC, which are discussed on pages 2-11 and 2-12.

External IPC

- FCLOSE** Breaks a process's connection with a message file. If the process reopens the same file later, it may do so as either a reader or a writer.
- FINSTATE** Enables or disables soft interrupts for the process. Finstate is issued some time before the process initiates a Soft Interrupt read or write.
- FINEXIT** Reinstates a process's soft interrupts. FINEXIT is issued at the end of an interrupt procedure. It will reenable soft interrupts that were automatically disabled on entering the interrupt procedure.

## Creating Message Files

Message files can be created in several ways. When a user process opens a new file and indicates in the *Options* that it will be a message file, the FOPEN intrinsic creates the new message file. In order to create a message file with the :BUILD command, use the MSG keyword. For example, to build a message file named MSG1, enter:

```
:BUILD MSG1;MSG
```

A new message file can be specified as job temporary or default to permanent and may also be defined with a :FILE command. Use the MSG keyword for a new file:

```
:FILE MSG2,NEW;MSG
```

A new message file named MSG2 is indicated.

When you perform a :LISTF,2 command, message files will be identified by an "M" in the third column of the TYP field; MSG1 is identified here:

FILENAME	CODE	-----LOGICAL RECORD-----				----SPACE----			
		SIZE	TYP	EOF	LIMIT	R/B	SECTORS	#X	MX
MSG1		128W	VBM	0	1031	1	258	1	8

Occasionally, you might create a message file and specify a certain number of records for the file to contain, only to discover that the file system has allocated more records than you requested. The reason for this is that the file system is maintaining the necessary internal structure for the message file. The file system has four basic rules for establishing this structure when the message file is created:

1. Open and close records are written to the message file every time a writer process opens or closes the file; therefore, the file system adds two records to the requested number to allow for a minimum of one open and one close operation.
2. The requested number of records is rounded up to fill an even number of blocks.
3. The file system adds an extra block to the message file for the file label to occupy. (This block is transparent to you.)
4. Each extent is the same size; that is, the file system assigns the same number of blocks to each extent.

External IPC

Suppose you want to create a message file named ODDSIZE:

```
:BUILD ODDSIZE;MSG;REC=,3;DISC=51,8
```

You have specified a message file with fifty-one records, three records per block, that occupies eight extents. The file system will adjust the number of records to conform to the rules for message file structure.

1. The file system adds two records to allow for one open and one close indication; the number of records goes from 51 to 53.
2. The number of records is rounded up to 54 to provide an even number of blocks. With three records per block, 54 records will fill 18 blocks.
3. An additional block is added to the file to accommodate the file label. Now the file contains 19 blocks.
4. The eight extents must all be the same size, so the number of blocks is increased from 19 to 24. Each extent now contains three blocks.

Of the 24 blocks in ODDSIZE, 23 are data blocks and one contains the file label, which is invisible to the user. With three records per block, 23 blocks contain a total of 69 data records.

**NOTE**

*In addition to adjusting the number of blocks in a message file, the file system adds a certain amount of space to each block for "overhead". Six bytes will be added to each record, and four bytes will be added for each block.*

A listf ODDSIZE,2 now looks like:

FILENAME	CODE	-----LOGICAL RECORD-----				----SPACE----			
		SIZE	TYP	EOF	LIMIT R/B	SECTORS	#X	MX	
ODDSIZE		128W	VBM	0	69 3	12	1	8	



## Intrinsics - A Detailed Description

Read this section carefully. Intrinsics manipulating message files do not act in a standard file system fashion.

### FOPEN

**Foptions:** (2:3) *File Type*. Determines the type of the file to create for a new file. If the file is old, this field is ignored.

000 - Ordinary file

001 - KSAM file

010 - Relative I/O file

100 - Circular file

110 - Message file

<b>NOTE</b>
-------------

*The Default Designator FOPTION, bits 10 through 12, offers several choices for default file designators. Any value used other than 0 for "filename" will override the File Type field*

(8:2) *Record format*. Message files are always internally formatted as variable-length record files and these bits will be changed by the file system. However, a message file can appear as a fixed length file to an opener. There is no difference for a writer, but a reader will have the portion of this target area which exceeds the record filled with blanks (for an ASCII file) or zeroes (for a binary file).

00 - Fixed

01 - Variable (default)

**Aoptions:** (3:1) *File Copy.* This feature permits a message file to be treated as a standard sequential file, so it can be copied by logical record or physical block to another file. Setting this bit on causes all the remaining file parameters to have their normal defaults.

- 0 - The file will be accessed in its native mode; that is, a message file will be treated as a message file.
- 1 - The file is to be treated as a standard, sequential file with variable-length records. This allows nondestructive reading of an old message file at either the logical record or physical block level. Only block level access is permitted if the file is opened with write access. These blocks are checked for proper message file format to prevent incorrectly formatted data from being written to the message file while it is unprotected.

**NOTE**

*In order to access a message file in copy mode, a process must have exclusive access to the file.*

(5:2) *Multiaccess mode.* This feature permits processes located in different jobs or sessions to open the same file.

- 00 - No multiaccess. The file system changes this value to "2" to allow global multiaccess.
- 01 - Only intra-job multi-access allowed; this is the same as specifying the MULTI option in a command.
- 10 - Inter-job multi-access allowed; this is the same as specifying the GMULTI option in a :FILE
- 11 - Undefined. If this is specified, the FOPEN will be rejected with an error code of 40: ACCESS VIOLATION.

(7:1) *Inhibit buffering.* For message files, the file system turns this bit off. Readers may open a message file with NOBUF if they are in copy mode; this determines whether they will be accessing the file record by record or block by block. Writers must open message files with NOBUF if they are in copy mode; they will access the file block by block.

- 0 - read by logical record
- 1 - read by physical block

- (8:2) *Exclusive.* The values for this field are the same for any disc file, but they have different meanings for the readers and writers of a message file.

BIT SETTING	USER VALUE	MEANING
01	EXCLUSIVE	One reader, one writer
10	SEMI	One reader, multiple writers
11	SHARE	Multiple readers and writers
00	Default	One reader, one writer (converts to a 01)

**NOTE**

*Multiple accessors to one message file need not be in the same process tree*

- (11:1) *Multirecord.* For message files, the file system sets this bit to 0.
- (12:4) *Access type.* These bits specify whether the user will be a reader or a writer process.
- 0000 - READ access only. The FWRITE intrinsic cannot reference this file. This access type requires both read and write access capability to the file. In the case where the writer is waiting for room to write to the message file a read to file will free the writer's wait by giving free space back to the file. The reader process will also serve the writer by performing the write itself, thereby necessitating read and write access. A process that has opened a file with this access type is a "reader."
- 0001 - WRITE access only. If this is the first accessor to the file and the process has write access capability, then the file's contents are purged. If this is not the first accessor to the file, the file system sets this access type to APPEND. The FREAD intrinsic cannot reference this file. A process that has opened a file with this access type is a "writer."
- 0010 - WRITE SAVE access. The file system sets this to APPEND access.

## External IPC

0011 - APPEND access only. The FREAD intrinsic cannot reference this file. This access type requires append capability to the file. A process that has opened a file with this access type is a "writer."

**Device:** This field is relevant only if this is a new file. The DEVICE field must either be omitted or specify a disc; specification of any device other than a disc opens the device. When this occurs, the file is no longer a message file.

**Numbuffers:** (0:11) Ignored.  
(11:5) Value between 2 and 31; default is 2. This parameter must not exceed the physical record capacity of the file.

**Filesize:** The number of records is rounded up to completely fill the last block and to make the last extent the same size as the other extents. Two additional records are included for the open and close records (see page 2-5 to reference the rules the file system uses to allocate the number of records in a file).\*\*

\*\*See the MPE Intrinsic Manual - Accessing and Altering Files.

FCONTROL

A few controlcodes deal specifically with IPC. Those not mentioned here are either invalid or not used with IPC applications.

CONTROLCODE	PARAM	DESCRIPTION
4	integer	Set timeout interval. This applies to both FREADs and FWRITEs. The timeout will be armed at the beginning of the I/O request and cleared when the I/O completes. PARAM specifies the length of the timeout in seconds. A value of zero disables time-outs on the file. This function will remain enabled until explicitly canceled.
6	-	Write End-Of-File. Used only to verify the state of the file by writing out the file label and buffer area to disc; this ensures that the message file can survive system crashes. No actual EOF is written.
43	-	Abort NOWAIT I/O. A CCG condition code is returned if an outstanding I/O operation has completed. An IOWAIT must be issued to finish the request.
45	TRUE	Enable extended wait. Permits a reader to wait on an empty file that is not currently opened by any writer, or a writer to wait on a full file that has no reader. This FCONTROL will remain in effect until FCONTROL 45 is issued with a PARAM value of FALSE.
	FALSE	Disable extended wait. Specifies that when an FREAD encounters an empty file that has no writer, or an FWRITE encounters a full file that has no reader, it will return an end-of-file condition. (Default.)
46	TRUE	Enable reading the writer's ID.
	FALSE	Disable reading the writer's ID. Only data is read to the reader's target area (default).

CONTROL CODE	PARAM	DESCRIPTION
47	TRUE	Nondestructive read. The next FREAD by this reader will not delete the record. Subsequent FREADs will be unaffected and will logically delete records as they are read.
	FALSE	The next FREAD by this reader will delete the record (default).
48	PARAM	<p>Arm/disarm soft interrupts. Param contains the external-type label (plabel) of your interrupt procedure. In SPL it is passed as a parameter by placing an "at" sign (@) before the procedure name.</p> <p>If the value of PARAM is 0, the interrupt mechanism is disabled for this file.</p>

**FCHECK**

**FCHECK will return: 151 CURRENT RECORD WAS LAST RECORD WRITTEN BEFORE SYSTEM CRASHED.**

when a record is read following system startup and the "crash bit" has been turned on the record's header. Refer to the information on message files and their header records page xx (3-4)

**FGETINFO**

The value returned in RECSIZE will indicate the user's data record size, and the value returned in EOF will indicate the number of data records, unless an FCONTROL 46 is in effect. Besides data records in a message file there are also two word of header information for each data record written and open and close records for each writer that opens and closes the file. An FCONTROL 46 will return a value in RECSIZE the size of the user's data records, including the two word data record header. The number of records returned in EOF will include open, close and and data records. (See Message Files and header records page 3-2 for further reference.)

The value returned in BLKSIZE reflects the actual blocksize of the file. When the file is created, the blocksize is computed by the following algorithm:

$$\text{BLOCKSIZE} = ((\text{RECORDSIZE} + 3) * \text{BLOCKING FACTOR}) + 2$$

where RECORDSIZE and BLOCKSIZE are in words. For example, with a recordsize of 100 words and a blocking factor of 10, the blocksize would be 1032 words.

**FFILEINFO**

Three values for ITEMVALUE are specifically for use with IPC:

Item #	Type	Description
34	integer	The current number of writers.
35	integer	The current number of readers.
49	logical	Flabel of the user's soft interrupt procedure. A value of zero implies that soft interrupts are not being used.

## External IPC

### FINSTATE

A True/False value turns on or off soft interrupts for the process.

### FINEXIT

Once the interrupt handler begins execution, file system interrupts are turned off. FINEXIT is to be placed at the end of the interrupt routine to reinstate soft interrupts (logical value true).

## Intrinsics Not Allowed

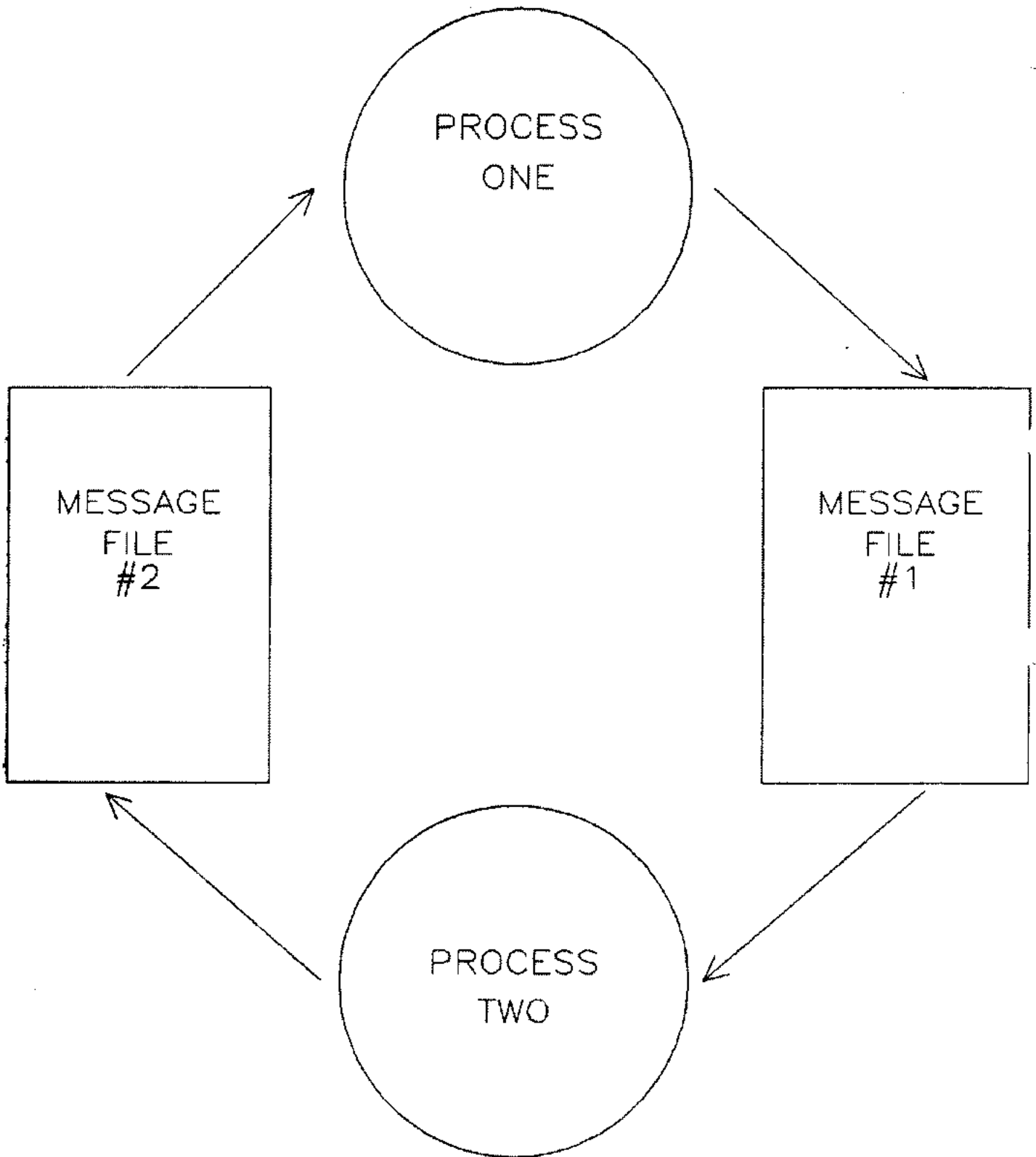
Certain intrinsics are not allowed for message files. (The FSETMODE intrinsic is permitted, but ignored.) Disallowed intrinsics are listed below:

FPOINT	FREADDIR	FDELETE
FREADSEEK	FSPACE	
FUPDATE	FWRITEDIR	



## Examples Using Message Files (Basic IPC)

In this example of IPC there are two processes. The first process reads data from a data file and passes the data to a second process via MSGFIL1 (refer to picture below). After process one finishes, it hangs a read on MSGFIL2 (which is empty) waiting for acknowledgement from process two that it has processed all the data. When process one receives the message from process 2 by completing a read, it terminates.



PROCESS ONE

```
$CONTROL USLINIT
```

```
<< Purpose: >>
```

```
<< Read data from a data file and send to another process. >>
```

```
BEGIN
```

```
LOGICAL EOF := FALSE;
INTEGER DATA'FILE, PIN, IN'FILE, OUT'FILE,ERR;
INTEGER LEN:=80
LOGICAL TRUE:=1;
BYTE ARRAY IN'FILE'NAME (0:8) := "MSGFILE2 ";
BYTE ARRAY OUT'FILE'NAME (0:8) := "MSGFILE1 ";
BYTE ARRAY DATA'FILE'NAME (0:8) := "DATAFILE";
ARRAY MESSAGE (0:39);
```

```
INTRINSIC CREATEPROCESS, FCLOSE, FOPEN, FREAD, FWRITE,
QUITPROG, FCONTROL;
```

```
<< Create entries for the message files in the directory. >>
<< Note that IN'FILE'NAME ("MSGFILE2") is opened with FOPTIONs >>
<< %30004: this indicates a new ASCII message file. >>
```

```
IN'FILE := FOPEN (IN'FILE'NAME, %30004);
IF < THEN QUITPROG (1);
FCLOSE (IN'FILE, 2, 0); << Save file as session temporary. >>
IF < THEN QUITPROG (2);
```

```
<< Note that OUT'FILE'NAME ("MSGFILE1") is also opened with FOPTIONs %30004>>
```

```
OUT'FILE := FOPEN (OUT'FILE'NAME, %30004);
IF < THEN QUITPROG (3);
FCLOSE (OUT'FILE, 2, 0); << Save file as session temporary. >>
IF < THEN QUITPROG (4);
```

```
<<Opening the message file we will write to. >>
<<Note that IN'FILE'NAME ("MSGFILE1") is opened with FOPTIONs >>
<<%106 and AOPTIONS %1100: %106 indicates an old temporary >>
<< ASCII file and %1100 indicates a reader process with >>
<< exclusive access and multiaccess capability. MSGFILE1 >>
<< has already been designated as a message file. Since >>
<< only one reader and one writer process will be accessing >>
<< the message file, exclusive access mode is specified. >>
```

External IPC

```
IN'FILE := FOPEN (IN'FILE'NAME, %106, %1100);
IF < THEN QUITPROG (7);
FCONTROL (IN'FILE, 45, TRUE);
```

```
<< Open message file that we will read from . >>
<< Note that OUT'FILE'NAME ("MSGFILE2") is opened with FOPTIONs >>
<< %106 and AOPTIONS %1101: %106 indicates an old temporary >>
<< ASCII file and %1101 indicates a writer process with >>
<< exclusive access and multiaccess capability. MSGFILE2 has >>
<< already been designated as a message file. Since only >>
<< one reader and one writer process will be accessing the >>
<< message file, exclusive access mode is specified. >>
```

```
OUT'FILE := FOPEN (OUT'FILE'NAME, %106, %1101);
IF < THEN QUITPROG (8);
FCONTROL (OUT'FILE, 45, TRUE);
```

```
<< Open data input file. >>
<< Note that DATA'FILE'NAME ("DATA") is opened with FOPTIONs %3 >>
<< and AOPTIONS 0: %3 indicates an old permanent or temporary >>
<< file and 0 indicates read only access. The file system >>
<< will change the FOPTIONs to specify an ASCII file. >>
```

```
DATA'FILE := FOPEN (DATA'FILE'NAME, %3, 0);
```

```
IF <> THEN QUITPROG (9);
```

```
WHILE NOT EOF DO BEGIN
```

```
LEN := FREAD (DATA'FILE, MESSAGE, -80);
```

```
IF < THEN QUITPROG (10);
```

```
IF > THEN BEGIN
```

```
EOF:=TRUE;
```

```
MESSAGE:= "FOUND EOF IN DATA FILE";
```

```
WRITE (OUT'FILE, MESSAGE, -LEN, 0);
```

```
END<<IF>>;
```

```
ELSE BEGIN
```

```
FWRITE (OUT'FILE, MESSAGE, -LEN, 0);
```

```
IF <> THEN QUITPROG (11);
```

```
END;
```

```
END << WHILE >>;
```

```
FCLOSE (OUT'FILE, 4, 0); << No more data to send: EOF >>
```

```
IF < THEN QUITPROG (12);
```

```
FREAD (IN'FILE, MESSAGE, 1); << Wait for writer to >>
```

```
IF <> THEN QUITPROG (13); << to process this information >>
```

```
<< Because we have issued an FCONTROL 45 against MSGFIL2 >>  
<< process one will terminate normally. If the second >>  
<< process ;however, closes without communicating to >>  
<< process one then we will end in error >>
```

```
FCLOSE (IN'FILE, 4, 0);  
IF < THEN QUITPROG (14);  
END.
```

External IPC

PROCESS TWO

\$CONTROL USLINIT

<< *Purpose:* >>  
<< *Receive data from other process and print it.* >>

BEGIN

LOGICAL EOF := FALSE, TRUE:=TRUE;  
INTEGER LEN, IN'FILE, OUT'FILE;  
BYTE ARRAY IN'FILE'NAME (0:8) := "MSGFILE1 ";  
BYTE ARRAY OUT'FILE'NAME (0:8) := "MSGFILE2 ";  
ARRAY MESSAGE (0:39);  
INTRINSIC FCLOSE, FOPEN, FREAD, FWRITE, QUITPROG, PRINT FCONTROL;

<< *Open message file from other process.* >>  
<< *Note that IN'FILE'NAME ("MSGFILE1") is opened with FOPTIONs* >>  
<< *%106 and AOPTIONS %1100: %106 indicates an old temporary* >>  
<< *ASCII file and %1100 indicates a reader process with* >>  
<< *exclusive access and multiaccess capability. MSGFILE1* >>  
<< *has already been designated as a message file. Since* >>  
<< *only one reader and one writer process will be accessing* >>  
<< *the message file, exclusive access mode is specified.* >>

IN'FILE := FOPEN (IN'FILE'NAME, %106, %1100);  
IF < THEN QUITPROG (13);  
FCONTROL (IN'FILE, %45, TRUE);

<< *Open message file to other process.* >>  
<< *Note that OUT'FILE'NAME ("MSGFILE2") is opened with FOPTIONs* >>  
<< *%106 and AOPTIONS %1101: %106 indicates an old temporary* >>  
<< *ASCII file and %1101 indicates a writer process with* >>  
<< *exclusive access and multiaccess capability. MSGFILE2* >>  
<< *has already been designated as a message file. Since only* >>  
<< *one reader and one writer process will be accessing the* >>  
<< *message file, exclusive access mode is specified.* >>

OUT'FILE := FOPEN (OUT'FILE'NAME, %106, %1101);  
IF < THEN QUITPROG (14);  
MESSAGE := "ERROR"  
WHILE NOT EOF DO BEGIN  
LEN := FREAD (IN'FILE, MESSAGE, -80);  
IF < THEN QUITPROG (15);

```
IF > THEN EOF= TRUE
    ELSE PRINT (MESSAGE, -LEN, 0);
END << WHILE >>;
```

```
<< Because an FCONTROL 45 has not been issued a CCL is generated >>
<< when process two tries to read past the EOF. >>
```

```
<<Now signal other process we are done.>>
```

```
Message="INFORMATION RECEIVED"
FWRITE(OUTFILE,MESSAGE,-LEN,0)
FCLOSE (OUTFILE, 4, 0);
IF < THEN QUITPROG (16);
FCLOSE (INFILE, 4, 0);
END. <<Process Two>>
```

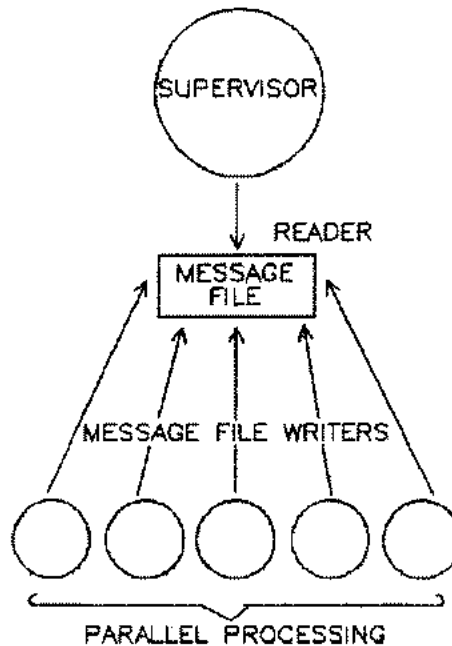
## Software Interrupts

Software Interrupts is an enhancement to BASIC IPC. Soft Interrupts allow the same extended wait conditions when reading or writing to a message file as BASIC IPC. When the I/O completes your program is interrupted and executes a predetermined interrupt procedure. The interrupt handler will complete the read or write and return your program to the line that was executing when the interrupt occurred.

**NOTE**

*Soft Interrupts cannot be implemented in Cobol.*

Typically a soft interrupt application is one in which a master process serves as a controller to parallel processes. One message file is used for "system" communication where the master is a reader and the slave processes are writers to the message file. The master process will usually wait for communication from other processes. It waits by hanging a read on the empty message file. It will be "soft interrupted" by one of the slave processes writing to the message file. The master process parses the command it received in the message and processes the command with a subroutine. The master process replies if necessary through the individual message file of one of its slaves.





**OPERATION**

Initially, software interrupts are disabled for your program. To enable soft interrupts, use the FINTSTATE intrinsic with a value of TRUE, as follows:

```
VALUE:=FINTSTATE(TRUE);
```

The FINTSTATE intrinsic with a value of FALSE will inhibit soft interrupts and VALUE is the old state of FINTSTATE. MPE will inhibit soft interrupts just before entering an interrupt procedure. This is done to prevent unwanted nesting of the interrupt procedures. Using the FINTEXT intrinsic also with a value of true at the end of the interrupt procedure, it will re-enable interrupts just before it exits.

```
PROCEDURE INTERRUPTPROC (FILENUM);
  VALUE FILENUM;
  INTEGER FILENUM;
  BEGIN
  .
  .
  .
  FINTEXT (TRUE);
  END;
```

**NOTE**

*Software interrupts are automatically inhibited just before a CONTROL-Y trap procedure. The trap procedure may elect to allow soft interrupts by calling the FINTSTATE intrinsic. If it does not call FINTSTATE, the RESETCONTROL intrinsic will restore the process' interrupt state to its pre-CONTROL-Y value.*

When you have enabled software interrupts for your program, you must also "arm" soft interrupts for a particular file by specifying the interrupt procedure's label in an FCONTROL 48.

**FCONTROL (File Number, 48, @Procedure Name)**

Calling FCONTROL 48 with a parameter of zero will disarm the software interrupt mechanism so that the file would be accessed by Basic IPC procedures.

**NOTE**

*The FFILEINFO intrinsic may be used to return the label of the interrupt handler. FFILEINFO 49 will return the label as an integer value: if it returns a value of zero, no interrupt handler has been armed.*

After the program interrupt has occurred, an IODONTWAIT must be issued against the file to complete the request. Your interrupt handling procedure should issue the IODONTWAIT intrinsic. IODONTWAIT will allow completion operations for the associated I/O or return a zero for an incomplete I/O operation. Multiple pending soft interrupts are queued in a FIFO manner. Open files with a soft interrupt pending (read or write) cannot be closed.

**NOTE**

*The message file and the interrupt must occur on the same system. Soft interrupt use is restricted to message file and may not be used with remote files.*

No more than one uncompleted FREAD or FWRITE may be outstanding from one process for one particular message file. Any additional FREAD's or FWRITE's will result in an unequal condition code being returned. An uncompleted FREAD or FWRITE request; however, may be aborted by issuing an FCONTROL 43 (abort nowait I/O).

An interrupt will not occur while you are executing within operating system code; that is, while you are processing an MPE intrinsic or procedure. Exceptions to this are the PAUSE, PAUSEX, and IOWAIT intrinsics that will allow the processing of a soft interrupt. MPE reinvokes these intrinsics when the interrupt procedure has completed. If the PAUSE intrinsic is soft interrupted two minutes into its overall time, after the soft interrupt has finished processing, it will be reinitialized to its original time. PauseX, however, was created to be used for soft interrupts. It will retain the amount of time it has used and continue after an interrupt, the remaining unused interval. If the PAUSEX was issued for five minutes and was interrupted again two minutes after it had begun, PAUSEX would continue to pause after the interrupt procedure for the remaining three minutes.

The validity of an interrupt procedure depends on the code domain and executing mode (privileged or non-privileged) of your code. It must be compatible with the code domain and the mode (privileged or non-privileged) of your interrupt procedure. The code domains are:

PROG (User program)  
 GSL (Group SL)  
 PSL (Public SL)  
 SSL (System SL, non-MPE segments)  
 MPSSL (System SL, MPE segments)

When the code of the caller is:	The label:
Non-privileged in PROG, GSL, or PSL.	Must be non-privileged in PROG, GSL, or PSL.
Privileged in PROG, GSL, or PSL.	May be privileged or non-privileged in PROG, GSL, or PSL.
Privileged or non-privileged PROG,GSL,PSL	May be in any non-MPSSL segment.

## Soft Interrupt Example

The following is a skeleton example of a program which is using soft interrupts. The reader process sets up the soft interrupt and then initiates a read on an empty message file. That read request will be queued. The writer process once created, writes to the message file and releases the reader. The reader continuing to execute mainline code will be interrupted and process the interrupt with procedure INTERRUPTPROC. Once INTERRUPTPROC completes, the reader will continue at the interrupted instruction in mainline code.

### READER PROGRAM

```

$CONTROL
BEGIN
ARRAY BUFFER(0;79);
INTEGER TRUE:=1, PLABEL, FILENUM, TCOUNT;
INTRINSIC FINSTATE, FCONTROL, FINTEXTIT, INDONTWAIT,
FOPEN, FREAD;
PROCEDURE INTERRUPTPROC(FILENUM); ← STEP 1
  VALUE FILENUM;
  INTEGER FILENUM;
  BEGIN
    IODONTWAIT(FILENUM,BUFFER,TCOUNT); ← STEP 2
    PRINT BUFFER TO DISPLAY ON THE TERMINAL;
    FREAD(FILENUM,BUFFER,TCOUNT); ← STEP 3
    FINTEXTIT(TRUE); ← STEP 4
  END;

BEGIN <<INITIALIZATION>>
  FINSTATE(TRUE); ← STEP 5
  FILENUM:=FOPEN(MSGFILE); ← STEP 6
  FCONTROL(FILENUM,48,PLABEL); ← STEP 7
  FREAD(FILENUM,BUFFER,TCOUNT); ← STEP 8
  CREATE(WRITER PROCESS);
  <<END OF INITIALIZATION>>

LOOP: <<BEGIN MAIN LINE>>

  PERFORM OTHER PROCESSING; ← STEP 9

  END; <<END OF MAIN LINE>>
END.

```

**INTERRUPT PROCEDURE FUNCTIONS:**

STEP 1 This is the beginning of the interrupt procedure which handles soft interrupts. Once the READ has started to complete this procedure will be executed.

**NOTE**

*Once the soft interrupt has occurred, the read has not actually transferred any data. A call to IOWAIT or IODONTWAIT is required to complete the read.*

STEP 2 IODONTWAIT or IOWAIT is needed to complete the read.

STEP 3 Call FREAD again to initiate another soft interrupt read.

STEP 4 FINTEXTIT will re-enable soft interrupts before exiting the interrupt procedure (possibly to process a pending soft interrupt).

**MAIN LINE FUNCTIONS:**

STEP 5 FINSTATE(TRUE); will enable soft interrupts to occur at the program level (only needs to be initiated once).

STEP 6 FOPEN the new file with following options:

A. Message file

B. Reader

**NOTE**

*The first message file opener must open the new message file twice. When the message file is closed for the first time, it will be put in the directory at that time. This would be needed only when more than one process will be accessing the message file.*

External IPC

- STEP 7 FCONTROL 48: This will enable soft interrupts for the message file specified in "FILENUM."
- STEP 8 This FREAD for FILENUM will initialize the first soft interrupted read.
- STEP 9 The reader process is free to perform other necessary processing.

## WRITER PROGRAM

```

$CONTROL
BEGIN
ARRAY BUFFER(0:79):=" A SOFT INTERRUPT OCCURED";
INTEGER FILENUM;
INTRINSIC FOPEN,FWRITE
BEGIN
  FILENUM:=FOPEN(MSGFILE); ←————STEP 1
LOOP:
  PAUSE(5 MINUTES); ←————STEP 2
  FWRITE(FILENUM,BUFFER,TCOUNT,%60); ←————STEP 3
  GOTO LOOP;
END;
END.

```

STEP 1 FOPEN the file with following options:

A. Message file

B. Writer

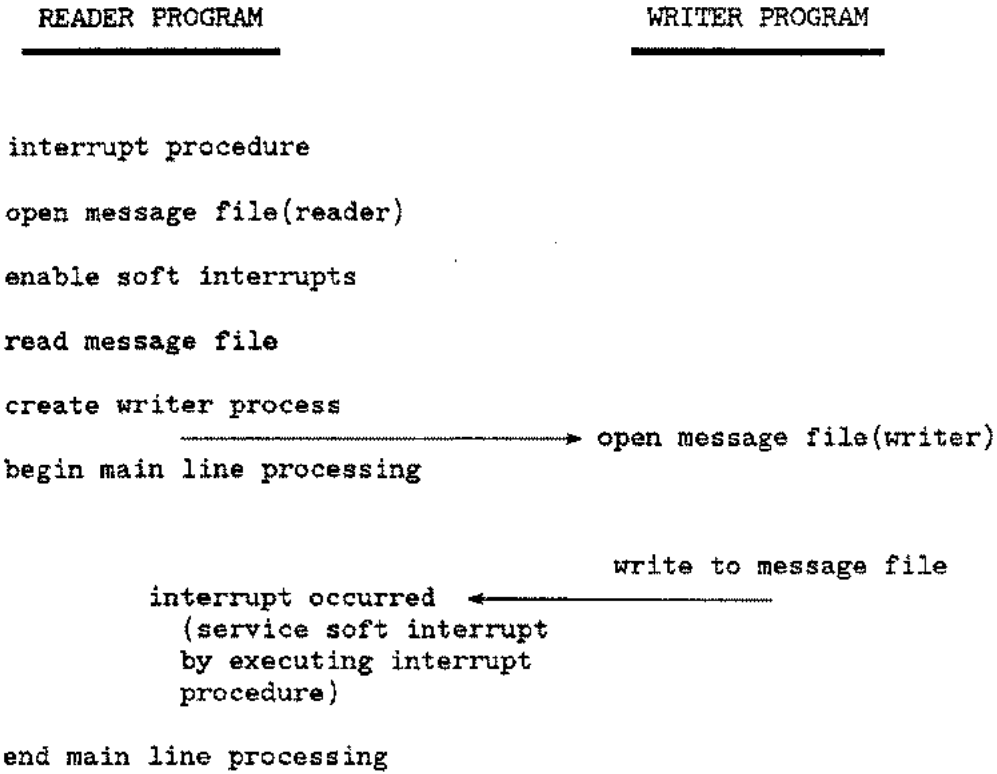
STEP 2 PAUSE for five minutes which will insure that the read posted by the reader program will be interrupted by a soft interrupt.

<b>NOTE</b>
-------------

*The pause is not needed in an actual application design. It is used only to delay the soft interrupt from occurring the instant the soft interrupt read is posted.*

STEP 3 Write to the message file causing the soft interrupt to occur on the reader process.

FUNCTIONAL FLOW



The READER program executes main line code until the writer program writes a message to the message file. Every five minutes the WRITER program writes a record and the READER program soft interrupts, reading the message.



# MESSAGE FILES AND THE FILE SYSTEM

SECTION

III

File System  
Callable  
Intrinsics



File System  
&  
IPC  
Access  
Procedures



Uncallable  
IPC  
Procedures

*IPC PROCEDURES ARE ACCESSED THROUGH STANDARD FILE SYSTEM INTRINSICS*

IPC was written sharing as much MPE File System code as possible ergo it uses many standard file system intrinsics with either added options or specially defined for manipulating message files. This section is essentially an explanation of how IPC works in conjunction with MPE's file system. IPC data structures and dynamics are added to the fundamental file system descriptions.

## MESSAGE FILES

Message files since they are the fundamental building blocks of IPC. They are the communication device between processes. Message files differ from standard files, because within the file block the data records lie opposite from the header information.

### FILE BLOCK

```
..... *****  
: First data record :  
: ..... : Same format as standard  
: Second data record : variable length blocks.  
: ..... :  
z z  
: ..... :  
: Last data record :  
: ..... :  
: Record delimiter (-1) :  
: ..... *****  
: Available space :  
: ..... :  
: Header delimiter (%77) :  
: ..... :  
: Last header record :  
: ..... :  
z z  
: ..... :  
: Second header record :  
: ..... :  
: First header record :  
: ..... :
```

Data is written at the top of the block while its header information is written at the bottom. A negative one delimits data records and a %77 delimits the headers.

### DATA RECORDS

Data records in a message file are always variable length records regardless of how you may have opened the file. Each record has a one word header that prefixes the record with a byte count of the records length.

```
.....  
: Number of bytes in record :  
:.....  
: First data word of record :  
:.....  
z ..... z  
:.....  
: Last data word of record :  
:.....
```

### Record Format

## HEADER RECORDS

The header records are built from the bottom of the block toward the top. There are three types of header records and they are written by a writer upon different events. An "Open" header record is written with the first write to the file. The "close" header record is written when an FCLOSE is issued, and a "Data" header record is written each time the writer writes to a message file. Header information is two words long. The type of header record is contained in the first word (8:15).

```
.....  
: C:LC:           : Header Type: 0  
.....  
: Writer's ID*    : -1  
.....
```

### Header Format

C (0:1)- Set on if this was the last record written before the system crashed. An FCHECK issued will return:

"Current Record Was Last Record Written Before System Crash

LC (1:1)- Valid only for close headers. Set to one if this was the last writer to close the file.

Type(8:8)- 0 data  
          1 open  
          2 close

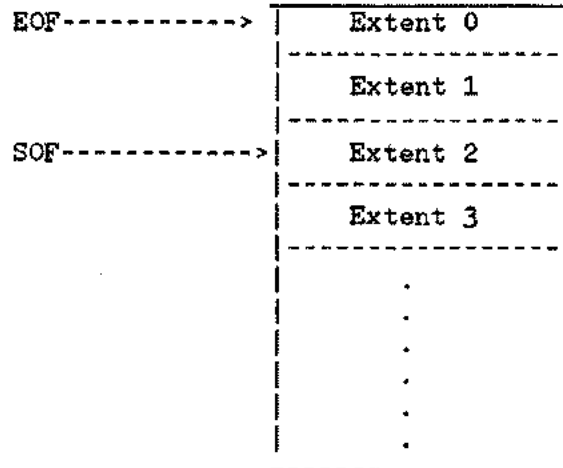
\*Recall that the writer's ID may be read with an FCONTROL 46.

**FILE EXTENTS**

Observe the message file below and note that the Start Of File is greater than the End Of File. Message files are circular files. Data written to a message file can wrap-around the file as we use or reuse allocated extents; therefore the figure below would not be an unusual one to see.

File extents are composed of a number of blocks. Extent zero also contains space for the file label and user labels in the exact same format as a standard file. Starting with block zero, sufficient blocks are allocated to the file label/user labels to satisfy their space requirements. The start of file moves into the extent map of a File Control Block<sup>2</sup> as records are read; hence records are logically deleted. The EOF moves when records are written.

Extents outside of the SOF/EOF range may not exist. They are deleted at close time when there are no more writers accessing the file.

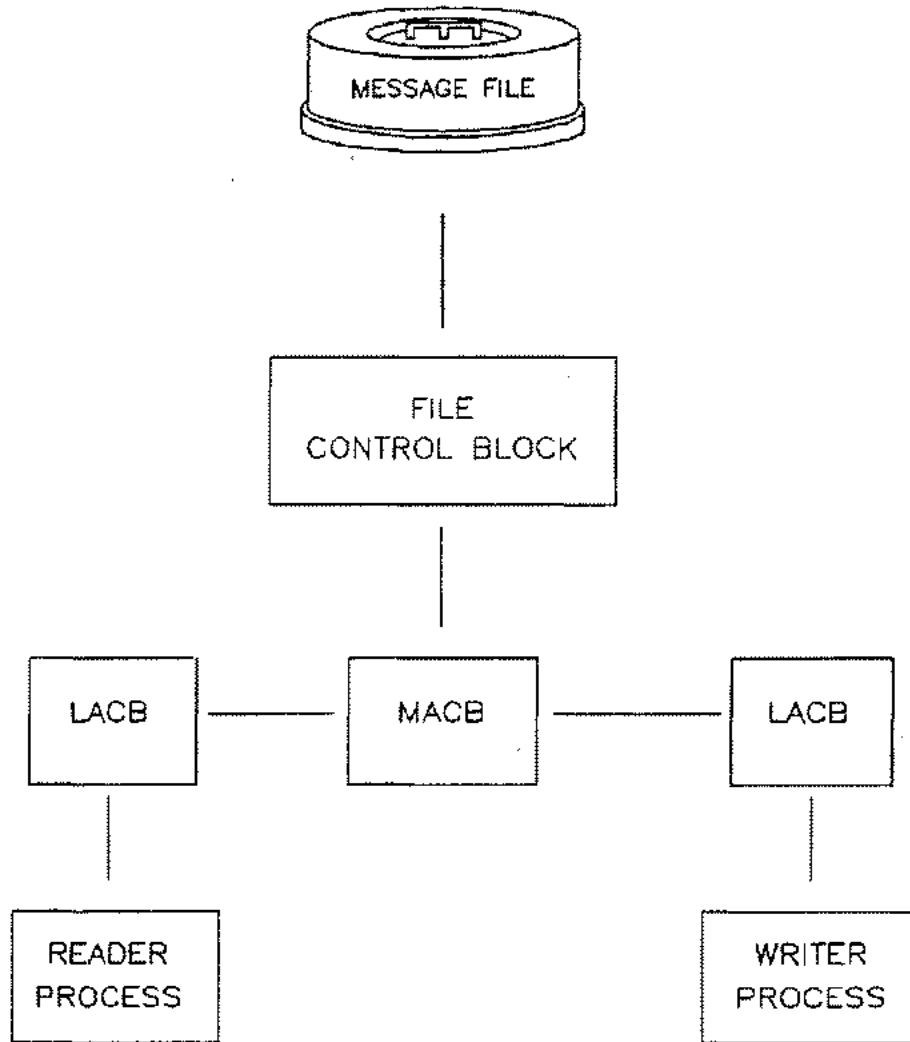


<sup>2</sup>See File Control Block in next section page 3-15.

## CONTROL BLOCKS

File access is coordinated by control blocks. Control blocks contain information pertaining to a file. Control blocks are dynamic; they are created when the file is opened (FOPEN); they are consulted when a file is accessed (FREAD, FWRITE); and pointers to these control blocks are deleted when file access is terminated (FCLOSE). A message file has three control blocks (excluding control block extensions): a **Logical Access Control Block** and its extension, **Message Access Control Block** (corresponding to a standard file's physical access control block) and its extension, and **File Control Block**.

WHAT ARE THE RELATIONSHIPS BETWEEN  
FILE --- CONTROL BLOCKS --- PROCESS



The figure above shows how we logically look at the message file and its control blocks. Message files have multi - access capability which permits the sharing of the access to a file. Each user shares the same SOF and EOF. File access information which is unique and local to each individual accessor is contained in the LACB. File access information which is common to each set of accessors is contained in the MACB. The MACB contains an FCB vector<sup>3</sup>, a pointer to the FCB and the FCB contains the file's extent map.

<sup>3</sup>See vector structure page 3-21.

## Logical Access Control Block

The logical access control block contains the foptions and aoptions used by a process to open the message file. It also contains the file number (returned by the FOPEN intrinsic), file name, record size and block size. Any file in our file system opened multi has an LACB. A message file's LACB is the same as a standard file's LACB.

```
.....  
:      : Size of the ACB including buffers (words)      : 0  
.....  
:      : File number                                     : 1  
.....  
: File name                                             : 2  
.....  
z      z  
.....  
: Foptions                                             : 6  
.....  
: Aoptions                                             : 7  
.....  
: Record size (bytes)                                  : 10  
.....  
: Block size (words)                                   : 11  
.....  
z      z  
.....  
: Carriage control code (writers)                     : 13  
.....  
z      z  
.....  
: Error code                                           : 16  
.....  
: Transmission log (units same as last read/write) : 17  
.....
```



## Logical Access Control Block Extension

The LACB extension contains information specific to IPC and soft interrupts. Both the LACB and the LACB extension are created at FOPEN time. These control blocks are not shared with other processes but are unique to the process that has created them. These control blocks are found in the process's own stack in the PXFILE area<sup>4</sup> (unless that process is running its message file application NOCB or there is no room remaining in the area- then both of these control blocks will be found in extra data segments).

.....	:	User's soft interrupt plabel	:	20
.....	:	Number of seconds to wait on boundary condition	:	21
.....	:	O:Ex:Nd:Vr:Bt:Cls :C : Carriage control	:	22
.....	:	Reply Port (basic IPC port)	:	23
.....	:	Writer ID	:	24
.....	:	Control block index for nowait writer record buf	:	26
.....	:	DST relative addr of nowait writer record buffer	:	27
.....	:	No wait I/O resultant error code	:	30
.....	:	No wait I/O resultant transmission log	:	31
.....	:		:	

<sup>4</sup>Where is the PXFILE area? See File Access later in this section.

## Message File Internals

Word	Field	Description
22		Accessor's local flags.
(0:1)	D 1	- have not yet issued an FREAD/FWRITE against the file.
(1:1)	ex 1	- extended wait mode.
(2:1)	nd 1	- do not destroy the next record read.
(3:1)	vr 1	- writer has not yet written his first record (ie., he is a virgin).
(4:1)	bt 0	- transmission log should be expressed in words.
	1	- " " " " " " bytes.
(5:1)	cls	- Not currently used (reserved for group IPC standard).
(6:1)	C	- No wait completion message is in LACB area.
(8:8)	car	- carriage control character to be used for the writer's record (a value of one indicates no carriage control character).
	ctl	

## Message Access Control Block

The MACB is a Physical Access Control Block<sup>5</sup> for message files. It is considered the "point of attachment" (from the process) to the disc file. It contains the "start of file" pointer for the reader/s and "end of file" pointer for the writer/s. The MACB is created by the first process to open the file. According to file system placement strategy the MACB and its extensions (they are contiguous) are placed in an extra data segment.

## MACB Extension

This extension (contiguous with the MACB) contains IPC specific information, the location of reader and writer wait queues.<sup>6</sup> Beyond the extension are the message file buffers (also contiguous).

\*Same fields as LACB

```

.....
:      : Size of the ACB including buffers (words)      : 0  *
.....
:      :                               : File number      : 1  *
.....
: File name                               : 2  *
.....
z                                           z
.....
: Foptions                               : 6  *
.....
: Aoptions                               : 7  *
.....
: Record size (bytes)                   : 10 *
.....
: Block size (words)                    : 11 *
.....
z                                           z      *
.....
: Carriage control code (writers)       : 13 *
.....
z                                           z      *
.....
: Error code                             : 16 *
.....
: Transmission log (units same as last read/write) : 17 *
.....
: Total number of unread records (includes opens : 20
.....
: and closes)                           : 21

```

<sup>5</sup>Reference MPE Tables Manual page 6-1.

<sup>6</sup>See Message Queue Mechanisms section 4.

Message File Internals

```

.....
: Block number of the file's tail (relative to the : 22
.....
: start of file block) : 23
.....
: Logical record transfer count : 24
.....
: : 25
.....
: Physical block transfer count : 26
.....
: : 27
.....
: Address of the head record's header : 30
.....
: Address of the next write header : 31
.....
: FCB control block vector : 32
.....
z z
.....
: Number readers : Number readers & writers : 34
.....
z z
.....
: : Records per block : 36
.....
:Wrt buf indx: : # buf - 1 : 37
.....
: Address of the head record's data : 40
.....
: Size of the buffer (words) : 41
.....
z z
.....
: : Logical device number : 46
.....
:0:# rd buf : # wt buf :er :qw :m :c :d :s :f : 47
.....
: Number of max sized free records : 50
.....
: : 51
.....
: Number of free words in the current free record : 52
.....
: Address of the next write record : 53
.....
: Number of nondata records in the file : 54
.....
: : 55
.....
: # of read requests that have a claim on file : 56
.....
: Last read error : Last write error : 57

```

Message File Internals

.....		
: DST number of the physical ACB	:	60
.....		
: Address of the physical ACB	:	61
.....		
: DST number of the logical ACB	:	62
.....		
: Address of the logical ACB	:	63
.....		
: DST rel address of the stack access control blk	:	64
.....		
: DST rel address of the DB area	:	65
.....		
: MACB vector table entry address	:	66
.....		
: MACB control block vector table address	:	67
.....		
: Target area's DST number	:	70
.....		
: Reserved for calling parameters	:	71
.....		
:	:	72
.....		
:	:	73
.....		
: Reserved for the stack marker from file system	:	74
.....		
: intrinsics	:	75
.....		
z		z
.....		
: User's soft interrupt plabel	:	100*
.....		
: Number of seconds to wait on boundary condition	:	101*
.....		
: O:Ex:Nd:Vr:Bt:Cls :C : Carriage control	:	102*
.....		
: Reply Port (basic IPC port)	:	103*
.....		
: Writer ID	:	104*
.....		
: Control block index for nowait writer record buf	:	105*
.....		
: DST relative addr of nowait writer record buffer	:	106*
.....		
: No wait I/O resultant error code	:	107*
.....		
: No wait I/O resultant transmission log	:	110*
.....		

BEGINNING OF MACB EXTENSION

.....	
: Write wait queue (basic IPC port)	: 111
.....	
: Read wait queue (basic IPC port)	: 112
.....	
: Head record's length (bytes)	: 113
.....	
: Head record's record type (same values as header)	: 114
.....	
: Head record's writer ID	: 115
.....	
: Head record's header word value	: 116
.....	
: Max size record plus its overhead (words)	: 117
.....	
: ACB wait queue message - contains same info as	: 120
.....	
: the wait queue message in the Message Queue	: 121
.....	
: Entry	: 122
.....	
:	: 123
.....	
:	: 124
: Waiter's reply port, 0 if using ACB compltn area	: 125
.....	
: ACB completion message area - *see Message Queue	: 126
.....	
: Entry for completion message format	: 127
.....	
: Waiting process' pin	: 130
.....	
: Waiting process' file number	: 131
.....	
: Waiting process' soft interrupt plabel	: 132
.....	
: DST rel address of buffer one	: 133
.....	
: DST rel address of buffer two	: 134
.....	
: Etc.	: 135
.....	

\* Value is private to a particular accessor.

Word	Field	Description
47		File's global flags.
	(9:1)	er 1 - extended read
	(10:1)	qw 1 - one or more writers has been queued on the wait queue.
	(11:1)	m 1 - wait msg is located in the ACB
	(12:1)	c 1 - completion msg is located in the ACB
	(13:1)	d 1 - the current write buffer has dirty bit set
	(14:1)	s 1 - the start of file is block zero
	(15:1)	f 0 - the ACB buffers have not been filled
102		Accessor's local flags.
	(0:1)	G 1 - have not yet issued an FREAD/FWRITE against the file.
	(1:1)	ex 1 - extended wait mode.
	(2:1)	nd 1 - do not destroy the next record read.
	(3:1)	vr 1 - writer has not yet written his first record (ie., he is a virgin).
	(4:1)	bt 0 - transmission log should be expressed in words. 1 - " " " " " " bytes.
	(5:1)	cls - Not currently used (reserved for group IPC standard).
	(6:1)	C - No wait completion message is in LACB area.
	(8:8)	car - carriage control character to be used for ctl the writer's record (a value of one indicates no carriage control character).

## File Control Block

The file control block data structure and placement strategy is the same as that for standard files. Information used by the file system is copied from the file label into the FCB. The FCB is created with the first opener of the file and when more than one process has opened the message file, the FCB will reside in a control block table in an extra data segment or Shared System Control Block Table (See MPE Tables Manual, The File System - Control Block Area). The file label and MACB have pointers to the File Control Block called the FCB vector (see "vector format" under Active File Table later in this section).

The FCB has the following format:

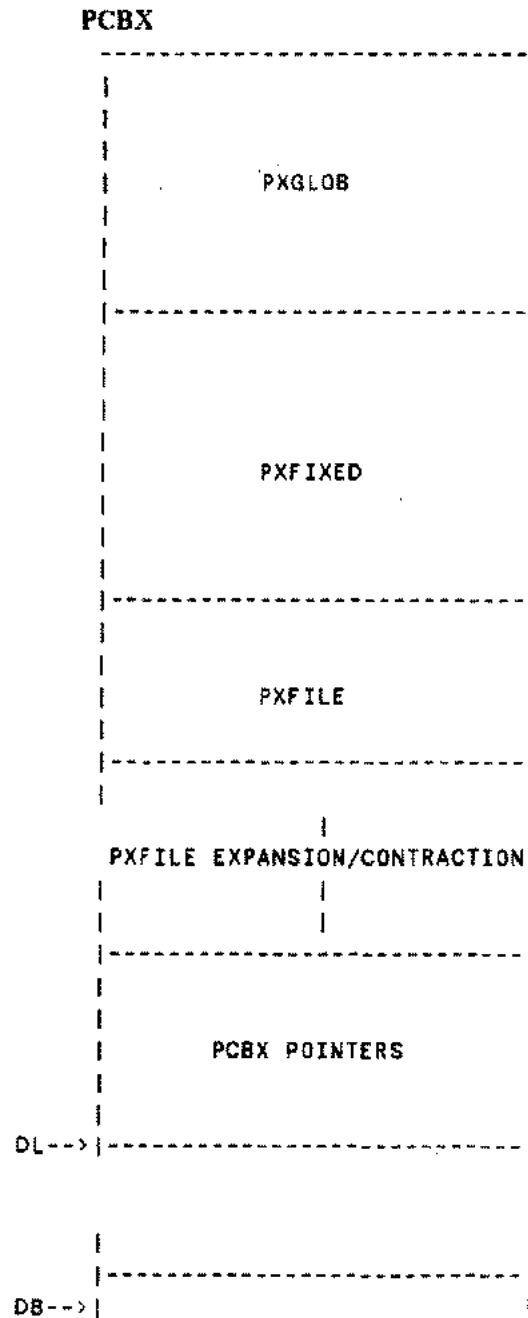
0	1	2	3	7	8	12	13	14	15	
1   Complete FCB size										0
New FCB vector										1
FOPTIONS										2
Device specification										3
Prev. lock		Dev. type		C		Device subtype				4
No. opens for output					No. opens for any mode					5
Creator ACB vector										8
RIN number										7
Exclusive status										10
Private volume information										11
File limit										12
Reserved for IMAGE										13
Reserved for IMAGE										14
Reserved for IMAGE										15
End of data pointer										16
End of data pointer										17
No. user labels written					No. user labels avail.					20
Extent size in sectors										21



Blocking factor	Sectors per block	22
Sector offset to data	Disp   No. extents - 1	23
Last extent size in sectors		24
	No. opens input mode	25
Group name - 1st char.	Group name - 2nd char.	26
Group name - 3rd char.	Group name - 4th char.	27
Group name - 5th char.	Group name - 6th char.	30
Group name - 7th char.	Group name - 8th char.	31
Acct name - 1st char.	Acct name - 2nd char.	32
Acct name - 3rd char.	Acct name - 4th char.	33
Acct name - 5th char.	Acct name - 6th char.	34
Acct name - 7th char.	Acct name - 8th char.	35
Start of file block number		36
		37
Current number of data blocks in the file		40
		41
Number of open and close records (message file)		42
		43
Logical device number		44
First extent sector number		45
Logical device number		
Last extent sector number		

## Control Block Access

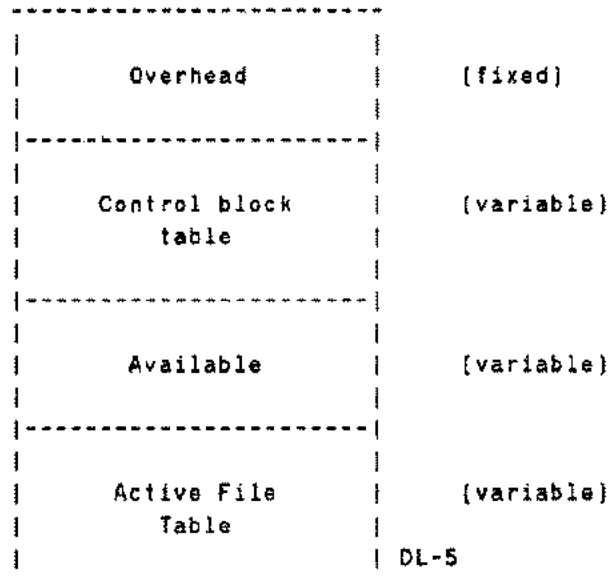
This is PCBX area (Process Control Block Extension) that sits below DL in the process's stack. The PXFILE area lies within the PCBX area and control blocks are contained in the PXFILE area. The PCBX is used to hold process specific information that can be swapped out with the stack.



PXFILE

The PXFILE area is a contiguous, expandable block of storage that is managed by the file system.

The overall structure of the PXFILE area is:



Available Block

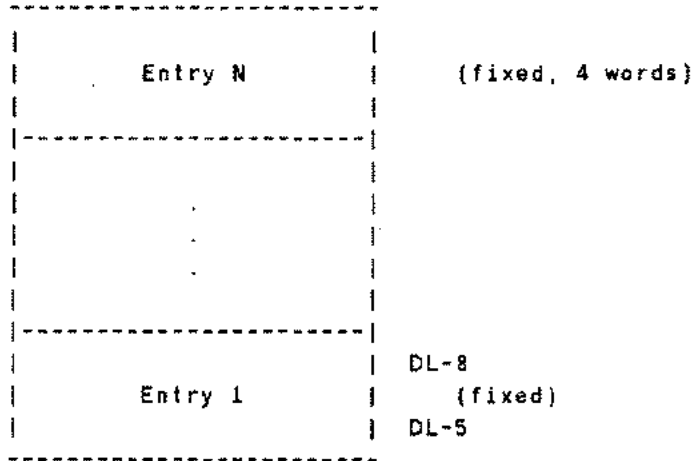
---

The part of the PXFILE area labeled "Available" is used to provide space when the Control Block Table or the Active File Table is expanded. If the Available area is exhausted, the entire PXFILE area is expanded, the AFT is relocated and new space is added to the Available Block. Currently the PXFILE area is only expanded; it is never contracted.

Active File Table (AFT)

The Active (or Available) File Table contains information used by the file system to coarsely characterize file access and, most importantly, to give the location of the file's control blocks.

The overall structure of the AFT is:



The AFT is negatively indexed by file number: the entry at DL-5 corresponds to file number 1, the entry at DL-9 corresponds to file number 2<sup>7</sup>, etc.

<sup>7</sup>For sessions, file #1 and #2, refer to \$STDIN and \$STDLIST respectively

The structure of a file system AFT entry is:

0	1	2	3	4	5	15	
* Entry type						<Not applicable to Msg Files>	0
Physical ACB Vector (MACB for Message Files)							1
Logical ACB Vector (LACB)							2
NO-WAIT I/O IOQX							3

\* Message Files have an entry type %10.

The file system will find file control blocks (e.g. during an FREAD or FWRITE) by negatively addressing into the Available File Table and using the appropriate vector to locate a control block.

The format for a vector; LACB, PACB or FCB (found in the PACB vector) is:

0	5	6	15
Entry No	DST number		
in File	where Control Block		
Control Block	is located		
Vector Table			

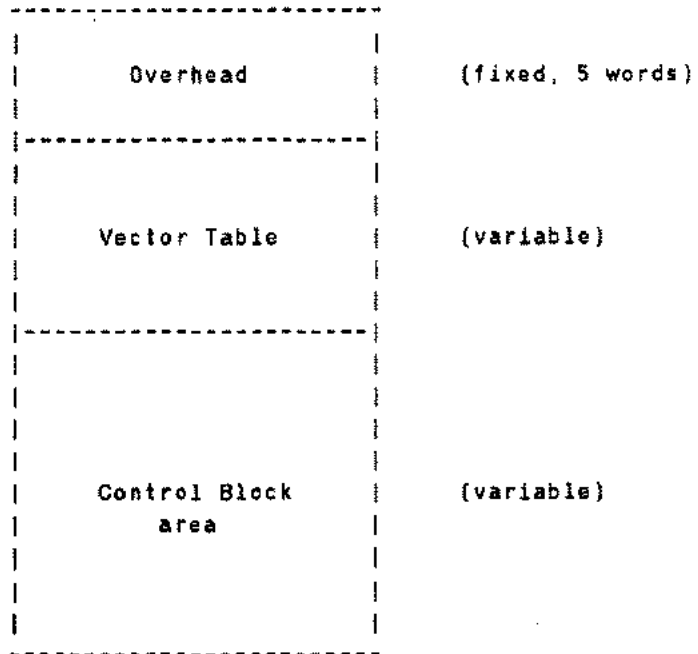
## Message File Internals

### Control Block Table

-----

The control block table holds either one control block or multiple control blocks. A control block table may either reside in the process's stack in the PCBX area or in an extra data segment (see file system narrative in the MPE Tables Manual File System; section 3.2 File Control Block Table).

The overall structure of a control block table is:

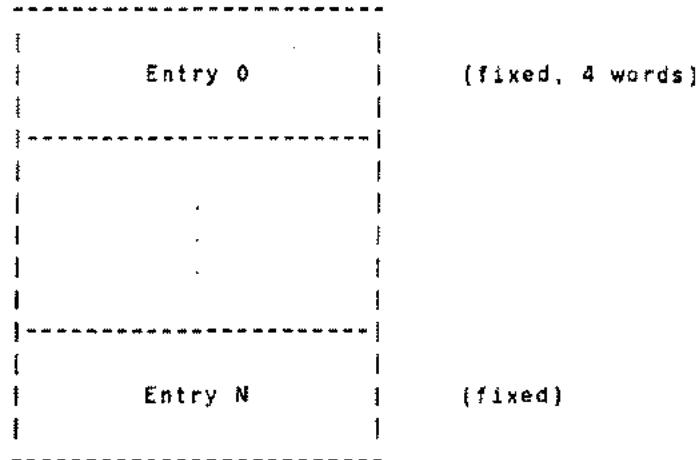


Vector Table

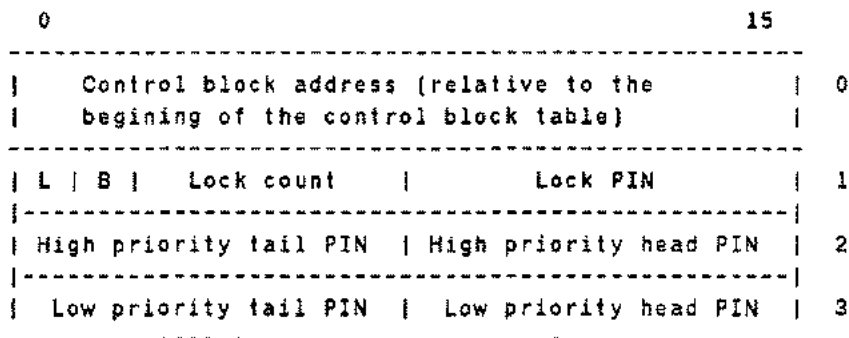
-----

Vectors in the AFT point to vector table entries in the Control Block Table. Vector Table entries contain the addresses of the control blocks (\*\*addresses are relative to the beginning of the Control Block Table) contained in the control block table. The file system also uses the fields in a vector table entry to lock or unlock control blocks in the control block table while they are being accessed.

The overall structure of the vector table is:



(An unused vector table entry will have zeroes in all the words of the entry.)  
 The structure of a vector table entry is:



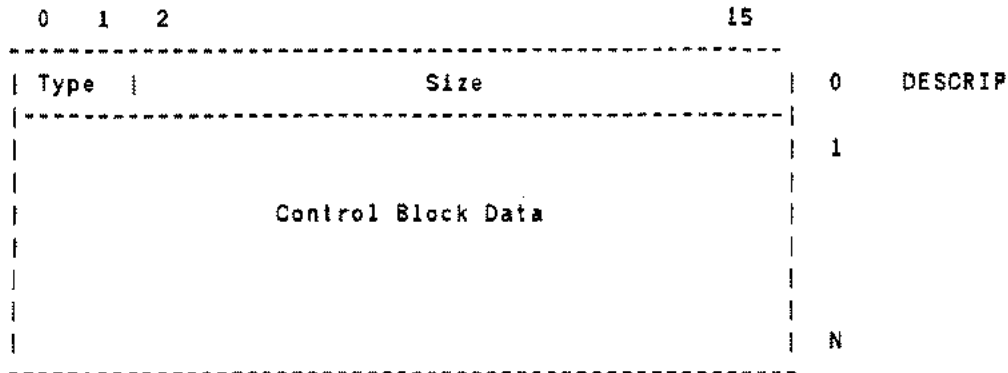
The lock pin field will have the pin number of the process that has locked the control block. Control blocks are locked while they are being accessed so no other processes may write to them. The High Priority head and tail pin in word two are the head and tail of processes that are waiting to access that control block and are impeded. The remaining links for this queue are the NIMP pin (next impeded process) and PIMP pin (previous impeded pin) found in the PCB entry for that process.

## Message File Internals

### Control Block Area

-----

The part labeled CONTROL BLOCK AREA contains the control blocks used by the file system. To facilitate storage management, all control blocks have the same overall structure:



**DESCRIP** This is the first word of a control block; the format is common for all control blocks.

**SIZE** This is the size (in words) of the control block. The size includes the descriptor word.

**TYPE** This is the type number of the control block. There are four types of control blocks:

- 0 - Garbage
- 1 - FCB
- 2 - PACB
- 3 - LACB

When a control block table is created, the initial control blocks are completely allocated to a single control block area of type garbage. When space is requested for a new control block, the control blocks are scanned (using a first fit algorithm) for a garbage control block that is as large as the size requested. The space for the new control block is taken from this garbage control block and the space remaining belongs to the new garbage control block size.

If space is requested and no garbage control block is large enough to contain the new control block, then the control block area and control block table are expanded by a sufficient amount. If expansion is not possible, some other control block table must be used.



## REVIEW

When a file system intrinsic is executed, the file access procedures need to get at file information. Information is sought by the file number which tells the file system to:

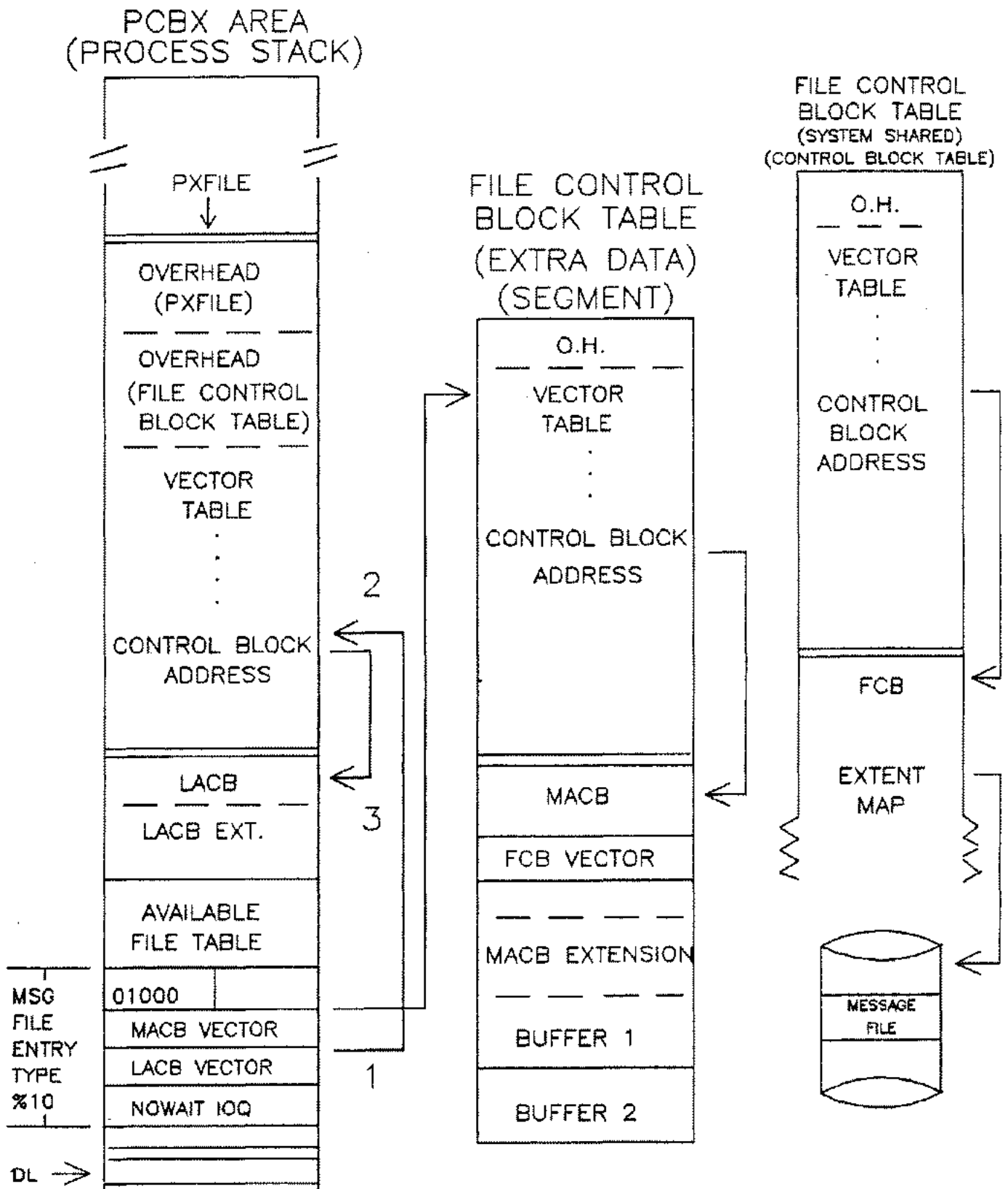
1. Negatively address into the Active File Table (4 x File Number).
2. Find the vector it needs for either an LACB or PACB (MACB) in the AFT entry. The vector will give a DST# where the vector table is contained and an index or entry into the vector table.
3. The vector table entry is an address relative to the beginning of the Control Block Table where the control block may be found.<sup>8</sup>

*(see picture on the following page)*

---

<sup>8</sup>See Appendix B for a dump example of finding control blocks.

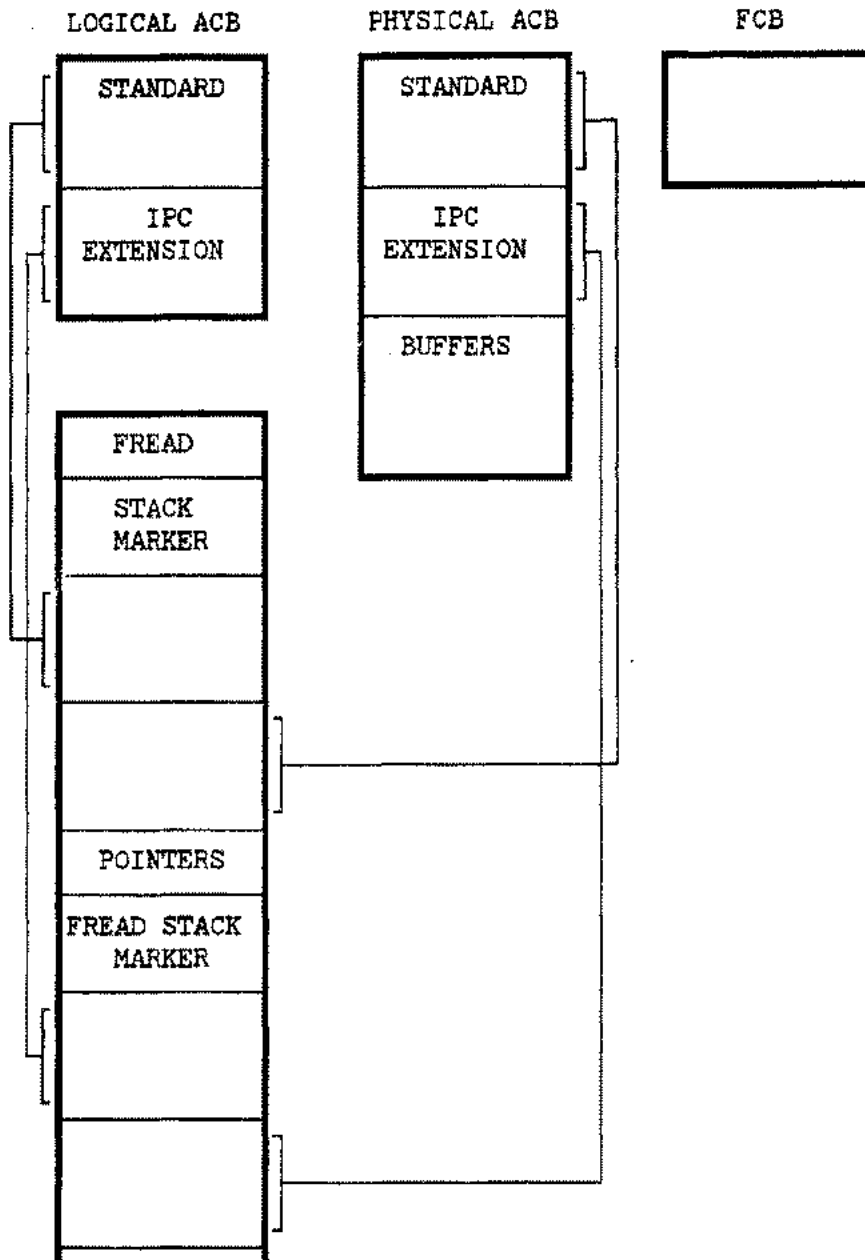
# CONTROL BLOCK ACCESS



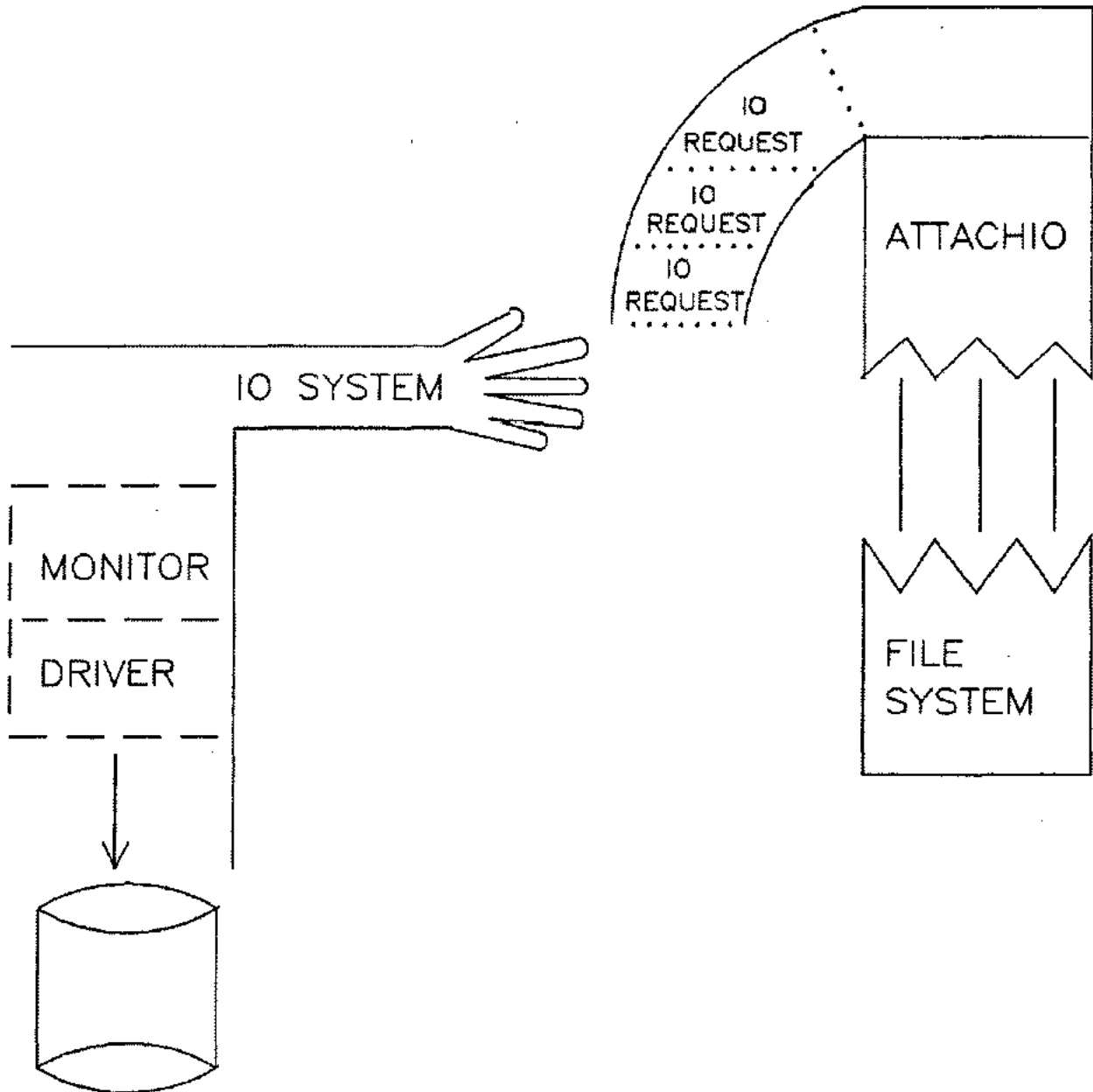
## File Access

To read or write to a message file, the file system must be able to obtain the information in the file's control blocks. Instead of doing several exchange DB's to get at different data segments that contain control blocks, the FREAD and FWRITE intrinsics lay the control blocks on the stack. You may have noticed already that the MACB contains the LACB ( words 0-%17) and LACB extension ( words %100-%110). When the MACB is laid on the stack it is overlaid by these two other control blocks unique to the caller of that intrinsic. We refer to these control blocks when they are on the stack as the Combined Access Control Block.

ACB - COMBINED ACCESS CONTROL BLOCK



When the control blocks are on the stack and accessible, then a call to the I/O system can be made for the physical I/O to be performed. Attachio is the procedure that bridges the file system and the I/O system. Some of the parameters it passes are function: read or write etc.; a byte count to be transferred; and a destination or source of the transfer. The I/O system creates a request for the I/O and puts the request on the IOQ until the driver can complete the request.



## Buffer Management and Anticipatory Reading

Message files support anticipatory reads. When the record number is passed in an FREAD, the file system uses the information in the complete ACB to calculate what block contains the record<sup>9</sup> and calculates which buffer to use for that block. The file system then makes a request for the transfer of the block containing that record. When the process "waits" for the record (with IOWAIT or IODONTWAIT) then the record is transferred from the buffer to the target DST. When a process writes to a message file the buffer is not "flushed" -- records are not actually transferred to the disc file until the buffer becomes filled.

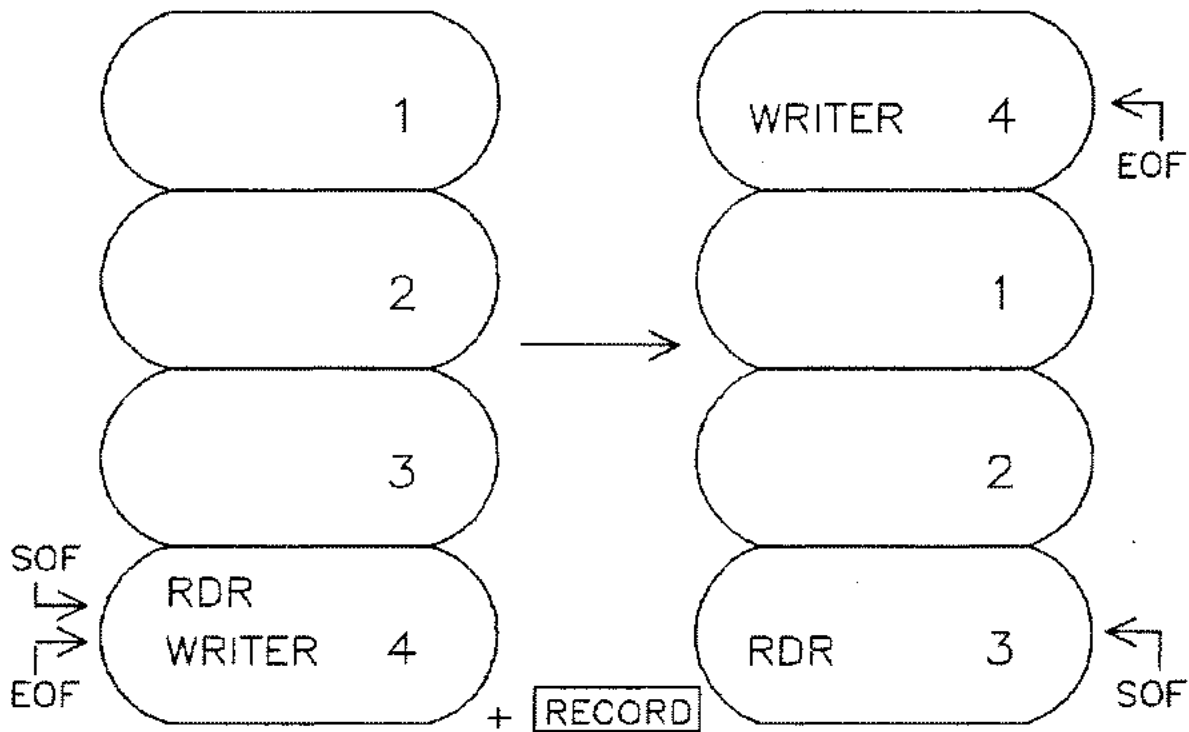
Both reader and writer buffers are contained in the same pool of buffers physically attached to the end of the MACB. Reader buffers are at the "head" of the queue whereas writer buffers comprise the "tail" of the queue. Buffer management runs in one of two modes.

---

<sup>9</sup>\*\*Note that variable length record files have a blocking factor of one record per block.

**COUPLED MODE**

When all active, unread records may be contained in the buffers buffer management is said to be in coupled mode.<sup>10</sup> A spent read buffer<sup>11</sup> is added to the tail of the write buffer list. Writers write to the tail buffer in the list. If this happens to be shared with a reader, then no disc I/O is initiated when the writer has filled the buffer. Otherwise when the buffer is full and only used by the writer on the next write, the buffer is flushed, copied to disc, and then the record is written. Note that if all the accessors were to close the file then all data would be written to the disc. When the number of active records can no longer be held in the buffers, uncoupled mode is entered.



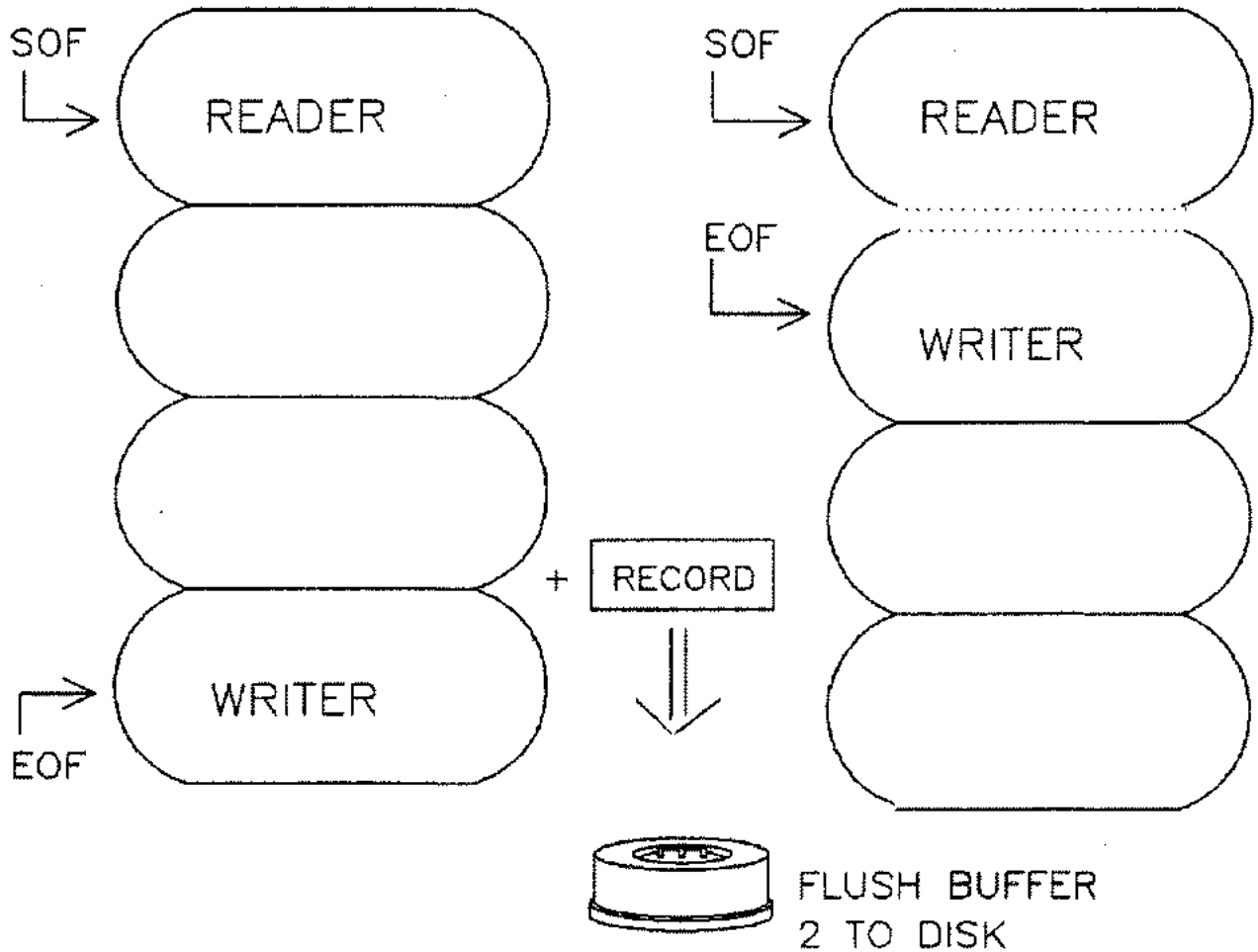
ALL ACTIVE RECORDS MAY BE  
CONTAINED IN THE BUFFERS

<sup>10</sup>Note that coupled mode is guaranteed with FOPEN.

<sup>11</sup>Remember that records are logically deleted after having been read. A spent read buffer is a block (physical record) in which all records have been deleted.

UNCOUPLED MODE

The active records cannot be contained in the buffer area. In uncoupled mode readers and writers use their own buffers independently. The writers will circularly use the buffers not being used by the reader and flush each buffer when it's necessary to add a record. Note that if the reader reads sufficient records such that the active records may be contained in the buffers, buffer management automatically reverts to coupled mode.



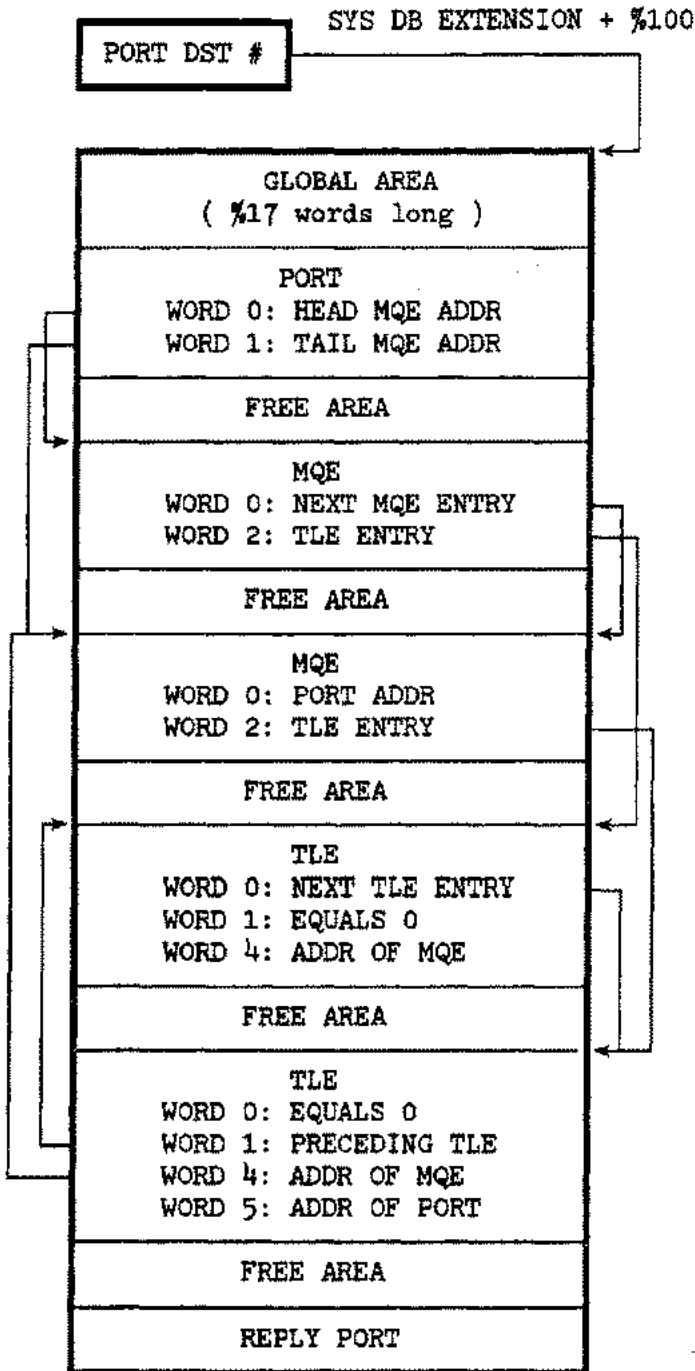








## Port Data Segment



MQE: Message Queue Entry  
 TLE: Timer List Entry

## PORT DATA SEGMENT GLOBAL AREA

The PDS global area is 17 octal words in length. It contains global information to manage the wait queues, reply ports and time list entries.

\* Processes may be impeded waiting for free blocks in the Port Data Segment. These processes will form a linked list with the head and tail PIN number at words %10 and %11 in the port data segment global area, and PIMP PIN (previous impeded process) and NIMP PIN fields (next impeded process) in the PCB. \* The TLE (the smallest time-out) address is in word %112. It heads a thread of TLE's through the Port Data Segment. The first TLE is also an entry in the system Timer Request List. That entry number is kept in word %13.

0	: Data segment number of this port data segment	: 0
1	: Block size in words	: 1
2	: Total number of blocks	: 2
3	: Maximum number of blocks	: 3
4	: Current number of free blocks	: 4
5	: Number of open ports	: 5
6	: Head of free list	: 6
7	: Tail of free list	: 7
10	: Head of impeded process list	: 8
11	: Tail of impeded process list	: 9
12	: Head of timeout thread (TQE address)	: 10
13	: TRLX of timeout	: 11
14	: Value returned by TIMER intrinsic when	: 12
15	: Timeout was initiated.	: 13
16	: Head of Port List	: 14
17	: Not Used	: 15

## Queuing Mechanisms

### PORTS

All entries in the Port Data Segment (ports and message queue entries and Timer List Entries) are ten octal words long. Ports contain overhead information to maintain queues. You will find a port for each wait queue that is open. And there is a reply or return for each message file opener.

#### NOTE

*There is one exception to this rule. In semi-exclusive mode there may be at most one reader, and in exclusive mode there may be only one writer. The generality of a wait queue is not needed. Instead when the process attempts to read from the empty file or write to a full file the reply port number is placed directly into the MACB. If the process has opened the file with wait, then no reply port is used.*

Information contained in a port is shown in the figure on the next page. Head and tail MQE pointers are used for access to queues. Soft interrupt subtype is always "1". The port manager (word 3 (8:8)) is the process that opens the port. Word 6, the number of sends, is an accumulated number of MQEs that have been attached to the port. And the label for the soft interrupt procedure is written when passed by an FCONTROL 48.

```
** LABEL      :I/E:  : SIT ENTRY NO  :CST/CSTX NO :
```

```
.....  
.....  
I/E: Internal or External Bit
```

When soft interrupts are armed with an FCONTROL 45, IPC procedure FCONTROL turns on the Enable Wake Up bit (word %2 (0:1)).<sup>11</sup> When the port is accessed, the Enabled Wake Type bit (word 2(1:1)) tells IPC procedures the action to take when a message is received.

---

<sup>11</sup>Note: This bit is also toggled by another IPC procedure - LONGWAIT).

PORT

```

    0  1  2  3  4  5  6  7  8  9 10 11 12 13 14 15
    .....
0 :  Head MQE address                               : 0
    .....
1 :  Tail MQE address                               : 1
    .....
2 :E : W : Next port number in port list thread: 2
    .....
3 :  Soft int subtype      ; Pin of port's owner  : 3
    .....
4 :  Soft interrupt parameter one                   : 4
    .....
5 :  Number of MQEs in the port's queue             : 5
    .....
6 :  Number of sends to this port                  : 6
    .....
7 :  Soft interrupt plabel                          : 7
    .....
:0 :1 :2 :3 :4 :5 :6 :7 :8 :9 :10:11:12:13:14:15:

```

W Type word 2(1:1)

- 0 - Awaken on port
- 1 - User soft interrupt
- 2 - System soft interrupt

## MESSAGE QUEUE ENTRIES

The Message Queue Entry records wait information when attached to the reader or writer wait queues: is there a timeout, the return port number, writers ID, DST number of the buffer, etc. When attached to a reply or return port an MQE stores the number of bytes transmitted or an error code for an incomplete read or write. See the parameter explanations on the following page for information contained in a wait message or completion message.

```

0  1  2  3  4  5  6  7  8  9  10 11 12 13 14 15 16
.....
0 : Next MQE entry; if last, (port addr) LOR 7 : 0
.....
1 : Port number of return port : 1
.....
2 : Time List Entry (TLE), 0=no timeout, -1=timed out: 2
.....
3 : Parameter zero : 3
.....
4 : Parameter one : 4
.....
5 : Parameter two : 5
.....
6 : Parameter three : 6
.....
7 : Parameter four : 7
.....
:0 :1 :2 :3 :4 :5 :6 :7 :8 :9 :10:11:12:13:14:15:

```

### Timer entry definitions

- 0 - No timeout
- 1 - Timeout expired
- 2 - TLE address for a pending timeout



Wait Message

parm#

- 0 - WRITER ID
- 1 - LOCAL FLAGS (differ with each accessor)
  - (0:1) - accessor just opened file
  - (1:1) - will wait on boundary condition if no symbiotic process
  - (3:1) - writer has not written a record
  - (4:1) - transmission log in bytes
  - (8:1) - carriage control code
- 2 - DST# of data buffer
- 3 - Address of data buffer (DST relative)
- 4 - Length of data buffer in bytes

Completion Message

- 0 - Resultant error code
- 1 - Resultant transmission log in bytes

**TIMER LIST ENTRY**

Below is a Time List Entry. Its use and operation has been explained on page 4-3. The captions inside the picture should be sufficient to explain the TLE fields.

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15																
	.....																															
0	:	Next TLE (sorted in incr time val), 0 if last:														0																
	.....																															
1	:	Preceding TLE entry (0 if first entry)														1																
	.....																															
2	:	Number of milliseconds the timeout value														2																
	.....																															
3	:	of this TLE is beyond the previous TLE.														3																
	.....																															
4	:	Address of the affected MQE														4																
	.....																															
5	:	Address of the MQE's port														5																
	.....																															
6	:	Value of TIMER* when this timeout expires														5																
	.....																															
7	:	{Milliseconds}														7																
	.....																															
	:	0	:	1	:	2	:	3	:	4	:	5	:	6	:	7	:	8	:	9	:	10	:	11	:	12	:	13	:	14	:	15

\* TIMER is an internal MPE intrinsic that returns the time in milliseconds.

### HOW DO WE FIND THE PORT DATA SEGMENT?

The DST number of the port data segment is found at Sysglob Extended +%100. The address of sysglob extension is located at %1377 (sysglob +%377). This is a sysglob relative address. Add it to %1000 and you will be at the beginning of the extension area. One-hundred and one-hundred-and-one octal words into the extension area is called the **DST Number Array**. The DST Number Array is simply two words, the first of which contains the DST number of the existing port data segment and the following word is reserved for a possible second port data segment.

#### PORT DST NUMBER ARRAY

Located in System DB Extension Area.

.....		
64	: Port data segment number	: 64
.....		
65	: Reserved for a second port segment	: 65
.....		

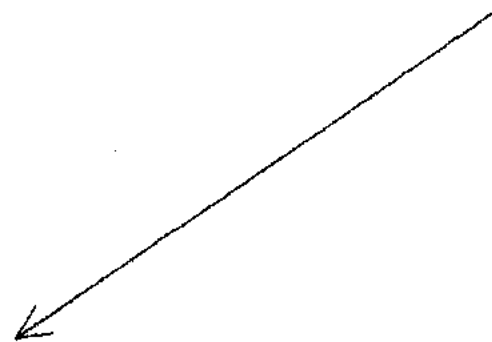
# MEMORY DUMP

POINTER TO  
SYSGLOB  
EXTENSION  
(SYSGLOB  
RELATIVE)

001370: 000000 000400 074163 000113 100561 000001 000504

001500: 030370 030370 030370 000000 1504

001600: 020103 030060 031060 000000 000130 000000



PORT DST NUMBER ARRAY  
IS AT SYSGLOB EXT  
+%100 AND %101

PORT DATA SEGMENT  
DST#

AVAILABLE FOR  
SECOND PORT  
DATA SEGMENT

## Port Data Segment Operation

This section explains Basic IPC operation and the fundamental operation of the port data segment. Soft interrupts and its message facility is explained later in this section. Below is a description of the port data segment operation a program will FOPEN a file, read from an empty message file or writing to a full message file.

### FOPEN

When the message file is opened a reply port is initialized for the opener. If it is the first time the file is opened both reader and writer wait queues are created by establishing a port for each queue and write pointers to those ports in words %110 and %111 of the message ACB.

Pointers to ports have the format shown below. The port index refers to the word to use in the DST number array. Since we have only one port data segment at this time the index is always "0". The remaining part of the word indexes into the port data segment by port number.

#### PORT NUMBER

```

0  1  2  3  4  5  6  7  8  9 10 11 12 13 14 15 16
.....
:Port index :  Port data segment relative addr/8  :
.....

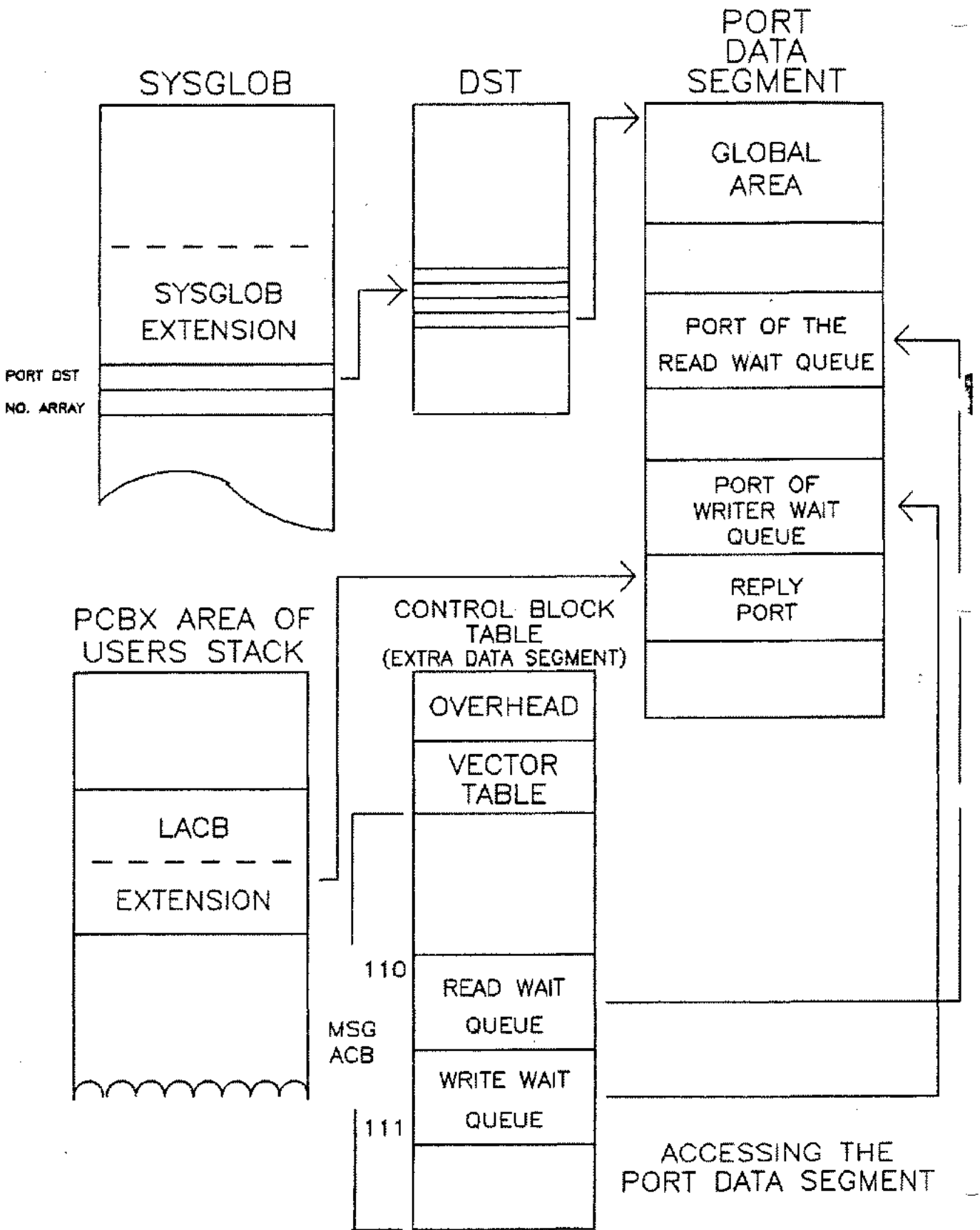
```

The address of the openers reply port is in the third word of the openers LACB extension. There is one small exception, however, if the opener is using waited I/O, a reply port is never built.<sup>13</sup>

---

<sup>13</sup>Most customers will use nowait I/O in IPC applications because of its throughput advantage.

Queueing Mechanisms



## WRITING

The usual course of a write is to 1) insure that there is room for the record, 2) write the record (usually just a data move to an ACB buffer), and, 3) if a reader is waiting send a successful message to the readers reply port.

The following exceptions might occur:

### A. The first write.

The writers first write after an FOPEN consists of writing the open record, the data record and allocating space for the closed record.

### B. Impeded writer.

When the file is first opened the amount of free file space (in max-size records) is calculated. Each write to the file first subtracts its record size and record overhead from the free space. If the record will not fit into the file, then the writer is \*waited if:

- 1) a reader has opened the file, or
- 2) this is the writers first write after the open, or
- 3) the writer has specified extended wait (FCONTROL 45).

\* Waiting consists of placing a message on the wait queue. Eventually a reader returns free space (a block's worth at a time). The actual write is performed by the reader at this time. This is to:

- 1) Expedite the no-wait writer's data, and
- 2) insure the released writers will fill the file in the same order as they were freed (any other sequence may result in the writers running out of file space).

When the transfer is completed a successful message is attached to the writer's reply port.





**READING**

Most read requests read a record and if the next record is in another block, issue an anticipating read of the block from the disc.

The following exceptions may occur:

**A) Empty File**

The reader decrements a counting semaphore (number of records in the file) when the semaphore goes to "0" there are no more records left. The reader is waited if:

- 1) One or more writers has opened the file
- 2) Or this is the readers first read to the file after the open
- 3) Or reader has specified extended wait (FCONTROL 45). [NOTE: Do not confuse this with extended read (FCONTROL 46)]

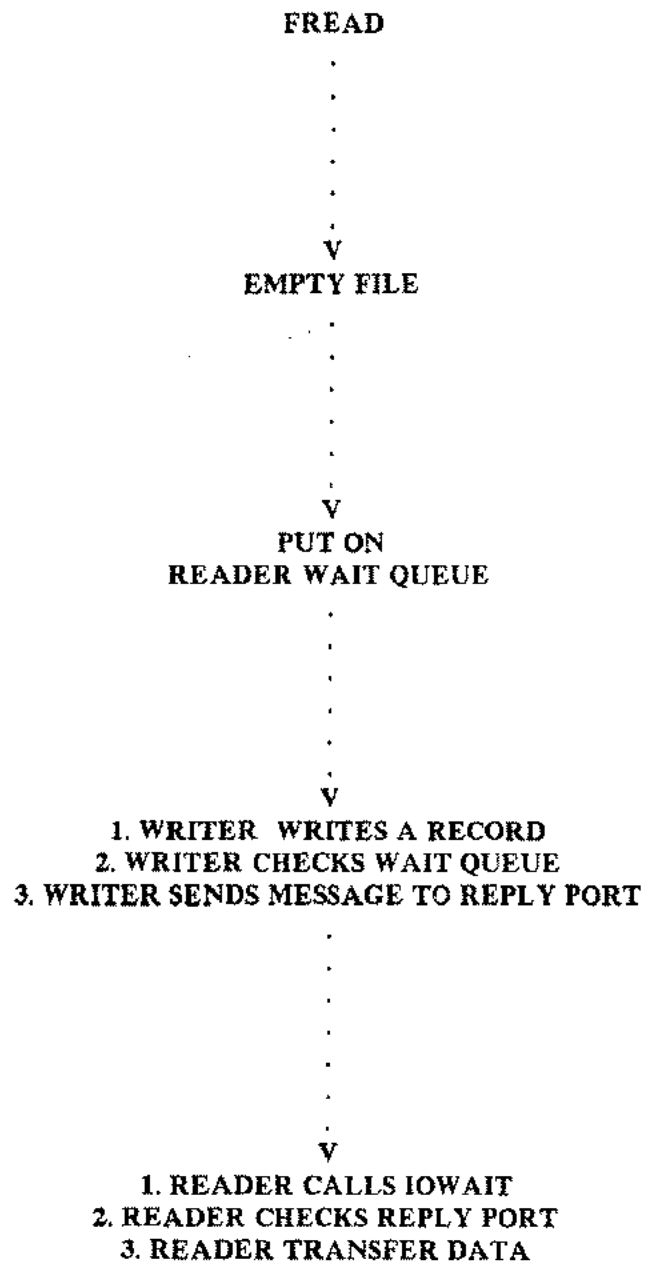
Again, parallel with writing to a message file \*waiting consists of placing a message on the reader's wait queue. Each time a writer writes a record the writer checks the reader wait queue. When the writer discovers an impeded reader it will send a successful message to the reader's reply port. The reader then (or at I/O wait time) performs the data transfers.

Waited readers, unlike waited writers, perform their own data movement upon being liberated. They can do this because they are freed with a claim on any record in the file -- not a particular one. Thus it's permissible for other readers to issues FREAD/IOWAITs between another reader's FREAD and IOWAIT. The only constraint is that a one record claim be set aside for the first reader when the process is liberated from the wait queue.

**B) Extended Read Mode (FCONTROL 46)**

Possible open/close records must be logically deleted before the actual read to the users target area. If the reader must wait due to an empty file, then when the reader is awakened by the writer this process must be repeated until the actual read can commence. To detect an empty file the reader decrements a counting semaphore. When the semaphore goes to "0" there are no more records left, hence an empty file. If not in extended read then only data records count in the counting semaphore. Extended read causes all records (data, open, and close) to be counted.

Queuing Mechanisms



## Soft Interrupts

Soft interrupts not only uses the port data segment for queuing I/O requests but additionally shares a memory resident message facility with system processes to implement interrupts and interrupt handling. IPC uses this same facility to queue interrupts. Since this message facility is memory-resident, it is fast and can be accessed while on the ICS. There are three data structures that make up the in-core message facility, the Message Harbor Table, Primary Message Tables, and the Secondary Message Table.

### MESSAGE HARBOR TABLE

The Message Harbor Table is DST %71. Each process running on the system including users and system process has an entry of 5 words in the Message Harbor Table. The entry is known as the process's "Message Harbor." The first four words are links or addresses into either the Primary Message Table or Secondary Message Table respectively where system and user process messages are queued. IPC uses word 1 in the process's Message Harbor and it heads a linked list of interrupts in the Secondary Message Table. The fifth word in a process's message harbor is a port mask of the ports being used, hence bit 1 is set when an interrupt is queued.

```

.....
: LINK TO FIRST MSG PORT 0 :
.....
: LINK TO FIRST MSG PORT 1 :
.....
: LINK TO FIRST MSG PORT 2 :
.....
: LINK TO FIRST MSG PORT 3 :
.....
: NON-EMPTY PORT MASK :
.....

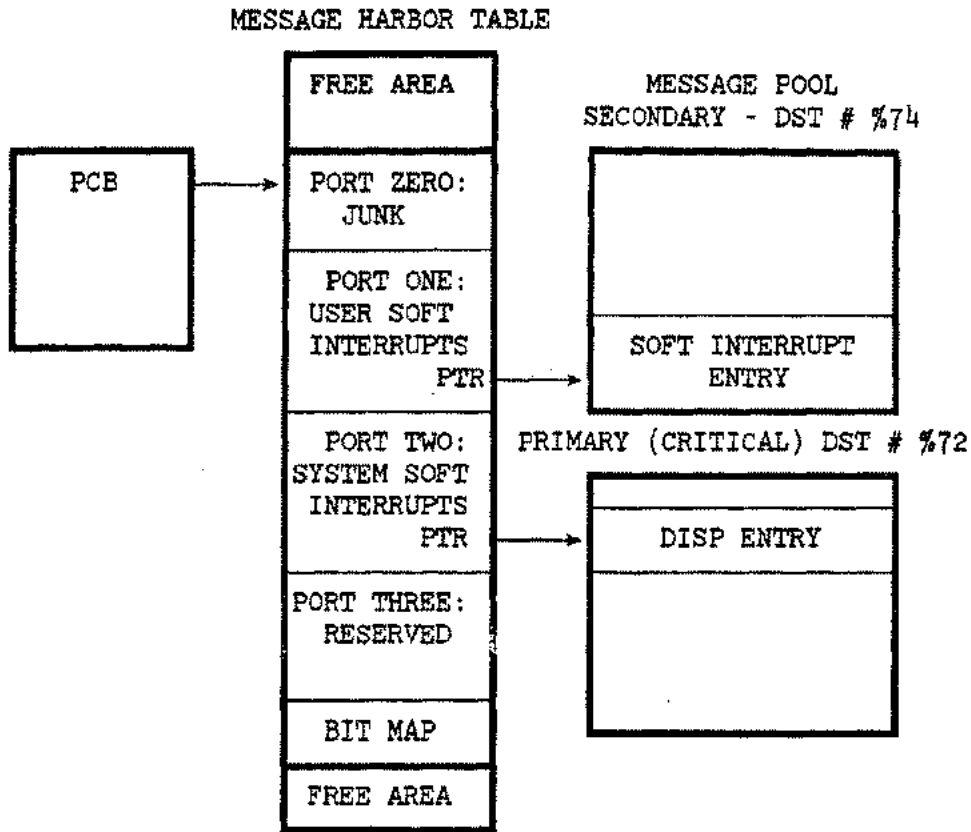
```

First Message Queue Link (0:1) = 1 ==> Next Message in Secondary Message Table  
 (1:15) = Index of next message in appropriate table

### PRIMARY AND SECONDARY MESSAGE TABLES

The primary message table is DST %72 and is used to queue system messages. The secondary message table is at DST %74 and is used to queue software interrupts in an FIFO manner. The messages are posted by the process that relieves the boundary condition and PCB wait flags are set. (i.e. The writer writing to an empty message file or the reader that frees space). Later when the dispatcher is going to launch the waited process it will note the PCB wait flags. The PLABEL in the secondary message table is used to launch this process into its interrupt handler.

WORD 0: NEXT MESSAGE IN QUEUE LINK
WORD 1: SOFT INTERRUPT MESSAGE WORD ZERO
WORD 2: SOFT INTERRUPT MESSAGE WORD ONE
WORD 3: INTERRUPT HANDLER'S PLABEL
WORD 4: SOFT INTERRUPT SUBTYPE

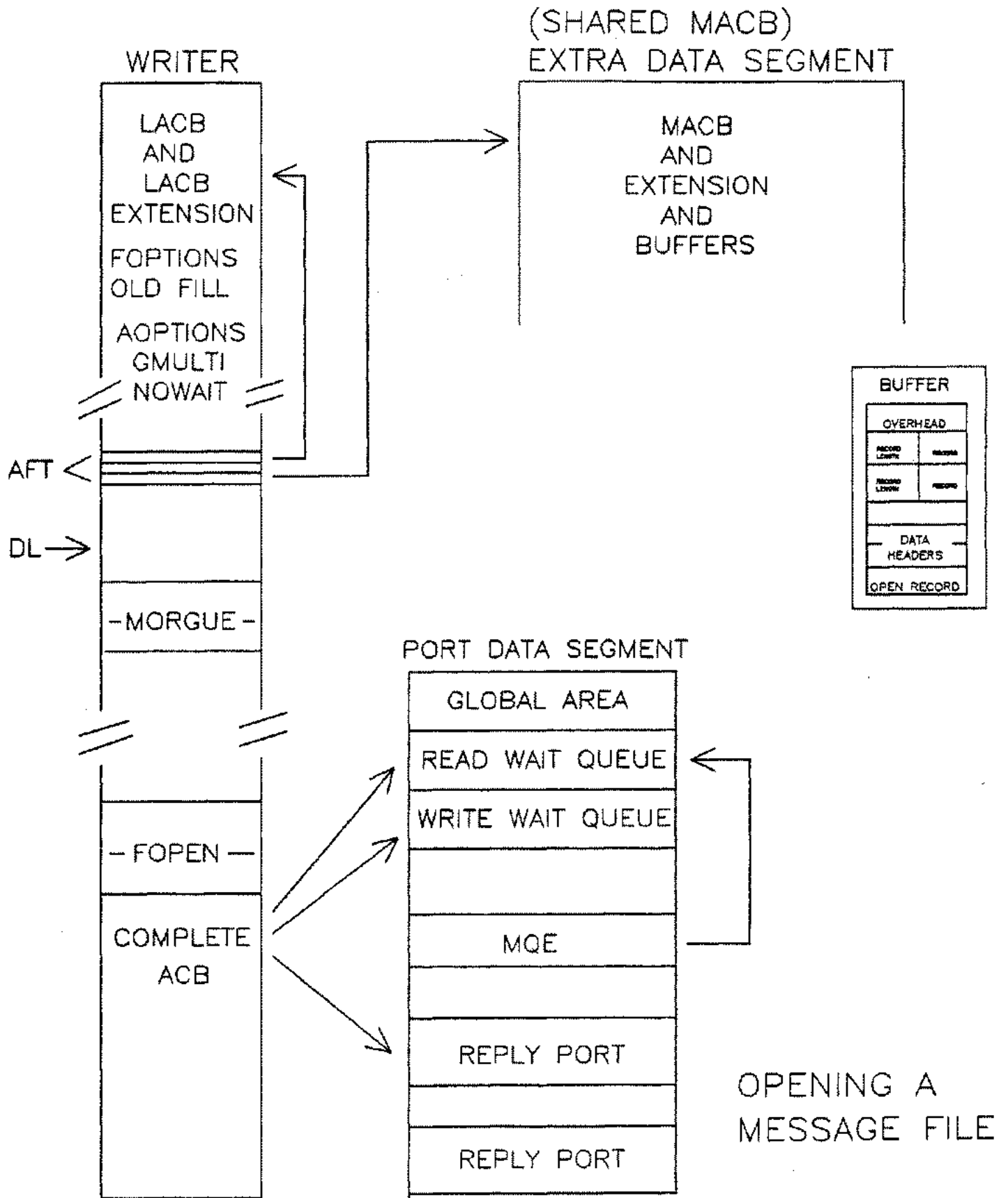


## Soft Interrupt Operation

In this example of a soft interrupt operation a reader has already opened and written to a message file. The reader has enabled soft interrupts with an `FINSTATE` and allows extended wait with an `FCONTROL 45`. The reader is still executing but has queued a request on the reader wait queue. We are going to follow the writer until the read completes.

### OPENING THE MESSAGE FILE

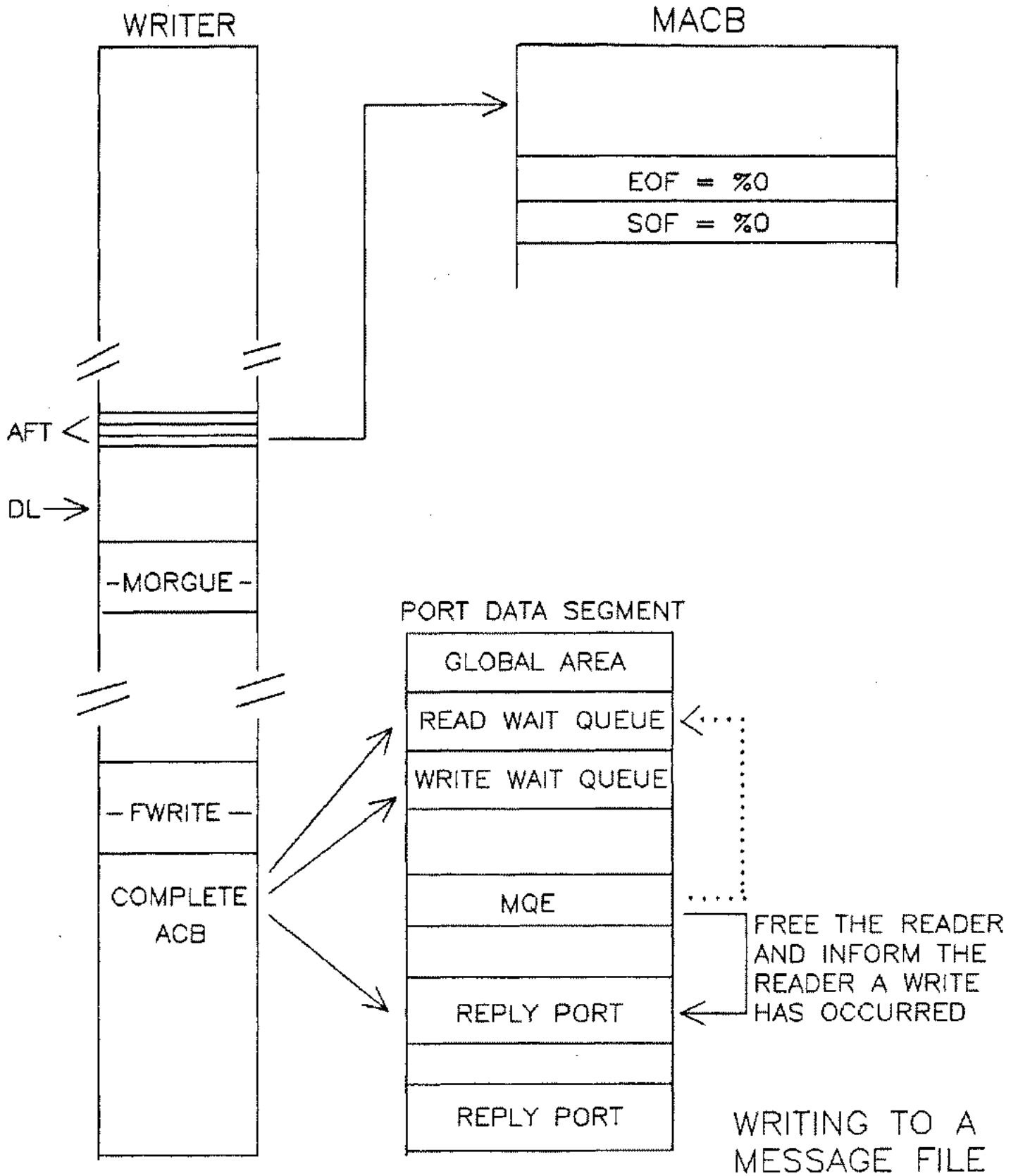
1. Open an old message file with write, `Gmulti` access and no-wait attributes.
2. The writer is assigned a 16 bit `Writer's ID` and a return port is allocated in the port data segment. This process's reply port is located via word 3 of its `AFT` entry and the reply port points back to the process in the right byte of word 3 - "PIN of the Owner". A writer and reader wait queue have already been established by the first opener of the message file (a reader in this case). Pointers to the reader and writer wait queues are found in the shared `Message ACB`.
3. A `Logical Access Control Block` and `Logical Access Control Block Extension` are created in the writer's own stack. These control blocks will contain local variables. A buffered message access control block and its extension and a `File Control Block` already exist in extra data segments. Again the process's `Available File Table` contains pointers to locate the `MACB` and `LACB`.



**AN FWRITE IS INVOKED**

1. After an Fwrite is issued, procedure LOC'ACB locates an LACB and MACB via an AFT entry and locks these control blocks with the locking mechanism in the vector table entry for each of these control blocks. The LACB and its extension are laid on the stack.
2. LOC'ACB also places an image of these control blocks on the user's stack.
3. IPC procedure FCWRITE is invoked and the MACB and MACB extension and local parameters are placed on the stack. The complete ACB now lies on the stack.
4. FCWRITE performs the write to the target buffer. Keep in mind data and its header information is written. When the buffer becomes full the actual I/O takes place.
5. Still within the writer process, the reader wait queue is checked for waiting readers (subroutine FREEREADER).





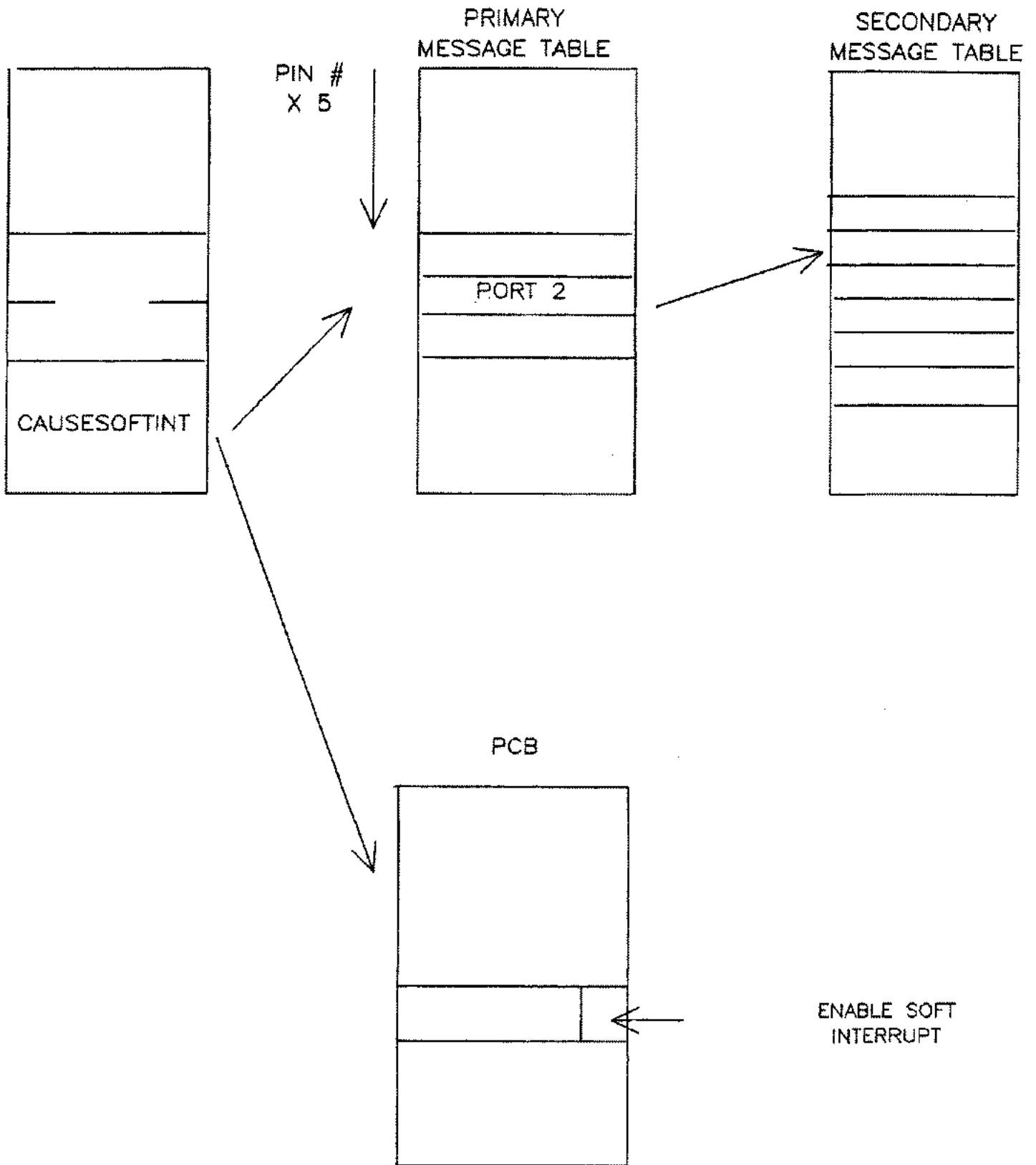
## THE SOFT INTERRUPT OCCURS

The reader wait queue is examined for waiting readers. An MQE is found linked to port of the reader wait queue. This reader process has enabled soft interrupts with an FINSTATE intrinsic. The "enable wake up" bit in the reply port is set (word 2 (0:1)) and the "allow soft interrupts" bit in the reader's PCB is turned on (PCB 13 (7:1)). The writer will make the necessary arrangements for the reader to be soft interrupted.

The writer process finds the address of the reply port from the MQE that was taken from the reader wait queue. The reply port gives us the information that identifies who the reader is (PIN number in word 3 (8:8)) and also contains the PLABEL (word 7) of the reader's soft interrupt routine. With this information the writer calls procedure CAUSESOFINT of KERNELC and passes the data as parameters to the procedure. The plabel must specify code which is within the interrupted process's code address space.

CAUSESOFINT first checks to see that the target process is healthy (not in soft or hard kill mode). It then sends a message to the process's appropriate soft interrupt port. That port in the secondary message table DST %21 is found by indexing into the process's Message Harbor within DST %71. How far do we index - pin number x 5 (each entry is 5 words long + base of DST %71. The second word in the process's message harbor points to the beginning of the process's queue in the secondary message table where CAUSESOFINT will queue its interrupt request.

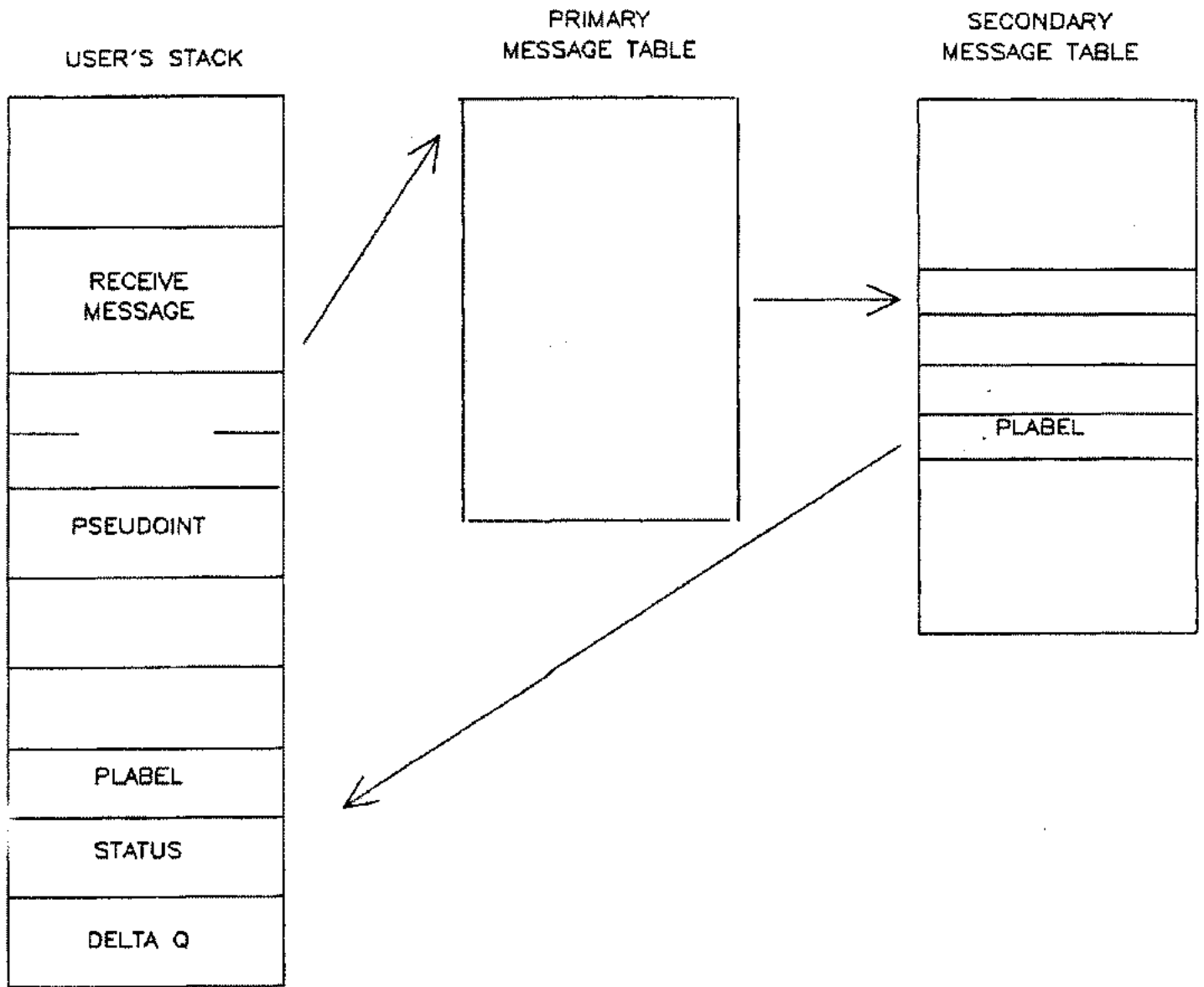
Since the reader we are dealing with is not critical, with a SIR, or impeded, or waiting, CAUSESOFINT will set the Soft Interrupt bit PCB 9 (9:1). Since the process is not waiting, it must already be on the dispatch queue so that no further action is needed until the dispatcher and procedure PSEUDOINT are called to service the interrupt.



## Queueing Mechanisms

The dispatcher scans the PCB of this process before it tries to launch the process. Since the soft interrupt bit is set; the dispatcher writes a stack marker on the process's stack to enter procedure PSEUDOINT. Then the process is launched. PSEUDOINT checks that the user has soft interrupts enabled (FINSTATE). If not, PSEUDOINT returns to the interrupted code and when the user reenables soft interrupts with either FINSTATE or FINTEXT, the process checks for any pending soft interrupts.

Next it is determined if the interrupt can be executed right away. That is, the process cannot be executing in system code or the user interrupt procedure must be privileged if we are interrupting privileged user code. If these two conditions are satisfied, then PSEUDOINT retrieves the message from the secondary message table, disables user soft interrupts, builds a stack marker for the interrupt handler and exits through it.



## DELAY OF SOFT INTERRUPTS

If the code was not compatible the delayed soft interrupt bit in the process's PCB is set again and bit 0 of the P-register save word is set to 1 in the first eligible stack marker. This will cause a bounds violation trap to ININ, which will note a soft interrupt is pending and will invoke PSEUDOINT thereby delaying the soft interrupt.

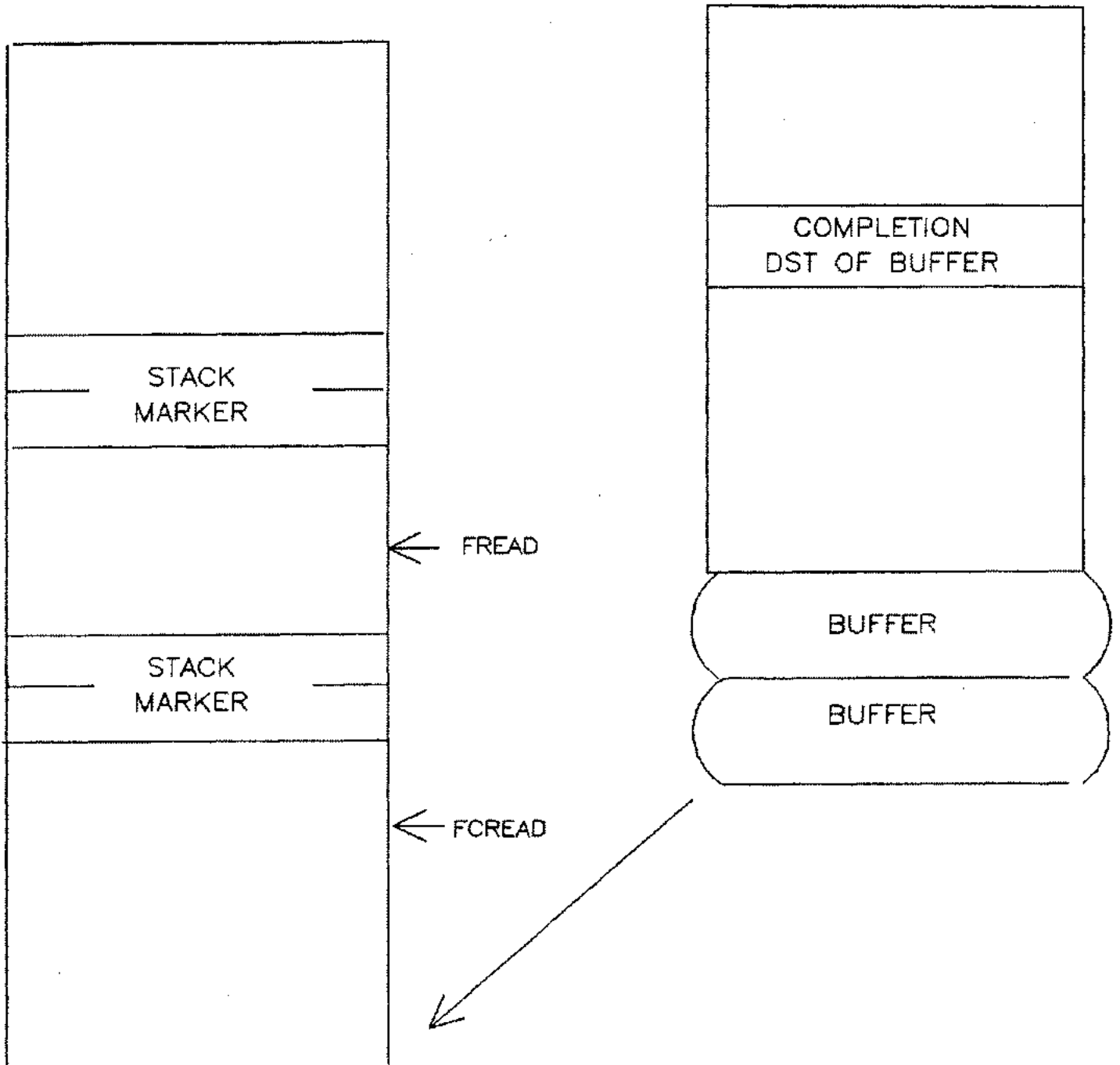
If the process is waiting but is not critical. It may have a SIR or be impeded on some other event on which the "delay soft interrupt bit" (PCB 0 (8:1)) is set. When the process is ready to launch the dispatcher will notice this bit and launch the process into PSEUDOINT so the interrupt will be served.

If the process has some combination of being critical, is with a sir or is impeded at the time the soft interrupt occurs, then the delayed soft interrupt bit and the psuedo-interrupt overflow bits (PCB 0 (4:1)) are set in the PCB. The dispatcher will not have to be recalled because after inhibiting conditions are removed, RELSIR and RESETCRITICAL (called by the awakening process) both look at the PSEUDOINT overflow flag in the PCB and call PSEUDOINT if the flag is set.

Upon entering the interrupt handler routine, soft interrupts are turned off. The soft interrupt routine calls IODONTWAIT to complete the I/O, does whatever soft interrupt processing it wishes to do, and then FINEXIT will reinstate soft interrupts and normal processing resumes.

AWAKE READER

MACB



## TIMEOUTS

A timed I/O request (initiated by issuing an FCONTROL4) upon timing out will return a condition code of CCL (error 22 is also set which can be retrieved by FCHECK). A timeout request will have created an MQE and a Timer List Entry in the port data segment. These two entries will have pointers to one another. All timer list entries in the port data segment will be linked to one another in ascending order (word 0) and the head TLE (smallest remaining interval) has an entry in the system's Timer Request List (DST %23). When the interval passes, Tick (internal interrupt procedure) updates the entries in the TRL, hence a timeout will be recognized on the ICS. Since the port data segment is not memory resident and there is no mechanism for waiting while on the ICS, timeouts of I/O requests for message files are handled by procedure OLD'TICK. TICK will set bit 12 in sysglob cell %121. This will inform UCOP that a timeout has occurred on a port data segment. OLD'TICK awakes UCOP. UCOP may wait (if necessary) for the port data segment to be brought into memory and does the necessary processing for the timeout. UCOP will send a message to the reply port and awaken the process from its timer wait.

## CLOSING A MESSAGE FILE

1. All outstanding requests are cancelled.
2. For readers, any outstanding claim to a record is returned.
3. If this is the last writer, then any reader waiting process (not in extended wait mode) is satisfied with a CCG.
4. A close record is written.
5. The reply port and record buffer are deleted.
6. If this is the last writer, then spent extents are deleted.



## Procedural Flow of a Soft Interrupt

### INITIATING A SOFT INTERRUPT READ

---

```

USER CODE  [ FREAD(FILENUM,TARGET,TCOUNT);
            << FILENUM = file number returned from fopen >>
            << TARGET = buffer which will receive data >>
            << TCOUNT = number of bytes/words to read >>

FILEIO     [ INTEGER PROCEDURE FREAD(FILENUM,TARGET,TCOUNT);
            ERRORON; << CHECKER - MODULE 69 >>
            << TURN TRAPS OFF >>

            SETCRITICAL; << KERNELC - MODULE 92 >>
            << SET CRITICAL IN PCB >>

            LOC'ACB(*,11,FILENUM,UMODE);
            << IF INVALID FILE NUMBER - RETURN CCL >>

            CASE FTYPE << FTYPE 8 = A MESSAGE FILE >>

            IF IOQX <> 0 THEN RETURN CCL
            << NOWAIT I/O PENDING ALREADY - AFT WORD 3 >>

            FCREAD(ACB'NOWAIT,TARGET,TCOUNT);

            << ACB'NOWAIT = AOPTIONS.(4:1) NOWAIT I/O
            MODE >>

IPC        [ FCREAD

            [ FCREAD(FUNCTION,ADDR,TCOUNT);
            INITIALIZE; << COMBINE LACB AND PACB AT Q+ AREA >>

            PUTWAITQUEUE(ACB,0,0,LENGTH);
            << THIS PROCEDURE PLACES A WAIT MESSAGE INTO THE
            APPROPRIATE READERS WAIT QUEUE. THIS IS
  
```

IPC

EITHER THE ACB RESIDENT QUEUE OR THE PORT READERS WAIT QUEUE.

ALGORITHM FOR DETERMINING THE PROPER QUEUE:

1. THE READER WAIT QUEUE MUST BE USED IF:
  - A. MORE THAN ONE ACCESSOR CAN QUEUE UP.
  - B. THE TIMEOUT VALUE IS SPECIFIED (ONLY AVAILABLE ON THE PORT WAIT QUEUE).

NOTE: IF THE PORT WAIT QUEUE IS USED, THEN THE REPLY PORT MUST ALSO BE USED. THE REPLY PORT MUST BE USED FOR NOWAIT I/O SO THAT IOWAIT/IODONTWAIT CAN EFFICIENTLY CHECK THE REQUEST'S STATUS.

```

FCPREPAFT(FILENUM,SOFT'INT'PEND = -2);
  << CHANGE THE MESSAGE FILE AFT ENTRY - WORD
    THREE (PORT NUMBER WORD FOR NOWAIT I/O)>>

READEXIT(ERROR);
PUTMYCOMPLMSG(ACB,ERRORCODE,0);
  << SEND A MESSAGE TO OWN LACB >>
CAUSESOFINT(0,USER'INT,FILE'SOFT'INT,
             SOFTINTPLABEL,2,0);

  << 0 = PIN NUMBER OF TARGET PROCESS - A
    ZERO IMPLIES OWN PROCESS>>
  << USER'INT:=0 = TYPE OF SOFT INTERRUPT
    0 - INTERRUPTS PROCESSED ONLY WHEN USER
      INTERRUPTS ARE ENABLED AND THE PROCESS
      IS EXECUTING IN USER CODE (SEG %3XX).
    1 - INTERRUPTS PROCESSED IN CURRENT STATE
      UNLESS THE TARGET PROCESS IS WAITING, HAS
      SIR, IMPEDED, OR CRITICAL.
    NOTE: THE ABOVE CONDITIONS ONLY DELAY THE
    SOFT INTERRUPT. ALL TYPE ONES ARE
    SERVICED BEFORE TYPE ZEROES.
  << FILE'SOFT'INT:= 1 = SUBTYPE, TYPE OF
    PRE-PROCESSING TO BE DONE BY PSEUDOINT
    PROCEDURE. A ZERO IMPLIES NO
    PREPROCESSING .
  << SOFTINTPLABEL = THE INTERRUPT HANDLER
    PLABEL>>
  << 2 = NUMBER OF WORDS IN THE MESSAGE >>
  << 0 = LOGICAL FLAGS;
    (0:14) RESERVED FOR FUTURE USE
    (0:1) 0 = IF APPROPRIATE PLACE PROCESS
      ON THE READY LIST
    1 = DO NOT PUT ON READY LIST
  
```

KERNELC [ CAUSESOFINT(PIN,TYPE,SUBTYPE,PLABEL,MSGLEN,FLAGS)

KERNELC [ END; << RETURN TO IPC >>

IPC [ UNLOCEXTENDACB(ACB);  
 << ROLLS BACK THE MESSAGE FILE EXTENSIONS TO THE  
 PACB AND THE LACB. >>  
 ASMB (EXIT 1).  
 END; << RETURN TO FILEIO >>

FILEIO [ RESETCRITICAL;  
 END. << RETURN TO USER CODE >>

USER CODE [ SOFT INTERRUPT ENABLED - USER CODE CAN CONTINUE  
 TO PROCESS IN MAINLINE CODE.

COMPLETING THE SOFT INTERRUPT READ

---

```

USER CODE  [ FWRITE(FILENUM,TARGET,TCOUNT,CONTROL);
            << FILENUM = FILENUMBER RETURNED FROM FOPEN >>
            << TARGET = BUFFER WHICH CONTAINS THE DATA >>
            << TCOUNT = NUMBER OF BYTES/WORDS TO WRITE >>
            << CONTROL = CARRIAGE CONTROL CODE >>

FILEIO     [ PROCEDURE FWRITE(FILENUM,TARGET,TCOUNT,CONTROL);

            ERRORON; << CHECKER - MODULE 69 >>
            TURN TRAPS OFF;

            SETCRITICAL; << KERNELC - MODULE 92 >>
            SET CRITICAL IN PCB;

            LOC'ACB(*,9,FILENUM,UMODE);
            << IF INVALID FILE NUMBER - RETURN CCL >>

            CASE FTYPE << FTYPE 8 = A MESSGE FILE >>

            IF IOQX <> 0 BEGIN << NO-WAIT I/O PENDING
                                RETURN CCL >>

            FCWRITE(ACB'NOWAIT,TARGET,TCOUNT);
            << ACB'NOWAIT = AOPTIONS.(4:1) NOWAIT I/O >>

```

```

PROCEDURE FCWRITE(FUNCTION,ADDR,TCOUNT);
  <<FUNCTION = 1 COMPLETE WRITE >>

  LOCEXTENDACB(ACB'LOC,AATARGETDST,ADDRESS);
    << BUILD CACB FROM LACB/PACB >>

  PUTRECORD(ACB,RECTYPE,CALLERDST,TARGET,LENGTH);
    << PUTS WRIT'ERS DATA INTO A WRITE BUFFER >>

  FREEREADER(ACB);
    << INFORMS A WAITING READER A WRITE HAS
      OCCURRED >>

  GETWAITQUEUE(ACB,ABREADQUEUE,MSG,DELETE'MSG);
    << GET WAIT QUEUE MESSAGE FROM ACB OR PORT
      WAIT QUEUE >>

  PUTCOMPLMSG(ACB,PORT,ERROR,TLOG,ID);
    << PLACES A COMPLETION MESSAGE INTO THE
      APPROPRIATE AREA ACB COMPLETION AREA OR
      THE REPLY PORT >>

  WRITEXIT;
    << RETURNS FROM THE ACCESS PROCEDURE >>

  UNLOEXTNEDACB(ACB);
    << ROLL BACK THE LACB/PACB >>
  ASMB(EXIT 1);
  
```

FILEIO

```

UNLOC'ACB(9,0)
  << RELEASE ACB - COPY ACB TO LACB/PACB >>

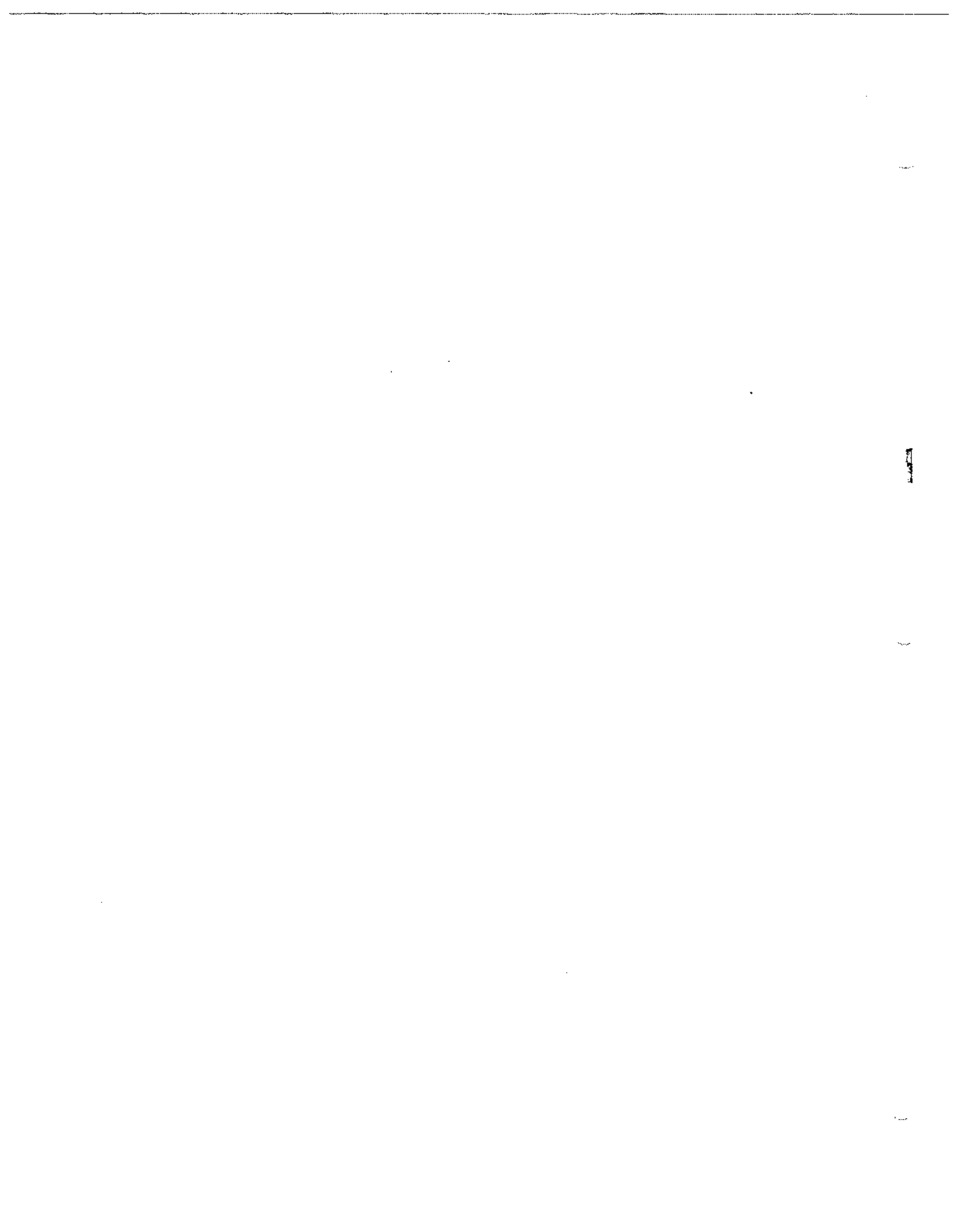
  IF SO <> 0 THEN AWAKE(*,PORTWAKE,0); ELSE DEL

  TOS:=ACB'ERROR
  
```

USER CODE

```

WRITE COMPLETE
  
```



# PCB FIELDS USED BY SOFTWARE INTERRUPTS

APPENDIX

A

Field Name	Location	Description
SPCBCritSir	PCB.(2:2)	Nonzero if the process is critical or with SIR.
SPCBIOvrFlag	PCB.(4:1)	Set when the execution of the soft interrupt must be postponed due to the process's s's being critical, with sir or impeded. Tells the appropriate kernel procedure to enter PSEUDOINT.
SPCBDelaySoft	PCB.(8:1)	Set when the execution of the soft interrupt must be delayed - critical impeded, with sir, system code, or privileged code. This bit tells the bounds violation procedure in ININ that a soft interrupt is pending. The dispatcher ignores this bit.
SPCBWaitField	PCB(4).(0:12)	Indicates which events (if any) the user process is waiting on. Does not include system events such as memory wait.
SPCBImpede	PCB(4).(12:1)	Set when the process is impeded.
SPCBPSim	PCB(8).(0:3)	Current value of process's pseudo interrupt state.
SPCBWakeSoft	PCB(8).(3:1)	Set on when the process will accept the soft interrupt even though it is waiting on other events.
SPCBSoftInt	PCB(9).(9:1)	Set when the dispatcher (and PSEUDOINT) should be aware of a pending soft interrupt.
SPCBPIFlags	PCB(9).(10:6)	Specifies which of the six, independent psuedo interrupts may be pending against the process.

PCB Fields

SPCBAllowSoft

PCB(13).(7:1)

On implies that the process will process user soft interrupts. A zero value postpones processing of user soft interrupts (but not system soft interrupts). This bit is managed by the user through FINSTATE and FINTEXT.



# MEMORY DUMP EXAMPLE

APPENDIX

B

## Summary

PCB 55 has opened the file MSG1.  
An FCONTROL 48 is enabled.  
Pin 55 has initiated a soft interrupt read.  
PCB 55 has not enabled soft interrupts (FINSTATE) at the program level.

PCB 54 has opened the file and written a record.  
It has triggered a soft interrupt on message file MSG1.  
A soft interrupt is now pending on PCB 55.  
Port 110 has a message pending soft interrupt.  
The Message Harbor contains the pending soft interrupt read  
that will not be processed because PCB 55 does not have interrupts  
enabled at the program level (FINSTATE).

## Message File: MSG1

### PCB 55:

OPENS FILE READ, EXCLUSIVE	(12:4,8:2)
FCONTROL 48 ENABLED	(LACB EXTENSION WORD 0)
SOFT INTERRUPT PENDING READ	(IOQX = -2 OR %177776)
FILE NUMBER = 11	(LACB)
REPLY PORT =110	(LACB IPC EXTENSION WORD 3)
AWAKEN THE PROCESS	(0:1)
GENERATE USER SOFT INTERRUPT	(1:2)
NO MQE	
PACB DST NUMBER = 124	(LACB)
WRITE WAIT QUEUE =100	(PACB)
READ WAIT QUEUE =70	(PACB)
PCB INT ENABLED @ PROG LEVEL	(PCB (13).(7:1)
ACCESSORS LOCAL FLAGS = 000000	(LACB IPC EXTENSION WORD 1)
FILE'S LOCAL FLAGS =004217	(PACB)

### PCB 54

OPEN FILE: WRITE, EXCLUSIVE	(12:4,8:2)
REPLY PORT = 120	(LACB IPC EXTENSION WORD 3)
WRITER ID = 1	(LACB IPC EXTENSION WORD 4)
ACCESSOR'S LOCAL FLAGS = 004001	(LACB IPC EXTENSION WORD 2)

Dump Example

FILE'S GLOBAL FLAGS =004217	(PACB)
FILE NUMBER = 10	(LACB)
PACB DST NUMBER = 124	(LACB)
WRITE WAIT QUEUE = 100	(PACB)
READ WAIT QUEUE =70	(PACB)



HP3000 III MEMORY DUMPC.00.02 OF SYS VER 0 UPDATE 01 FIX 08 DUMP TIME 10/16/11, 4:01PM  
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\*\*\*\*\* DST TABLE \*\*\*\*\*

SEGMENT NUMBER	SEGMENT DESCRIPTION	REFERENCE BIT	SEGMENT LENGTH	ABSOLUTE ADDRESS	BANK/ /LDEV	DISC ADDRESS	D V	R C	I Y	S K	M D	F ILE	1 6 T S	1 6 B R E A K	VH ALLOC
1	CODE SEGMENT TABLE	OFF	1400	007754	0										00
2	DATA SEGMENT TABLE	OFF	1440	008314	0										00
3	PROCESS CONTROL BLOCK	OFF	1400	013014	0										00
4	CST EXTENSION	OFF	1440	013354	0										00
5	SYSTEM GLOBAL AREA	OFF	700	001000	0										00
6	FIXED LOW CORE	ON	2000	000000	0										00
7	INTERRUPT CONTROL STACK	OFF	1400	014414	0										00
10	SYSTEM BUFFERS	OFF	2020	022704	0										00
11	UCOP REQUEST QUEUE	ON	104	134423	0										01
12	PROCESS-PROCESS COMMUNICATION TABLE	OFF	140	017223	1										00
13	I/O QUEUE	OFF	1030	015514	0										00
14	TERMINAL BUFFERS	OFF	2010	001704	0										00
15	LOGICAL-PHYSICAL DEVICE TABLE	ON	170	037510	0										00
16	LOGICAL DEVICE AND CLASS TABLE	ON	1204	061423	3										02
17	DRIVER LINKAGE TABLE	OFF	40	000194	0										00
20	I/O RESOURCE TABLE	OFF	24	000174	0										00
21	DISK FREE SPACE	OFF	20000	040023	1										21
22	LOADER SEGMENT TABLE	ON	3144	174423	1										14
23	TIMER REQUEST LIST	OFF	204	000450	0										00
24	DIRECTORY	ON	2000	007823	1										03
25	DIRECTORY SPACE	ON	600	151423	0										01
26	RIN TABLE	OFF	1304	036823	1										00
27	SWAP TABLE	OFF	7250	024724	0										00
30	JOB PROCESS COUNT	ON	70	030654	0										00
31	JOB MASTER TABLE	ON	700	174023	1										14
32	TAPE LABEL TABLE	OFF	1750		1	17074	D								02
33	LOG TABLE	OFF	170		1	3148	D								00
34	REPLY INFORMATION TABLE	OFF	2000		1	18304	D								03
35	VOLUME TABLE	ON	34	063423	3										01
36	BREAKPOINT TABLE	OFF	674		1	17170	D								01
37	LOG BUFFER 1	OFF	400		1	17174	D								01
40	LOG BUFFER 2	OFF	400		1	17200	D								01
41	LOG ID TABLE	OFF	150		1	3144	D								00
42	ASSOCIATION TABLE	OFF	640		1	17120	D								01
43	CST BLOCK	OFF	44	000220	0										00
44	JOB CUTOFF TABLE	OFF	74	000674	0										00
45	SYSTEM #1	OFF	100	081423	3										01
46	SPECIAL REQUEST TABLE	OFF	144	034174	0										00
47	VIRTUAL DISK SPACE TABLE	OFF	184	035320	0										00
51	ARSBM TABLE	OFF	44	000404	0										00
52	ILI	OFF	1020	021684	0										00
53	SIN TABLE	OFF	170	037700	0										00
54	FILE MULTI-ACCESS VECTOR	ON	200	177423	3										00
55	INPUT DEVICE DIRECTORY	ON	400	144223	1										40
56	OUTPUT DEVICE DIRECTORY	OFF	400	082023	3										40
57	WELCOME MESSAGE #1	OFF	1750		1	17034	D								2

Dump Example

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\*\*\*\*\* DST TABLE \*\*\*\*\*

SEGMENT NUMBER	SEGMENT DESCRIPTION	REFERENCE BIT	SEGMENT LENGTH	ABSOLUTE ADDRESS	BANK / LDEV	DISC ADDRESS	D C C O I S F U R M I N D	VH ALLOC
60	WELCOME MESSAGE #2)	OFF	1750		1	17044	D	2
61	CS SYSTEM SEGMENT)	OFF	1220		1	18170	D	2
62	JOB-PROCESS CROSS REFERENCE)	ON	60	177823	2			1
63	SYSTEM JDT)	OFF	34		1	18340	D	1
64	COMMAND INTERPRETER LOG-ON DST)	OFF	1000		1	17054	D	1
65	IMOUNTED VOLUME TAB)	OFF	520		1	17124	D	1
66	PRI. VOL. USER TABLE)	OFF	200		1	17130	D	10
67	AVAILABLE REGION LIST)	OFF	2004	035504	0			0
70	DISC REQUEST TABLE)	OFF	3120	016544	0			0
71	MSG HBR TABLE)	OFF	380	034340	0			0
72	PRIMARY MSG TABLE)	OFF	200	034720	0			0
73	MEASUREMENT INFO TABLE)	OFF	120	000264	0			0
74	SECONDARY MSG TABLE)	OFF	200	035120	0			0
75		ON	3244	047423	3			0
76		OFF	3674		1	18134	D	7
77		OFF	3804		1	17204	D	7
100		OFF	13144		1	17240	D	16
101		OFF	2554		1	17330	D	16
102		ON	2310	175423	0			13
103		OFF	2260		1	17434	D	13
104		OFF	4770		1	17464	D	13
105		ON	6784	054223	2			43
106		OFF	4720		1	17754	D	17
107		OFF	4010		1	20150	D	5
110		OFF	10374		1	20514	D	37
111		ON	204	075023	3			1
112		ON	1640	172623	0			12
113		OFF	1404		1	20264	D	22
114		OFF	5374		1	20404	D	22
115		ON	104	177823	1			1
116		ON	300	075423	3			5
117		ON	100	165223	3			3
120		ON	2520	127223	3			3
121		ON	500	076023	3			3
122		OFF	2520		1	20324	D	3
123		OFF	404	174623	0		D R	1
124		ON	610	132823	3			1
125		ON	500	151623	2			1
126		ON	610	012023	1			1
127		OFF	2520		1	20344	D	2
130		ON	460	020023	2			1
131		ON	7640	141623	2			10
132		ON	14164	067423	1			54
133		ON	11784	185423	2			54
134		ON	610	175623	3			1
135		ON	610	081423	3			1

# Dump Example

HP3000 III MEMORY DUMPC.00.02 OF SYS VER C UPDATE 01 FIX 08 DUMP TIME 10/16/81, 4.01PM  
 (C) HEWLETT-PACKARD CO. 1980

\*\*\*\*\* DST TABLE \*\*\*\*\*

SEGMENT NUMBER	SEGMENT DESCRIPTION	REFERENCE BIT	SEGMENT LENGTH	ABSOLUTE ADDRESS	BANK/ /LDEV	DISC ADDRESS	D C V	R C I	I N T	S H O W	F P P I S	S S Y S	S E R V I C E	VN ALLO
137		ON	644	173023	1									71



PCB

- \* ONLY ONE PROCESS RUNNING
- \* CI SON WAIT PCB 16
- \* PCB 54 IN PAUSE
- \* PCB 55 IN PAUSE





HP3000 III MEMORY DUMPC.00.02 OF SYS VER C UPDATE DT FIX 09 DUMP TIME 10/18/81, 4:01PM  
 (C) HEWLETT-PACKARD CO. 1980

\*\*\*\*\* PROCESS CONTROL BLOCK (2ND HALF) \*\*\*\*\*

----- SCHEDULING INFORMATION -----														---RESOURCES---				LIFE/DEATH		----- MISCELLANEOUS -----				
PIN	NPIN	POPIN	D P C Q D Q E R E	I C N O R E	H U I S P R D R S I O W W L M P C P P P P S A O V	P R I	P R I	P R I	P R I	P R I	P R I	P R I	P R I	C H R I T R	P R E V P I N	N E X T I M P D P I N	L I F E D E A T H	B M S	P P C	P C S T	P R X P T R	S L L P T R	B P T L N K	S Y S T E M P R O C N A M E
1			L			81											L	SNF	NUL			10	24003	PROGEN
2			L			82											L	SNF	NUL	CST	41		23731	SYSIO
3			L			170											L	SNF	NUL	CST	37		23743	IONESS
4			L			82											L	SNF	NUL			1	23755	LOG
5			L			175											L	SNF	NUL			2	23787	MEMLOG
6			L			175											L	SNF	NUL			3	24001	
7			L			175											L	SNF	NUL			4	24013	UCOP
10			L			12											L	SNF	NUL			5	24025	PFail
11			L			175											L	SNF	NUL			6	24037	DEVREC
12			L			218											L	SNF	NUL			7	24051	LOAD
14			L			230											L	SNF	NUL				24551	
16			C			230											L	SNF	NUL				25073	
54			C			230											L	SNF	NUL	F		15	21171	
55			C			230											L	SNF	NUL			27	21532	

80 ENTRIES  
 41 UNASSIGNED ENTRIES  
 17 ASSIGNED ENTRIES

# Dump Example

HP2000 III MEMORY DUMPC.00.02 OF SYS VER C UPDATE 01 FIX CR DUMP TIME 10/10/81, 4:01PM  
 (C) HEWLETT-PACKARD CO. 1980

\*\*\*\*\* PCBK AND STACK MARKERS FOR DST 132 (PCB 54) \*\*\*\*\*

SEG REL DL	SEG REL DB	JMAT INDEX	JPCNT INDEX	JOB INPUT LOG DEV #	JOB OUTPUT LOG DEV #	JDT DST INDEX	JIT DST INDEX	JOB TYPE	DUPLICAT	INTERACT	INIT Q	JCUT INDEX
001044	001200	1	2	20	20	120110	115	052	YES	YES	000302	0
ADDRESS	BANK	X	DELTA P	STATUS	DELTA Q	SEGMENT						
077238	1	177750	020003	101074	000012	74						
077224	1	000021	010030	141074	000014	74						
077210	1	000337	001725	142430	000020	30						
077170	1	000041	004074	082304	000007	304	USER SEGMENT					
077161	1	000041	000207	082304	000004	304	USER SEGMENT					
077155	1	000000	000053	082304	000024	304	USER SEGMENT					
077131	1	000000	000000	140041	000004	41						

\*\*\*\*\* PCBK AND STACK MARKERS FOR DST 133 (PCB 55) \*\*\*\*\*

SEG REL DL	SEG REL DB	JMAT INDEX	JPCNT INDEX	JOB INPUT LOG DEV #	JOB OUTPUT LOG DEV #	JDT DST INDEX	JIT DST INDEX	JOB TYPE	DUPLICAT	INTERACT	INIT Q	JCUT INDEX
000044	001000	1	2	20	20	120110	115	052	YES	YES	000302	0
ADDRESS	BANK	X	DELTA P	STATUS	DELTA Q	SEGMENT						
175036	2	177750	020003	101074	000012	74						
175024	2	170520	010030	141074	000014	74						
175010	2	000337	001725	142430	000020	30						
174770	2	000041	004074	082304	000007	304	USER SEGMENT					
174761	2	000041	000207	082304	000004	304	USER SEGMENT					
174755	2	000000	000053	082304	000024	304	USER SEGMENT					
174731	2	000000	000000	140041	000004	41						

HP1000 III MEMORY DUMP: 00 02 OF SYS VER C UPDATE 01 FIX 00 DUMP TIME 10/16/81, 4:01PM  
(C) HEWLETT-PACKARD CO. 1980

\*\*\*\*\* SIX TABLE \*\*\*\*\*

NO LOCKED S1AS

\*\*\*\*\* MONITOR TABLE \*\*\*\*\*

LOCATION	PIN	EVENT	PIN	EVENT	PIN	EVENT
40735	0	QUIESCE	031171	000200	122238	0 QUIESCE
40731	0	QUIESCE	031171	004000	122234	0 QUIESCE
40705	0	INTERRUPT	001132	000000	043177	0 INTERRUP
40471	54	SPECIALRQ	000121	141100	000001	54 SPECIALRQ
40455	0	SPECIALRQ	000132	000023	000000	0 SPECIALRQ
40441	54	SIDOMEEXIT	001420	000413	001052	54 SIDOMEEXIT
40425	0	SIDOME	000126	018024	000000	0 SIDOME
40411	0	SIDOMEEXIT	001260	000413	002770	0 SIDOMEEXIT
40375	0	ALLOCMEM	000004	000001	012023	0 ALLOCMEM
40361	54	QONSEG	000128	031171	000017	54 QONSEG
40345	0	INTERRUPT	001132	000000	042745	0 INTERRUPT
40331	54	SPECIALRQ	000132	005100	000001	54 SPECIALRQ
40315	0	INTERRUP	001132	000000	042717	0 INTERRUPT
40301	54	SIDOMEEXIT	001340	000413	002882	54 SIDOMEEXIT
40285	0	QUIESCE	031171	004000	122234	0 QUIESCE
40271	0	QUIESCE	031171	004000	122232	0 QUIESCE
40255	0	INTERRUPT	001132	000000	042330	0 INTERRUPT
40241	0	QUIESCE	031171	002000	122232	0 QUIESCE
40225	0	DEALLOC	000000	000001	107023	0 DEALLOC
40211	0	DEALLOC	000000	000001	107023	0 DEALLOC
40195	54	SIDOMEEXIT	001000	000000	001852	54 SIDOMEEXIT
40181	0	INTERRUPT	001132	000000	041645	0 INTERRUPT
40165	0	COMBAGE	000002	000200	000000	0 COMBAGE
40151	54	SIDOMEEXIT	001200	000413	001813	54 SIDOMEEXIT
40135	0	SIDOME	000132	018144	000000	0 SIDOME
40121	0	SEGIO	000132	015724	100001	0 SEGIO
40105	0	SWAPIN	000054	100000	000000	0 SWAPIN
40091	0	DEALLOC	000002	000001	077023	0 DEALLOC
40075	0	MAKEOC	104402	000000	000000	0 MAKEOC
40061	0	SEGIO	000132	015724	100001	0 SEGIO
40045	54	SPECIALRQ	000132	000011	000200	54 SPECIALRQ
40031	0	INTERRUPT	001132	000000	041471	0 INTERRUPT
40015	0	SEGIO	000134	015706	000001	0 SEGIO
40001	0	FETCHSEG	000134	000054	000003	0 FETCHSEG
39985	0	SIDOMEEXIT	001000	000000	131412	0 SIDOMEEXIT
39971	0	QUIESCE	031171	004000	122232	0 QUIESCE
39955	0	SIDOMEEXIT	001000	000000	131384	0 SIDOMEEXIT
39941	0	QUIESCE	031171	004000	122232	0 QUIESCE
39925	54	SPECIALRQ	000132	043480	000001	54 SPECIALRQ
39911	0	SPECIALRQ	000121	000003	000000	0 SPECIALRQ
39895	54	SIDOMEEXIT	001400	000413	001432	54 SIDOMEEXIT
39881	0	QUIESCE	031171	002000	122230	0 QUIESCE
39865	0	INTERRUPT	001132	000000	041410	0 INTERRUPT
39851	54	SPECIALRQ	000132	005100	000001	54 SPECIALRQ
39835	0	INTERRUP	001132	000000	041382	0 INTERRUPT
39821	54	SIDOMEEXIT	001240	000413	001327	54 SIDOMEEXIT
39805	0	QUIESCE	031171	004000	122230	0 QUIESCE
39791	0	INTERRUP	001132	000000	041301	0 INTERRUPT
39775	0	QUIESCE	031171	002000	122230	0 QUIESCE
39761	54	SPECIALRQ	000124	000023	000000	54 SPECIALRQ
39745	0	QUIESCE	031171	004000	122230	0 QUIESCE
39731	54	SIDOMEEXIT	001000	000000	130750	54 SIDOMEEXIT

Dump Example

HP3000 III MEMORY DUMPC. 09 02 OF SYS VER C UPDATE 01 FIX 00 DUMP TIME 10/10/81. 4:01PM  
 (C) HEWLETT-PACKARD CO. 1980

41711	0 INTERRUPT	001132	000000	040747	0 QUIESCE	031171	004000	122230	54 SIOOMEKIT	001100	000413	000744
41875	54 SPECIALRQ	000132	141040	000001	0 SIOOMEKIT	001000	000000	130722	0 SPECIALRQ	000132	000023	000000
41861	0 INTERRUPT	001132	000000	040720	0 QUIESCE	031171	004000	122230	54 SIOOMEKIT	001440	000413	000705
41845	54 SPECIALRQ	000132	064500	000001	0 SIOOMEKIT	001000	000000	130834	0 SPECIALRQ	000132	000023	000000
41831	0 INTERRUPT	001132	000000	040872	0 QUIESCE	031171	004000	122230	54 SIOOMEKIT	001100	000413	130443
41615	54 SIOOMEKIT	001300	000413	000842	54 SPECIALRQ	000132	043400	000001	0 QUIESCE	031171	004000	122230
41601	0 SIOOMEKIT	001000	000000	130484	0 SPECIALRQ	000132	000003	000000	0 INTERRUPT	001132	000000	040482
41585	0 QUIESCE	031171	004000	122230	54 SIOOMEKIT	001040	000413	000420	54 SPECIALRQ	000121	177780	000001
41551	0 QUIESCE	031171	004000	122230	0 QUIESCE	024013	000400	140175	0 SIOOMEKIT	001000	000000	131124
41535	0 SIOOMEKIT	000031	015544	000000	0 INTERRUPT	001132	000000	021132	0 SWAPIN	000007	100000	000000
41521	0 SIOOMEKIT	001020	000413	001105	0 SEQIO	000031	015544	000001	0 ALLOCHEM	000002	000001	174023
41505	0 FETCHSEQ	000031	000007	000003	0 QUIESCE	024013	000001	140175	7 QOWSEQ	000001	024013	000023
41471	0 SIOOMEKIT	001000	000000	131077	0 SIOOMEKIT	102001	020124	000000	0 INTERRUPT	001132	000000	021075
41455	0 SWAPIN	000007	100000	000000	0 SIOOMEKIT	001300	000413	001031	0 SEQIO	102001	020124	000001
41441	0 ALLOCHEM	000012	000003	102823	0 FETCHSEQ	102001	000007	000003	0 QUIESCE	024013	000001	140175
41425	7 QOWSEQ	102001	024013	000030	0 SWAPIN	000007	100000	000000	0 DEALLOC	000000	000000	174023
41411	0 FETCHSEQ	000102	000007	000001	0 FETCHSEQ	000011	000007	000001	0 QUIESCE	031532	000200	122234
41375	0 SIOOMEKIT	001000	000000	135405	0 SPECIALRQ	000125	000003	000000	0 INTERRUPT	001132	000000	015403
41361	0 QUIESCE	021532	002000	122234	53 SIOOMEKIT	001020	000413	005371	53 SPECIALRQ	000125	000000	000001
41345	0 SIOOMEKIT	001000	000000	135202	0 SPECIALRQ	000125	000003	000000	0 INTERRUPT	001132	000000	013300
41331	0 QUIESCE	021532	002000	122234	53 SIOOMEKIT	001700	000413	005244	53 SPECIALRQ	000125	141100	000001
41315	53 SIOOMEKIT	001000	000000	005227	0 SIOOMEKIT	001000	000000	135223	0 SPECIALRQ	000133	000023	000000
41301	0 INTERRUPT	001132	000000	015221	0 QUIESCE	031532	004000	122234	53 SIOOMEKIT	001740	000413	005172
41285	53 SPECIALRQ	000133	001520	000001	0 SIOOMEKIT	001000	000000	135157	0 SIOOMEKIT	000124	020824	000000
41251	0 INTERRUPT	001132	000000	015150	0 SWAPIN	000055	100000	000000	0 ALLOCHEM	000100	000413	005113
41235	0 SEQIO	000124	020824	000001	0 DEALLOC	000000	000003	132023	0 ALLOCHEM	000004	000003	132023
41221	0 FETCHSEQ	000124	000055	000003	0 QUIESCE	021532	000001	122234	53 QOWSEQ	000124	021532	000017
41205	0 SIOOMEKIT	001000	000000	135072	0 SPECIALRQ	000133	000023	000000	0 INTERRUPT	001132	000000	015070
41171	0 QUIESCE	021532	004000	122234	53 SIOOMEKIT	001720	000413	005056	53 SPECIALRQ	000133	005100	000001
41155	0 SIOOMEKIT	001000	000000	135044	0 SPECIALRQ	000133	000023	000000	0 INTERRUPT	001132	000000	013042
41141	0 QUIESCE	021532	004000	122234	53 SIOOMEKIT	001000	000413	005004	53 SPECIALRQ	000133	043400	000001
41125	0 QUIESCE	031171	000200	122230	54 SIOOMEKIT	001000	000000	134715	54 SPECIALRQ	000125	000003	000000
41111	54 INTERRUPT	001132	000000	014713	0 QUIESCE	031532	002000	122232	53 SIOOMEKIT	001000	000413	004457
41075	53 SPECIALRQ	000125	000000	000001	0 QUIESCE	031171	004000	122230	0 QUIESCE	021532	000000	122232
41061	53 SIOOMEKIT	001000	000000	004524	0 SIOOMEKIT	001000	000000	134520	0 SPECIALRQ	000133	000023	000000
41045	0 INTERRUPT	001132	000000	014516	0 QUIESCE	031532	004000	122232	53 SIOOMEKIT	001020	000413	004451
41031	53 SPECIALRQ	000133	001520	000001	0 SIOOMEKIT	001000	000000	134431	0 SIOOMEKIT	000137	020444	000000
41015	0 INTERRUPT	001132	000000	014427	0 SWAPIN	000055	100000	000000	0 SIOOMEKIT	001700	000413	004400
41001	0 SEQIO	000127	020444	000001	0 DEALLOC	000000	000001	174023	0 ALLOCHEM	000004	000001	173023
40765	0 FETCHSEQ	000137	000055	000003	0 QUIESCE	021532	000001	122232	53 QOWSEQ	000137	021532	000011
40751	0 SIOOMEKIT	001000	000000	134380	0 SIOOMEKIT	000135	020404	000000	0 INTERRUPT	001132	000000	014358

HP0000 III MEMORY DUMPC. 00. 02 OF SYS VER C UPDATE 01 FIX 09 DUMP TIME 10/18/81. 4:01PM  
 (C) HEWLETT-PACKARD CO. 1980

\*\*\*\*\* TIMER REQUEST LIST \*\*\*\*\*

FREE LIST POINTER 000034  
 NUMBER OF ENTRIES 000040  
 ENTRY SIZE 4  
 TRACE WORD 010020  
 QUANTUM/100MS 000000  
 POINTER TO MOST ACTIVE REQ 000024  
 DATE 10/18/81. 4:01PM

ENTRY	REQUEST STATUS	TYPE OF REQUEST	POINTER TO NEXT REQUEST	REQUEST POINTER	TIME TO SERVICE REQ IN FRONT (SEC/10)
18	ACTIVE	DELAY	0	PCBB IX = 000120	15532
20	ACTIVE	32TURNAROUND	14	DITP = 001320	8
24	ACTIVE	32TURNAROUND	30	DITP = 001320	9883
30	ACTIVE	HANGUP	20	DITP = 000001	107
34	INACTIVE	HANGUP	40	DITP = 000000	0
40	INACTIVE	HANGUP	44	DITP = 000000	0
44	INACTIVE	HANGUP	50	DITP = 000000	0
50	INACTIVE	HANGUP	54	DITP = 000000	0
54	INACTIVE	HANGUP	58	DITP = 000000	0
60	INACTIVE	HANGUP	64	DITP = 000000	0
64	INACTIVE	HANGUP	70	DITP = 000000	0
70	INACTIVE	HANGUP	74	DITP = 000000	0
74	INACTIVE	HANGUP	100	DITP = 000000	0
100	INACTIVE	HANGUP	104	DITP = 000000	0
104	INACTIVE	HANGUP	110	DITP = 000000	0
110	INACTIVE	HANGUP	114	DITP = 000000	0
114	INACTIVE	HANGUP	120	DITP = 000000	0
120	INACTIVE	HANGUP	124	DITP = 000000	0
124	INACTIVE	HANGUP	130	DITP = 000000	0
130	INACTIVE	HANGUP	134	DITP = 000000	0
134	INACTIVE	HANGUP	140	DITP = 000000	0
140	INACTIVE	HANGUP	144	DITP = 000000	0
144	INACTIVE	HANGUP	150	DITP = 000000	0
150	INACTIVE	HANGUP	154	DITP = 000000	0
154	INACTIVE	HANGUP	160	DITP = 000000	0
160	INACTIVE	HANGUP	164	DITP = 000000	0
164	INACTIVE	HANGUP	170	DITP = 000000	0
170	INACTIVE	HANGUP	174	DITP = 000000	0
174	INACTIVE	HANGUP	200	DITP = 000000	0



Dump Example

SYSTEM GLOBAL PORT DATA SEGMENT

\* FIND %1377 SYSTEM GLOBAL (WHICH IS ADDRESSED TO SYSTEM GLOBAL  
EXTENSION)

\* ADD %1000 TO IT, SYSGLOB

\* PORT DATA SEGMENT IN SYSTEM GLOBAL EXTENSION + %100

SYSTEM GLOBAL EXTENSION CELL + %1100 = DST # FOR PORT DATA  
SEGMENT.



HP3000 III MEMORY DUMPC.00.02 OF SYS VER C UPDATE 01 FIX 09 DUMP TIME 10/18/81. 4:01PM  
(C) HEWLETT-PACKARD CO. 1980

BANK 0

000404: 070000 040000 000000 000000 000000 040000 000000 000000 000414: 000000 000000 000000 000000 000000 000000 000000 000000  
000424: 000000 000000 000000 000000 000000 000000 000000 000000 000434: 000000 000000 000000 000000 000000 000000 000000 000000  
000444: 000000 000000 000000 000000

\$\$\$\$\$\$ DST 23 (TIMER REQUEST LIST) \$\$\$\$\$\$

000450: 000034 020004 010020 000000 177747 001500 050300 121441 000460: 000024 000000 000000 000000 112000 000120 000000 036254  
000470: 124014 001300 000000 000000 124030 001320 000000 023207 000500: 020020 000001 000000 000153 000040 000000 000000 000000  
000510: 000044 000000 000000 000000 000050 000000 000000 000000 000520: 000054 000000 000000 000000 000000 000000 000000 000000  
000530: 000084 000000 000000 000000 000070 000000 000000 000000 000540: 000074 000000 000000 000000 000100 000000 000000 000000  
000550: 000104 000000 000000 000000 000110 000000 000000 000000 000560: 000114 000000 000000 000000 000120 000000 000000 000000  
000570: 000124 000000 000000 000000 000130 000000 000000 000000 000580: 000134 000000 000000 000000 000140 000000 000000 000000  
000580: 000144 000000 000000 000000 000150 000000 000000 000000 000590: 000154 000000 000000 000000 000160 000000 000000 000000  
000590: 000164 000000 000000 000000 000170 000000 000000 000000 000600: 000174 000000 000000 000000 000200 000000 000000 000000  
000650: 000000 000000 000000 000000

\$\$\$\$\$\$ DST 30 (JOB PROCESS COUNT) \$\$\$\$\$\$

000654: 011422 000370 177376 177376 177376 177376 177376 177376 000664: 177376 177376 177377 000000 000000 000000 000000 000000

\$\$\$\$\$\$ DST 44 (JOB CUTOFF TABLE) \$\$\$\$\$\$

000674: 000000 011400 000071 000000 000000 000000 000011 000000 000704: 000000 000014 000000 000000 000017 000000 000000 000022  
000714: 000000 000000 000025 000000 000000 000030 000000 000000 000724: 000020 000000 000000 000000 000030 000000 000041 000000  
000734: 000000 000044 000000 000000 000047 000000 000000 000000 000744: 000000 000000 000000 000000 000000 000000 000000 000000  
000754: 000000 000000 000000 000000 000000 000000 000071 000000 000764: 000000 000000 000000 000000 000004 000001 030370 030370  
000774: 000000 030370 030370 030370

\$\$\$\$\$\$ DST 5 (SYSTEM GLOBAL AREA) \$\$\$\$\$\$

001000: 000000 000554 005314 012014 177404 014514 021704 013514 001010: 000510 000000 177450 177874 036700 177654 000704 037071  
001020: 000000 128023 001560 010733 000000 023724 034320 000004 001030: 000000 015544 001440 000040 000000 003724 000000 000000  
001040: 034340 034360 033174 000000 034504 000121 000000 000003 001050: 000001 177220 000400 010001 000400 177174 177202 177210  
001060: 000000 000000 000000 000000 000000 000000 000000 000000 001070: 000000 000020 000000 000001 000024 018015 000111 000000  
001100: 000001 172000 000000 000000 000800 040000 000077 075000 001110: 002280 077777 000004 000107 030111 030071 020103 000000  
001120: 000170 000000 103440 007232 000430 002340 000400 000100 001130: 000000 000130 177770 102115 000000 101117 000000 000000  
001140: 104110 000020 000000 000100 000200 000220 000000 000000 001150: 000100 000240 000000 000040 000120 100441 000000 101415  
001160: 001864 104054 101554 001245 002343 111054 004700 007703 001170: 000000 000000 002037 000040 000002 000282 000000 000000  
001200: 000000 000000 000000 000000 001777 000284 000000 000001 001210: 000000 000000 000000 000000 000001 000000 000000 000000  
001220: 000304 000000 000000 000001 000001 031344 002001 000787 001230: 000400 000200 000000 001300 000713 048517 051122 044523  
001240: 020040 044120 042440 020040 020040 000000 000000 000000 001250: 000000 000000 000000 000000 000000 000000 000000 000000  
001260: 000000 177284 000000 000000 000000 000001 104223 000000 001270: 000000 000000 000000 000000 000000 000000 000000 000000

Dump Example

HP3000 III MEMORY DUMP.00.02 OF SYS VER C UPDATE 01 FIX 00 DUMP TIME 10/18/81. 4:01PM BANK . 0  
(C) HEWLETT-PACKARD CO. 1980

001300: 000000 000000 000005 000000 004402 000020 003400 000020 001310: 002001 000020 001403 000020 002404 000020 000000 000000  
001320: 000000 000000 000000 112132 000000 000000 000000 000000 001330: 000000 000000 000000 000310 000310 104130 110530 106130  
001340: 100527 105532 101134 100534 100531 000000 000001 000000 001350: 000000 000000 022855 000000 001750 000001 000000 000000  
001360: 000000 000003 177777 000140 000000 000000 000000 000000 001370: 000000 000401 137724 000110 100547 101144 000000 000000  
001400: 000000 000000 000000 000000 000000 000000 000000 000000 001410: 177520 000000 034172 000000 000000 000000 000000 000000  
001420: 000005 000012 000047 000300 000000 000000 000000 000000 001430: 000000 000000 000000 000000 000000 000000 000000 000000  
001440: 030370 030370 030370 030370 030370 030370 030370 030370 001450: 030370 030370 030370 030370 030370 030370 030370 030370  
LINES 001400 - 001477 SAME AS ABOVE  
% 1604 BEGIN 2ND GLOB EXTENSION  
001500: 000370 030370 030370 030370 000000 000000 120023 000001 001510: 000000 000200 000000 000000 001520 041504 000764 000000 000000  
001520: 000000 000000 000000 000000 000000 000000 000000 000000 001530: 000000 000500 000000 000000 000000 000000 000000 000000  
LINES 001500 - 001577 SAME AS ABOVE  
% 1604 PORT DATA SEGMENT NUMBER = 137  
001580: 000000 000000 000000 000000 110547 005254 104472 002876 001570: 005202 100072 000000 000000 000000 000000 000000 000000  
001600: 020103 030111 030071 000000 000000 000000 000000 000000 001610: 000000 000000 000000 000000 000000 000000 000000 000000  
001620: 000000 000000 000000 000000 000000 000000 000000 000000 001630: 000000 000000 000000 000000 000000 000000 000000 000000  
LINES 001600 - 001677 SAME AS ABOVE  
001700: 000000 000000 000000 000000

##### DST 14 I TERMINAL BUFFERS #####  
001700: 000000 040070 000020 001750 007730 004400 000000 000000 005078 000030 005040 020040 020445 001704: 08... 15  
001720: 000014 020040 020040 020113 048117 047105 008412 020040 020081 033040 020040 020040 051473 001720: 08... 16 S:  
001734: 000030 000050 020508 020113 030415 005040 020040 030470 020040 020040 020120 040525 051505 001734: 08... 17 PAUSE  
001750: 000044 020081 030015 005040 020040 000070 030470 020040 020040 020103 008412 002440 020040 001750: 10... 18 C  
001764: 000080 030471 020040 020040 020117 050105 047040 048523 043481 000110 035501 036481 008412 001764: 10... 19 OPEN MSG1.H.A+1..  
002000: 000074 020040 020082 020040 020040 020040 053522 044524 042415 005040 020040 031008 020040 002000: 20... 20 WRITE... 21  
002014: 000110 000130 020040 020103 008412 020040 020082 031040 020040 020040 047520 042518 020115 002014: X G... 22 OPEN M  
002030: 000124 051507 051415 005040 020040 000190 002402 031440 020040 020040 051505 052111 047124 002030: 30... 23 SETI  
002044: 000140 020115 051507 031415 005040 020040 031084 020040 020040 000170 020103 047518 052122 002044: MSG3... 24 X CONTR  
002058: 000154 047514 020115 051507 031473 041475 022073 050075 020480 030080 008412 020040 020082 002058: 0L MSG3.C:4.P=1000... 2  
002074: 000170 000210 032440 020040 020040 051105 040504 020115 051507 031415 005040 020040 002482 002074: .S READ MSG3... 2  
002110: 000204 033040 020040 020040 041415 000230 005040 020040 031087 020040 020040 020117 050105 002110: 6... 27 OPE  
002124: 000220 047040 048523 043484 008412 020040 020082 030400 020040 006250 020040 051505 052111 002124: N MSG4... 28 SETI  
002140: 000234 047124 020115 051507 032015 005040 020040 031071 020040 020040 020122 042501 042040 002140: NY MSG4... 29 READ  
002154: 000250 000270 048523 043484 008412 020040 020082 032400 020040 020040 020120 040525 051505 002154: MSG4... 30 PAUSE  
002170: 000264 020081 020080 030015 005040 000210 020040 031481 020040 020040 020108 047518 008412 002170: 1000... 31 EOF  
002204: 000300 027480 044524 027483 030080 030040 020108 051111 020040 000330 048448 052040 030486 002204: /ORT/3000 FBI... 32 X ? IS  
002220: 000314 020040 030471 030408 020040 020083 035085 034440 050115 008412 024103 024440 044105 002220: 1981.3.58 PM.(C) HE  
002234: 000330 000350 048481 008412 052055 050110 041513 040522 042040 041517 027040 030471 002470 002234: .ML. HP-PACKARD CO. 19...  
002250: 000344 030015 005080 030040 030080 000410 040401 020101 040122 042501 062131 020108 054111 002250: 0... 00 00. M1 ALREADY EXI  
002264: 000360 051524 051440 020440 051109 051520 047518 042040 054505 000530 020015 008075 020120 002264: STS - RESPOND YE... P  
002300: 000374 052522 043505 020115 051507 032015 005040 020040 030481 020040 002440 020040 041125 002300: URGE MSG4... 11 BU  
002314: 000410 000435 051440 052117 020120 052522 043505 020117 048104 020101 047104 020113 042505 002314: .S TO PURGE OLD AND KE  
002330: 000424 050040 047105 053415 005120 000450 052522 043505 020117 048104 037480 030080 030086 002330: P NEW. P URGE OLD TOOOO  
002344: 000440 030440 030480 033483 030084 008412 042101 030482 032453 000470 054720 027481 031000 002344: I 107304. 04125... 12  
002360: 000454 020060 030060 030084 033440 030060 030080 030084 020061 033487 033487 033460 030060 002360: 000047 000004 177777 00  
002374: 000470 000519 027480 020040 034040 020040 020040 041125 044514 042040 048523 043460 035515 002374: M/O 8 BUILD MSG3.M  
002410: 000504 051507 008412 020081 032487 000370 042440 020040 034440 020040 020040 043111 046105 002410: SG 177. E 9 FILE  
002424: 000520 020115 051507 031454 047514 042015 005040 020040 030480 000550 042530 044524 048523 002424: MSG3 OLD... 10. HEXITMS  
002440: 000534 043484 035515 051507 008412 020040 020081 031040 020040 020040 043111 049105 020115 002440: G4 MSG... 12 FILE M  
002454: 000550 008570 008412 032054 047514 042015 005057 030080 020060 030080 030060 022040 030060 002454: .X. 4. OLD... 100 000004 00



\*SPCBDELAYSOFT PCB (8:1) 000100 = 0 000 000 001 000 000  
0 1 4 7 10 13

NOT SET

PROCESS IS CRITICAL, HOLDING SIR, IMPEDED

\*SPCBWAKESOFT PCB 8.(3:1) 172400 = 1 111 010 100 000 000  
0 1 4 7 10 13

SET 172000 = 1 111 010 000 000 000  
PROCESS ALLOWS SOFT INT 0 1 4 7 10 13  
EVEN IF WAITING ON OTHER EVENTS

\*SPCBSOFTINT PCB 9.(9:1) 100200 = 1 000 000 010 000 000  
0 1 4 7 10 13  
NOT SET 100000 = 1 000 000 000 000 000  
0 1 4 7 10 13

SET TO INFORM DISP/PSUEDOINT  
SHOULD BE AWARE OF SOFT INT

\*SPCBALLOWSOFT PCB 13.(7:1) 02220 = 0 010 010 010 011 000  
0 1 4 7 10 13

NOT SET

SET, THE PROCESS WILL PROCESS  
SOFT INT (FINSTATE)

Dump Example

HP1000 III MEMORY DUMPC.00.02 OF SYS VER C UPDATE 01 FIX 08 DUMP TIME 10/15/81. 4:01PM BANK 0  
 (C) HEWLETT-PACKARD CO. 1980  
 012414: 100540 000000 000401 041378 100244 000000 000401 041411 012424: 100570 000000 000401 041417 100480 000000 000401 041433  
 012434: 100508 000000 000401 041445

012440: 000004 125252 000001 000000 021808 000000 000003 152723 012450: 021153 000000 000003 037223 021051 000000 000000 152423  
 012480: 021485 000000 000002 030623 000004 125252 000001 000000 012470: 021808 000000 000001 080223 021153 000000 000000 144423  
 012500: 021051 000000 000002 043023 021485 000000 000002 075823 012510: 100000 001140 000000 000000 100000 001144 000000 000000  
 012520: 100000 001150 000000 000000 100000 001154 000000 000000 012530: 100000 001160 000000 000000 100000 001164 000000 000000  
 012540: 100000 001170 000000 000000 100000 001174 000000 000000 012550: 100000 001200 000000 000000 100000 001204 000000 000000  
 012560: 100000 001210 000000 000000 100000 001214 000000 000000 012570: 100000 001220 000000 000000 100000 001224 000000 000000  
 012800: 100000 001230 000000 000000 100000 001234 000000 000000 012810: 100000 001240 000000 000000 100000 001244 000000 000000  
 012820: 100000 001250 000000 000000 100000 001254 000000 000000 012830: 100000 001260 000000 000000 100000 001264 000000 000000  
 012840: 100000 001270 000000 000000 100000 001274 000000 000000 012850: 100000 001300 000000 000000 100000 001304 000000 000000  
 012860: 100000 001310 000000 000000 100000 001314 000000 000000 012870: 100000 001320 000000 000000 100000 001324 000000 000000  
 012700: 100000 001330 000000 000000 100000 001334 000000 000000 012710: 100000 001340 000000 000000 100000 001344 000000 000000  
 012720: 100000 001350 000000 000000 100000 001354 000000 000000 012730: 100000 001360 000000 000000 100000 001364 000000 000000  
 012740: 100000 001370 000000 000000 100000 001374 000000 000000 012750: 100000 001400 000000 000000 100000 001404 000000 000000  
 012760: 100000 001410 000000 000000 100000 001414 000000 000000 012770: 100000 001420 000000 000000 100000 001424 000000 000000  
 013000: 100000 001430 000000 000000 100000 001434 000000 000000 013010: 100000 000000 000000 000000

DUPLICATE DST 3 (PROCESS CONTROL BLOCK) \*\*\*\*\*

013014: 000000 024083 000000 004300 000400 000002 000000 000000 013024: 100000 101000 000000 000000 000000 040012 000000 000000  
 013034: 000000 023731 100000 003640 000400 000400 001400 000000 013044: 100000 101000 000000 000000 000000 040012 000000 000000  
 013054: 000000 023743 000000 003700 002300 000400 002000 000000 013064: 100000 101000 000000 000000 000000 040012 000000 000000  
 013074: 000000 023755 000000 003740 003400 000400 002400 000000 013084: 100000 101000 000000 000000 000000 040012 000000 000000  
 013114: 000000 023787 000000 004000 000010 000400 003000 000000 013124: 100000 101000 000000 000000 000000 040012 000000 000000  
 013134: 000000 024001 000000 004040 000400 000400 003400 000000 013144: 100000 101000 000000 000000 000000 040012 000000 000000  
 013154: 000000 024013 000440 004100 000400 000410 004000 000000 013164: 100000 101400 000000 000000 000000 040012 000000 000000  
 013214: 100000 024025 000000 004140 000401 000400 004400 000000 013224: 100000 101000 000000 000000 000000 040012 000000 000000  
 013234: 000000 024037 000000 004200 002400 000400 005000 000000 013244: 100000 101000 000000 000000 000000 040012 000000 000000  
 013254: 000000 024051 000000 004240 000400 000400 000000 000000 013264: 100000 101000 000000 000000 000000 040012 000000 000000  
 013274: 000000 000320 000440 004400 001000 000400 000000 000000 013284: 100000 101000 000000 000000 000000 040012 000000 000000  
 013314: 000000 024551 000000 004600 000000 000400 000000 000000 013324: 100000 101000 000000 000000 000000 040012 000000 000000  
 013334: 000000 000360 000440 004420 001000 000400 000000 000000 013344: 100000 100400 000000 000000 000000 040012 000000 000000  
 013354: 000000 024572 000000 004420 000000 000400 000000 000000 013364: 100000 100400 000000 000000 000000 040012 000000 000000  
 013374: 000000 000400 000440 005000 001000 000400 000000 000000 013384: 100000 000400 000000 000000 000000 040012 000000 000000  
 013414: 000000 000440 000440 005500 001000 007000 000000 000000 013424: 100000 000200 000000 000000 000000 040012 000000 000000  
 013434: 000000 000460 000440 005500 001000 007000 000000 000000 013444: 100000 000200 000000 000000 000000 040012 000000 000000  
 013454: 000000 000420 000440 005200 001000 010400 000000 000000 013464: 100000 000000 000000 000000 000000 040012 000000 000000  
 013474: 000000 000520 000440 005500 001000 007000 000000 000000 013484: 100000 000200 000000 000000 000000 040012 000000 000000  
 013514: 000000 000560 000440 005500 001000 007000 000000 000000 013524: 100000 000200 000000 000000 000000 040012 000000 000000  
 013534: 000000 000500 000440 005540 001000 012000 000000 000000 013544: 100000 000000 000000 000000 040012 000000 000000 000000  
 013554: 000000 000800 000440 005500 001000 007000 000000 000000 013564: 100000 000200 000000 000000 040012 000000 000000 000000  
 013574: 000000 000540 000440 005240 001000 013000 000000 000000 013584: 100000 000000 000000 000000 040012 000000 000000 000000  
 013614: 000000 000120 000440 005500 001000 007000 000000 000000 013624: 100000 000200 000000 000000 000000 040012 000000 000000  
 013634: 000000 000640 000440 005500 001000 007000 000000 000000 013644: 100000 000200 000000 000000 000000 040012 000000 000000

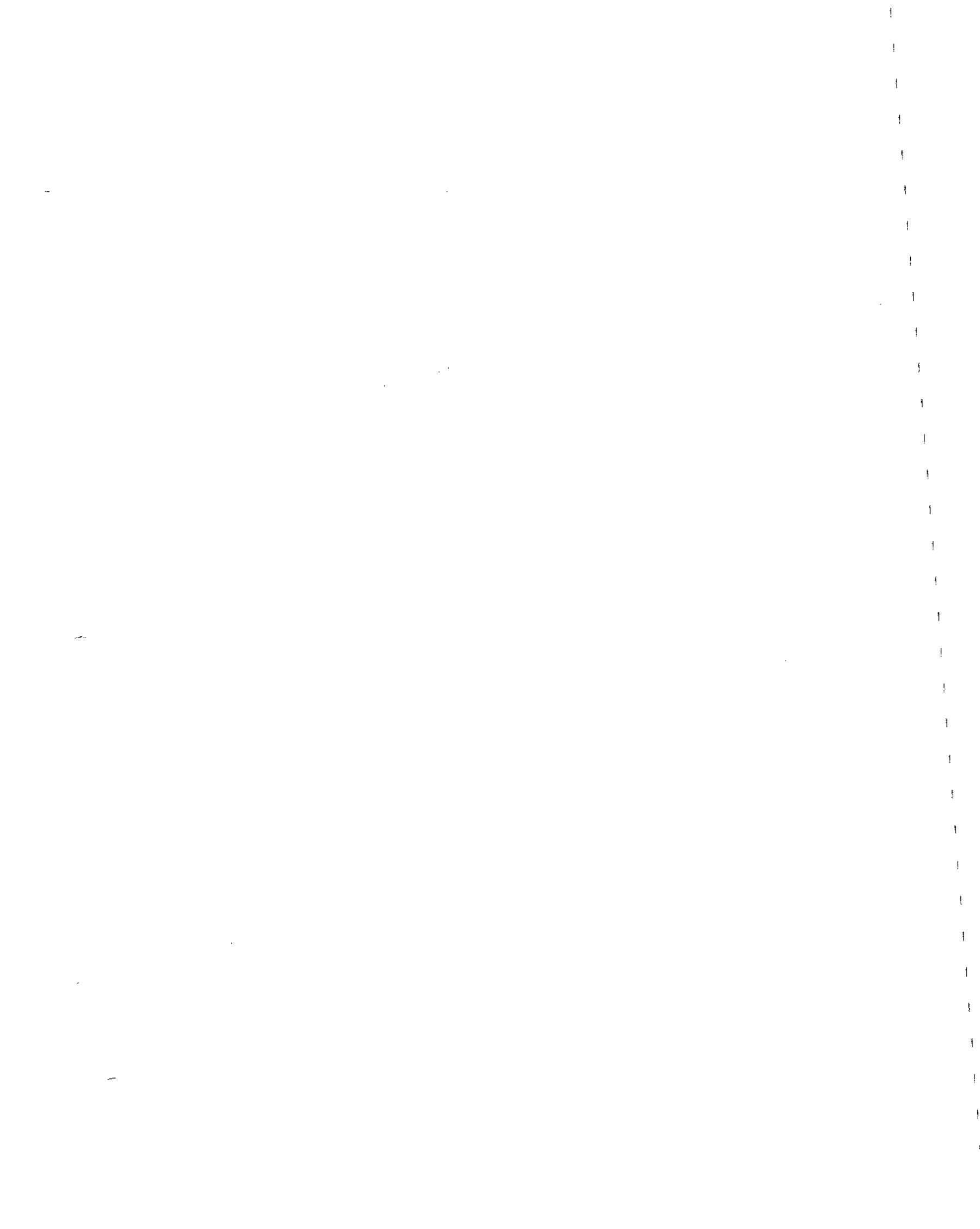
HP3000 III MEMORY DUMPC.00.02 OF SYS VER C UPDATE 01 FIX 06 DUMP TIME 10/19/81, 4:01PM  
(C) HEWLETT-PACKARD CO. 1980

BANK 0

013654:	000000	000800	000000	000000	000000	000000	000000	000000	013664:	000500	000000	000000	000000	000000	000000	000000	177777
013674:	000000	000700	000440	005500	001000	007000	000000	000000	013704:	182400	000200	000000	000000	000013	022230	000000	177777
013714:	000000	000740	000440	005500	001000	007000	000000	000000	013724:	182400	000200	000000	000000	000020	022230	000000	177777
013734:	000000	000780	000440	005500	001000	007000	000000	000000	013744:	182400	000200	000000	000000	000015	022230	000000	177777
013754:	000100	000720	000000	005200	000040	018400	000000	000000	013764:	043400	000000	000000	000000	000027	022230	000000	177777
013774:	000000	001000	000440	005500	001000	007000	000000	000000	014004:	182400	000200	000000	000000	000013	022230	000000	177777
014014:	000000	001040	000440	005500	001000	007000	000000	000000	014024:	182400	000200	000000	000000	000011	022230	000000	177777
014034:	000100	001120	000000	005520	000040	007000	000000	000000	014044:	043400	000200	000000	000000	000015	022230	000000	177777
014054:	000100	001020	000000	005360	000040	020400	000000	000000	014064:	043000	000000	000000	000000	000027	022230	000000	177777
014074:	000010	001220	000440	005500	001000	007000	000000	000000	014104:	182400	000200	000000	000000	000015	022230	000000	177777
014114:	000100	001180	000000	005220	000040	021400	000000	000000	014124:	043000	020000	000000	000000	000027	022230	000000	177777
014134:	000010	001100	000440	005540	001000	021400	000000	000000	014144:	182400	000000	000000	040044	000011	022230	000000	177777
014154:	000100	001080	000000	005020	000040	021400	000000	000000	014164:	043000	020000	000000	000000	000027	022230	000000	177777
014174:	000000	001140	000440	005500	001000	021400	000000	000000	014204:	182400	000000	000000	040044	000011	022230	000000	177777
014214:	000000	001280	000440	005500	001000	007000	000000	000000	014224:	182400	000200	000000	000000	000015	022230	000000	177777
014224:	000100	001200	000000	005360	000040	024000	000000	000000	014244:	053400	000000	000000	000000	000027	022230	000000	177777
014234:	000000	000000	000440	005500	001000	007000	000000	000000	014264:	182400	000200	000000	000000	000015	022230	000000	177777
014274:	000000	001240	000440	005500	001000	007000	000000	000000	014304:	182400	000000	000000	040044	000011	022230	000000	177777
014314:	000010	001170	000000	005520	000040	007000	000000	000000	014324:	182400	000200	000000	000000	000015	022230	000000	177777
014334:	000100	001340	000000	005360	000040	024000	000000	000000	014344:	182400	000200	000000	000000	000015	022230	000000	177777
014354:	100000	001380	001320	000000	000000	000000	000000	000000	014364:	000000	000000	000000	000000	000000	000000	000000	177777
014374:	100000	000200	001340	000000	000000	000000	000000	000000	014404:	000000	000000	000000	000000	000000	000000	000000	177777

PLS  
W4  
88

014414:	000000	000000	000000	000000	000000	000000	000000	000000	014424:	000000	000000	000000	000000	000000	000000	000000	000054
014434:	000001	000000	000000	000000	001750	001750	000144	000144	014444:	000264	000454	000000	000380	000312	000230	000375	000356
014454:	000010	000000	000000	000000	000000	000000	000000	000000	014464:	000000	000000	000000	000000	000000	000000	000000	000000
014474:	000132	100074	023133	177777	000000	000074	001300	007423	014504:	012500	177844	006415	000001	070823	000000	001512	100074
014514:	000000	000000	001000	000000	013314	008084	000230	000000	014524:	013314	000000	014748	101023	000000	000000	024812	000000
014534:	000000	000001	007423	002714	000000	000001	002678	000000	014544:	000120	000000	040037	000400	000000	013314	000000	001580
014554:	042454	001580	041504	000000	000784	000238	000000	000704	014564:	000000	001300	000000	000303	001000	000000	000000	000003
014574:	005105	100423	000022	177758	002688	141074	100000	000000	014604:	001000	000003	000000	000000	000304	000354	000003	000362
014614:	014748	101023	000014	022270	100023	000017	001300	000200	014624:	000000	015203	025228	101033	000007	013314	004000	000000
014634:	001000	013314	000001	000000	013314	017277	100074	000012	014644:	000000	000000	000238	000000	177777	000235	000000	000000
014654:	000303	001000	000000	133201	000003	005105	100423	000031	014664:	000000	000000	000303	000000	013887	000002	000121	000003
014674:	000000	000007	014133	103074	000020	000000	000011	000002	014704:	000000	013704	020884	000000	000000	012023	000128	000000
014714:	000531	012784	100074	000015	001300	002714	000000	000001	014724:	012000	000054	000000	010000	000000	031171	014200	103074
014734:	000015	015914	000001	000000	001000	102074	012000	000001	014744:	013022	000000	005177	000000	17134	020884	000000	000012
014754:	018024	017178	000001	017201	000001	002418	000000	001280	014764:	000413	002770	000003	005195	103033	000031	000000	000000
014774:	000303	000000	013775	170000	012023	002360	141155	000001	015004:	020884	177777	000004	000043	000011	000000	000000	000000
015014:	020884	000000	002714	020711	177777	177775	002360	141155	015024:	000031	020884	000004	000004	021711	040000	000004	000007
015034:	024854	100032	000012	087423	000038	177777	007254	004914	015044:	144472	000068	000000	000000	007406	000000	000000	000000
015054:	002714	020711	177777	177775	002360	141155	000031	020884	015064:	000004	000004	021711	040000	000007	024854	100013	
015074:	000012	000000	000000	000000	041401	000013	000000	000002	015104:	000015	000002	000015	000000	000001	104200	000001	107023
015114:	000000	003177	011851	100474	000031	000002	000002	000001	015124:	107378	000001	107401	000001	107401	000001	107377	000001
015134:	107023	000002	000000	000013	012225	100074	000022	000000	015144:	157000	000002	006000	000003	082400	000001	000001	020884
015154:	000000	000000	002071	177777	177775	002360	141155	000031	015164:	020884	000004	000004	021711	040000	000004	000007	024854
015174:	100033	000012	000012	000000	000000	000000	000000	000000	015204:	000000	000000	000000	000000	000000	000000	000000	000000



Dump Example

MESSAGE HBR TABLE:  $(5 * 55) + 34340 = 34701$   
(ENTRY LENGTH \* PIN #) + MSG HBR TABLE BASE

WORD 1: JUNK PORT

WORD 1: USER SOFT INTERRUPTS,  
(0:1) = 1 = NEXT MESSAGE IN SECONDARY MESSAGE TABLE  
(1:15) = 12 = INDEX OF ENTRY IN TABLE

WORD 2: SYSTEM SOFT INTERRUPTS

WORD 3: RESERVED

WORD 4: BIT MAP = 040000 = 0100 000 000 000 000

BIT ONE BEING SET CORRELATES WITH AN ENTRY EXISTING FOR  
WORD 1 = USER SOFT INTERRUPTS

NOTE: ENTRY EXISTS FOR EACH PCB BUT NOT PRINTED BECAUSE ONLY  
ONE ENTRY IS NON-ZERO



HP3000 III MEMORY DUMP.00.02 OF SYS VER C UPDATE 01 FIX 09 DUMP TIME 10/16/81, 4:01PM BANK 0  
(C) HEWLETT-PACKARD CO. 1980

##### DST 48 (SPECIAL REQUEST TABLE) #####  
034174: 000023 000005 000023 000024 000017 000009 000012 001375 034204: 000000 000644 000000 000017 002375 000200 000440 000000  
034214: 000000 003035 000200 000840 000000 000031 001375 000200 034224: 000844 000000 000038 001435 000200 000844 000000 000043  
034234: 002375 000200 000440 000000 000050 002375 000000 000000 034244: 000000 000055 003028 000200 000440 000000 000042 004228  
034254: 000000 000000 000000 000007 004088 000200 000840 000000 034264: 000074 004268 000000 000200 000000 000101 004268 000000  
034274: 000000 000000 000108 001375 000200 000844 000000 000113 034304: 001435 000200 000844 000000 000120 002375 000200 000440  
034314: 000000 000125 002375 000000 000000 000000 000132 002775 034324: 000000 000000 000000 000137 001375 000200 000844 000000  
034334: 000005 001435 000200 000844

##### DST 71 (MSG HBR TABLE) #####  
034340: 000000 000000 000000 000000 000000 000000 000000 000000 034350: 000000 000000 000000 000000 000000 000000 000000 000000  
LINES 034380 - 034477 SAME AS ABOVE  
034700: 000000 000000 000012 000000 000000 000000 000000 000000 034710: 000000 000000 000000 000000 000000 000000 000000 000000

##### DST 72 (PRIMARY MSG TABLE) #####  
034720: 000000 000005 000030 000005 000170 100000 000012 013314 034730: 000000 000000 100000 000017 000000 000000 000000 100000  
034740: 000024 000012 000000 000000 100000 000031 000017 000000 034750: 000000 100000 000038 000024 000000 000000 100000 000043  
034760: 000031 000000 000000 100000 000050 000038 000000 000000 034770: 100000 000055 000843 000000 000000 100000 000082 000050  
035000: 000000 000000 100000 000007 000053 000000 000000 100000 035010: 000074 000062 000000 000000 100000 000101 000087 000000  
035020: 000000 100000 001108 000074 000038 000000 100000 000113 035030: 001101 000000 000000 100000 000120 000108 000000 000000  
035040: 000000 000125 000113 000000 000000 100000 000132 000120 035050: 000000 000000 100000 000137 000123 000000 000000 100000  
035060: 000144 000132 000000 000000 100000 000151 000137 000000 035070: 000000 000000 100000 000156 000144 000000 000000 100000 000183  
035100: 000000 000000 100000 000170 000156 000000 000000 035110: 100000 000000 000183 000000 000000 000000 000000 000000

OVERHEAD

##### DST 74 (SECONDARY MSG TABLE) #####  
035120: 000010 000005 000023 000017 000000 000000 177777 000011 035130: 107304 000001 100005 000004 000010 107304 000001 100000  
035140: 000024 000012 000000 000000 100000 000031 000017 000000 035150: 000000 100000 000038 000024 000000 000000 100000 000047  
035160: 000031 000000 000000 100000 000050 000038 000000 000000 035170: 100000 000055 000043 000000 000000 100000 000082 000050  
035200: 000000 000000 100000 000007 000053 000000 000000 100000 035210: 000074 000062 000000 000000 100000 000101 000087 000000  
035220: 000000 100000 000108 000074 000038 000000 100000 000113 035230: 000101 000000 000000 100000 000120 000108 000000 000000  
035240: 000000 000125 000113 000000 000000 100000 000132 000120 035250: 000000 000000 100000 000137 000123 000000 000000 100000  
035260: 000144 000132 000000 000000 100000 000151 000137 000000 035270: 000000 000000 100000 000156 000144 000000 000000 100000 000183  
035300: 000151 000000 000000 000000 100000 000170 000156 000000 035310: 100000 000000 000183 000000 000000 000000 000000 000000

FIRST WORDS OF FREE ENTRIES

##### DST 47 (VIRTUAL DISK SPACE TABLE) #####  
035320: 000184 000001 000020 001000 000004 000020 001402 001347 035330: 000000 000000 000000 000000 000000 000000 000000 000000  
035340: 000020 000001 000000 001000 000000 012000 002400 001402 035350: 000121 001500 001347 000000 000000 000000 000000 000000  
035360: 000000 000000 000000 000000 000000 000000 000000 000000 035370: 000000 000000 000000 000000 000000 000000 000000 000000  
035400: 000000 000000 000000 000000 000000 000000 000000 003777 177777 035410: 177770 000000 000000 000000 000000 000000 000000 000000



HP3000 III MEMORY DUMPC.00.02 OF SYS VER C UPDATE 01 FIX 09 DUMP TIME 10/18/81. 4:01PM (C) HEWLETT-PACKARD CO. 1980

BANK 1

\*\*\*\*\* PCBX AND STACK MARKERS FOR DST 132 (PCB 54) \*\*\*\*\*

SEG REL	SEG REL	JMAT	JPCNT	JOB INPUT	JOB OUTPUT	JOB DST	JIT DST	JOB TYPE	DUPLICAT	INTERACT	INIT Q	JOB INDEX
DL	DB	INDEX	INDEX	LOG DEV #	LOG DEV #	INDEX	INDEX		YES	YES		
001044	001200	1	2	20	20	126118	115	852	YES	YES	008302	0

ADDRESS	BANK	X	DELTA #	STATUS	DELTA Q	SEGMENT
077236	1	177756	020003	101074	000012	74
077224	1	000021	018636	141074	000014	74
077210	1	008337	001725	142436	006020	38
077170	1	000041	004074	082304	006007	384 USER SEGMENT
077161	1	000041	000207	082304	006004	384 USER SEGMENT
077155	1	000000	000053	082304	000024	384 USER SEGMENT
077131	1	000000	000000	140041	000004	41

\*\*\*\*\* DST 132 \*\*\*\*\*

\*\*\*\*\* PCBX: \*\*\*\*\*

\*\*\*PGLOBAL:

087423: 001044 001200 173407 000424 001024 126118 018115 000000

\*\*\*FIXED:

087433:	000120	008413	012580	008302	000134	000812	000000	000004	087443:	000000	000000	000000	000000	000004	051045	000000	000000
087453:	000000	000000	112783	840002	054000	000000	000000	010714	087480:	000000	005253	000344	002202	000000	000000	000000	000000
087473:	000000	000000	000000	000001	000000	000000	000000	000000	087503:	000000	000000	000000	004641	005255	000000	000000	000000
087513:	000000	000000	000000	000000	000000	000000	000000	000000	087523:	000000	000000	000000	000000	000000	000004	000000	000000
087533:	000000	000000	000000	000000	000000	000000	000000	000000	087543:	000000	000000	000000	000000	000000	000000	000000	000000

\*\*\*PHFILE: (ZERO TABLE ENTRIES ARE NOT PRINTED)

087553: 000116 000000 000000 000000 000000 000050 000000 000000 087583: 000000 000000 000000 000000 000000 000000 000000 000000

087573: 000170 000132 000000 000000 000000

----- FILE VECTOR TABLE -----

ENTRY	ADDRESS	LOCK	BRK	LOCK COUNT/PIN	HIPRI TAIL	HIPRI HEAD	LOPRI TAIL	LOPRI HEAD
087800:	000108	100454	000000	000000	0	108	LOCK	
087804:	000128	100454	000000	000000	1	128	LOCK	
087810:	000148	100454	000000	000000	2	148	LOCK	
087814:	000168	000000	000000	000000	3	168		
087820:	000246	000000	000000	000000	4	246		
087824:	000355	100454	000000	000000	5	355	LOCK	
087830:	000408	100454	000000	000000	6	408	LOCK	
087834:	000437	100454	000000	000000	7	437	LOCK	

----- CONTROL BLOCKS -----

087700(080105):	000001	140020	000000	022123	052104	044518	020040	002244	001700	000126	000050	000000	087700:	.....	STDIN	.....	P
087714(080121):	000430	080020	000000	000000	177773	140020	000000	022123	052104	044518	051524	002814	087714:	.....	STDLE3	.....	P
087730(080135):	001701	000121	000051	000000	000000	000000	000000	000000	177788	010041	000000	041817	087730:	.....	Q	.....	P
087744(080151):	048515	048518	047040	000255	021400	000120	000050	000000	000000	000000	000000	087744:	.....	WARD	.....	P	
087760(080165):	000000	100080	000000	041501	052101	046117	043440	000001	000428	000120	001200	000000	087760:	.....	C. CATALOG	.....	F
087774(080201):	000000	000000	000000	000000	000000	000000	000000	000000	000000	000000	000000	000000	087774:	.....		.....	
070010(080215):	000000	177777	177777	002112	000000	000000	000000	000020	000000	000000	000000	000000	070010:	.....	J	.....	
070024(080231):	000000	000000	000000	037001	000000	000000	000000	000000	000000	000000	000000	000000	070024:	.....		.....	
070040(080245):	000000	130080	000000	048517	047183	040524	020040	000005	000420	000120	001200	000000	070040:	.....	G. MONCAT	.....	F
070054(080261):	000000	080000	000000	000000	001200	000000	000040	000000	000000	000000	000270	000000	070054:	.....		.....	
070070(080275):	000000	177777	177777	008112	000000	000000	000001	000020	000000	000000	000000	000000	070070:	.....		.....	
070104(080311):	000000	000000	001000	077401	000000	000000	000000	000000	000000	000000	000000	000000	070104:	.....		.....	
070120(080325):	000000	000000	000000	045514	047516	042440	020040	002001	000787	000400	000270	000000	070120:	.....	KLONE	.....	
070134(080341):	000000	004000	000000	000000	006300	000000	000037	000000	000000	000000	000033	000000	070134:	.....		.....	

----- ADDR INTO CONTROL BLOCK -----

Dump Example

HP3000 III MEMORY DUMP. 00.02 OF SYS VER C UPDATE. 01 FIX 09 DUMP TIME 10/16/81. 4:01PM  
 (C) HEWLETT-PACKARD CO. 1980

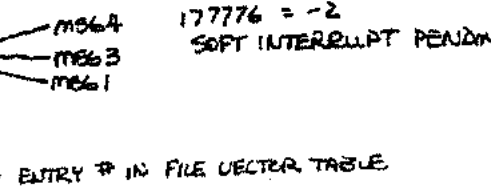
BANK 1

```

070150:000355: 1A0831 000310 048523 043481 020048 070040 032001 002103 000000 000200 000000 000000 070150: MSG1 4.C.....
070164:000371: 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 070164: MSG2 4.C.....
070200:000405: 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 070200: MSG3 4.C.....
070214:000421: 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 070214: MSG4 4.C.....
070230:000435: 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 070230: MSG4 4.C.....
070244:000451: 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 070244:
070260:000465: 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 070260:
070282:000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 070282:
070283:000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 070283:
LINES 070303 - 070402 SAME AS ABOVE
    
```

```

070403:000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 070403:
***** AVAILABLE FILE TABLE *****
070417: 100000 000126 010132 177776 11 10 0 124 0 132 177776 070417:
070417: 100000 000134 014132 177778 11 10 0 124 0 132 177778 070417:
070423: 100000 000124 012132 000000 10 10 0 124 0 132 070423:
070427: 000000 000121 000000 000000 7 FILE 0 121 0 0 0 070427:
070433: 000000 000127 000000 000000 6 FILE 0 127 0 0 0 070433:
070437: 000000 010132 000000 000000 5 FILE 0 132 0 0 0 070437:
070443: 000000 000132 000000 000000 4 FILE 0 132 0 0 0 070443:
070447: 000000 000117 004132 000000 2 FILE 0 117 0 0 0 070447:
070453: 000000 000117 002132 000000 2 FILE 0 117 0 0 0 070453:
070457: 000000 000117 000132 000000 1 FILE 0 117 0 0 0 070457:
**RPOINTER:
070463:000000 000714 001034 001044
*****DI REGISTER*****
070467:177844: 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 070467:
LINES 070503 - 070812 SAME AS ABOVE
070813:177776: 100704 000000 177777 000000 000000 177777 000000 177777 070813:
*****OB REGISTER*****
070823:000000 001750 002400 001776 000007 002008 004022 004040 004204 004300 004474 004610 004822 070823:
070837:000014 004846 004710 004714 004722 004740 005004 005038 005072 005122 005140 005152 005168 070837:
070853:000030 005204 005228 005242 005288 005314 005358 005374 005414 005448 000001 003700 000041 070853:
070867:000044 000000 177777 000000 177777 000000 002827 000000 001750 000001 000001 000003 001750 070867:
070703:000000 000004 000005 005730 004823 005722 006034 002751 003018 000003 000110 000012 000000 070703:
070717:000074 008148 000000 006186 008188 000036 007777 000007 000110 000000 000000 003140 003140 070717:
070733:000110 000212 000212 000008 000004 000005 006500 003240 000000 000000 000000 010510 001557 070733:
070747:000124 157617 000000 000000 000000 000000 000000 000000 000000 000000 000000 010520 070747:
070753:000140 010534 006325 008308 000003 001304 112703 004311 004447 011118 177777 000000 000002 070753:
070777:000154 000000 000000 014354 014358 014360 014362 000001 000001 008205 014484 014504 014522 070777:
071013:000170 014542 014562 025040 025040 025040 040517 050124 044517 047123 020145 071162 067582 071013:
071027:000204 028040 080543 072185 080554 028045 074170 074170 074170 028040 082570 070145 081584 071027:
071043:000220 082548 028045 074170 074170 074170 058000 000000 028040 025040 025040 043117 050124 071043:
071057:000234 044517 047123 020145 071162 067582 024540 060543 072185 080554 020045 074170 074170 071057:
071073:000250 074170 028040 082570 070145 081584 082544 020108 051105 048101 052105 020162 062584 072562 071073:
071107:000264 052558 082570 070145 081584 082544 020108 051105 048101 052105 020162 062584 072562 071107:
071123:000300 087040 073141 086185 082454 020145 074180 062543 072145 082040 022570 074170 074170 071123:
071137:000314 074054 020141 086184 072541 088040 022570 074170 074170 074040 058000 045514 047518 071137:
071153:000320 042456 048517 051122 044523 027110 050105 020000 024082 034054 023445 023454 047488 071153:
071167:000344 028111 034054 031130 028111 030482 028047 042047 028040 022447 028117 030481 071167:
071201:000360 024400 000000 000000 000000 000000 000000 000000 000000 000000 024082 034054 033480 071201:
071217:000374 024101 030451 024524 044517 047123 020145 071162 087582 000000 024101 034054 030530 071217:
071233:000410 028111 032454 030530 028111 032454 030530 028111 031454 031130 028062 023532 032051 071233:
071247:000424 024062 032130 026082 033450 040481 028081 054051 024524 024047 052111 048505 023454 071247:
071263:000440 043081 030058 031454 023454 020120 042322 020122 042324 023454 043081 030058 031454 071263:
071277:000454 023473 020040 041520 052440 052111 048505 023454 043081 030058 031454 023454 020120 071277:
071313:000470 042522 020122 042521 023508 030490 027093 024584 072582 067040 073141 024047 052111 071313:
    
```





DST 137

PORT: HEAD THREAD (%16 OVERHEAD) =	220	22
NEXT =	210	21
	200	20
	150	15
	140	14
	130	13
	120	12
	110	11
	100	10
	70	7
	40	4
	30	3
	20	2
	00	

PORT 40: MQE = 50 (TAIL/HEAD SAME)  
          PORT MGR PIN = 55  
          SOFT INT SUBTYPE = 1  
          SOFT INT PLABEL = 107304

MQE 50: NEXT MQE = 47 (40 LOR 7)  
          RETURN PORT = 4  
          TIME LIST ENTRY = 1, TIME OUT OCCURED

PORT 130: MQE = 160 (TAIL/HEAD SAME)  
          RETURN PORT = 15  
          TIME LIST ENTRY = 170

TLE 170: NEXT TLE 0 IF LAST

Dump Example

MP1000 IIT MEMORY DUMPC.00 02 OF SYS VER C UPDATE 01 FIX 09 DUMP TIME 10/10/81, 4:01PM BANK 1
(C) HEMLETT-PACKARD CO. 1980
124200: 100000 000104 000000 110001 020544 100000 025453 000000 124210: 000000 000104 100000 000000 040052 000000 000000 000400
124220: 010473 000000 000000

\*\*\*\*\* CST 52
\*\*\*\* [124223 TO 145173 NOT PRINTED] \*\*\*\* \*\*\*\*\*
145174: 000032 000104 100000 000104

145200: 100000 000055 000000 110001 020424 100000 025453 000000 145210: 000000 000055 100000 000000 040070 000000 000000 000400
145220: 011375 000000 000000

WAIT QUEUES ARE PORTS

\*\*\*\*\* CST 70
\*\*\*\* [145223 TO 180373 NOT PRINTED] \*\*\*\* \*\*\*\*\*
180374: 000043 000055 100000 000055

180400: 100000 000052 000000 110001 020504 100000 025453 000000 180410: 000000 000052 100000 000000 040072 000000 000000 000400
180420: 011735 000000 000000

\*\*\*\*\* CST 72
\*\*\*\* [180423 TO 172773 NOT PRINTED] \*\*\*\* \*\*\*\*\*
172774: 000000 000052 100000 000052

173000: 100000 000004 000000 110001 020444 100000 026455 000000 173010: 000000 000004 100000 000000 000137 000000 000000 000400
173020: 021784 000000 000000

HEAD OF PORT (NEEDS P) HEAD FREE LINE TAIL FREE LINE HEAD OF TRAILER

Table with columns for address, data, and pointer. Includes annotations for 'HEAD OF PORT (NEEDS P)', 'HEAD FREE LINE', 'TAIL FREE LINE', and 'HEAD OF TRAILER'. The table contains multiple rows of hexadecimal data and pointers.

MQE PORT

Dump Example

```

HP3000 III MEMORY DUMPC.00.02 OF SYS VER C UPDATE 01 FIX 06 DUMP TIME 10/18/61. 4:01PM BANK 2
(C) HEWLETT-PACKARD CO. 1960

151783(000140): 020040 020040 020040 020040 020040 020040 020040 020040 020040 030060 030061 030060 030060 151783: 00010000
151777(000154): 042517 042040 020040 020040 020040 020040 020040 020040 020040 020040 020040 020040 020040 151777: EOF
152013(000170): 020040 020040 020040 020040 020040 020040 020040 020040 020040 020040 020040 020040 020040 152013:
LINES 152027 - 152042 SAME AS ABOVE
152043(000220): 030080 030081 030080 030080 030080 030080 030080 030080 030080 030080 030080 030080 030080 152043: 00011000000000000000000000
152057(000234): 030080 030080 030080 030080 030080 030080 030080 030080 030080 030080 030080 030080 030080 152057: 000000000000000000000000
LINES 152073 - 152100 SAME AS ABOVE
152107(000284): 030080 030080 030080 030080 030080 030080 030080 030080 030080 030080 030080 030080 030080 152107: 0000000000000000
152123(000300): 177777 177777 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 152123:
152137(000314): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 152137:
LINES 152153 - 152172 SAME AS ABOVE
152373(000550): 000000
152374(000551): 000000 000003 100000 000003 152373:
152374:

```

```

152400: 100000 000038 000000 110001 013744 100000 020054 000000 152410: 000000 000038 100000 000000 040137 000000 000000 000400
152420: 031405 000000 000000

```

```

***** CST 137
*** (152423 TO 151773 NOT PRINTED) **** *****
161774: 051407 000038 100000 000038

```

```

162000: 100000 000018 000000 110001 018124 100000 020054 000000 162010: 000000 000018 100000 000000 040014 000000 000000 000400
162020: 004050 000000 000000

```

```

***** CST 14
*** (162023 TO 165373 NOT PRINTED) **** *****
165374: 031801 000018 100000 000018

```

```

165400: 100000 000051 000000 110001 017404 100000 020455 000000 165410: 000000 000051 100000 000000 000132 000000 000000 000400
165420: 021504 000000 000000

```

[\*\*\*\*\* PCBX AND STACK MARKERS FOR DST 133 (PCB 55) \*\*\*\*\*]

SEG REL DL	SEG REL DB	JMAT INDEX	JPCNT INDEX	JOB INPUT LOG DEV #	JOB OUTPUT LOG DEV #	JDT DST INDEX	JIT DST INDEX	JOB TYPE	DUPLICAT	INTERACT	INIT Q	JCUT INDEX
000644	001000	1	2	20	20	120118	115	852	YES	YES	000302	0
ADDRESS	BANK	X	DELTA P	STATUS	DELTA Q	SEGMENT						
175028	2	177758	020003	101074	000012	74						
175029	2	170520	018038	141074	000014	74						
175010	2	000117	001725	142438	000020	36						
174770	2	000041	004074	062304	000007	30#	USER SEGMENT					
174781	2	000041	000207	062304	000004	30#	USER SEGMENT					
174755	2	000000	000053	062304	000024	30#	USER SEGMENT					
174731	2	000000	000000	140041	000004	41						





```

HP3000 III MEMORY DUMPC.00 02 OF SYS VER C UPDATE 01 FIX 09 DUMP TIME 10/10/81, 4:01PM BANK 2
IC) HENLETT-PACRARD CO. 1980

050228(000127): 177777 001200 001200 023730 100074 000007 000000 000000 001000 012214 000001 000001 050228:.....
050242(000153): 001003 024124 100074 000012 000000 000000 000230 000000 177777 000000 000000 000000 050242:.....
050256(000187): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 050256:.....
LINES 050272 - 052881 SAME AS ABOVE
052882(002573): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 052882:.....

052887: 000000 000000 000000 000000 000000 000000 000000 000000 052877: 000000 000000 000000 000000 000000 000000 000000
052707: 000000 000000 000000 000000 000000 000000 000000 000000 052717: 000000 000000 000000 000000 000402 000000 000000 000461
052727: 034735 000000 000000 000000 000000 000000 000000 000000 052737: 000000 000000 000000 000000 000000 000000 000000
052747: 000000 000000 000000 000000 000000 000000 000000 000000 052757: 000000 000000 000000 000000 000000 000000 000000
052767: 000000 000000 000000 000000 000000 000000 000000 000000
052774: 000000 000018 100000 000018

053000: 100000 000032 000000 110001 017344 100000 024050 000000 053010: 000000 000032 100000 000000 040031 000000 000000 000400
053020: 007475 000000 000000

#####
** 53023 IC 81273 NOT PRINTED) **** #####
061374: 041605 000032 100000 000032

081400: 100000 000004 000000 110001 020404 100000 028455 000000 081410: 000000 000004 100000 000000 000125 000000 000000 000400
081420: 020400 000000 000000
#####

##### DST 15 #####
081423(000000): 000004 000115 100004 000000 000000 000000 000000 000000 000000 000000 000000 000000 081423:.....
081427(000014): 000000 000000 000010 000000 000000 000000 000000 000000 000000 000000 000000 000000 081427:.....
081453(000030): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 081453:.....
081467(000044): 000000 000188 000188 027112 000000 000000 000000 000000 000000 000000 000000 000000 081467:.....
081503(000060): 000010 000001 001000 037401 000000 000000 002007 000000 000182 000000 000000 000000 081503:.....
081517(000074): 000000 177777 177777 177777 177777 177777 177777 177777 177777 177777 177777 177777 081517:.....
081531(000110): 177777 177777 177777 177777 177777 177777 000000 000000 000000 000000 000000 000000 081531:.....
081547(000124): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 081547:.....
081563(000140): 000000 000000 000000 000000 000000 000000 000000 000000 000182 000177 000000 000000 081563:.....
081577(000154): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 081577:.....
LINES 081813 - 082722 SAME AS ABOVE
082223(000600): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 082223:.....
082237(000614): 044120 042440 020040 020040 000000 000000 000000 000000 000000 000000 000000 000000 082237:.....
082253(000630): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 082253:.....
082267(000644): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 082267:.....
082303(000660): 000000 000002 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 082303:.....
082317(000674): 050125 041040 020040 020040 051531 051440 020040 020040 000000 000000 000000 000000 082317:.....
082333(000710): 000000 000000 000401 031201 000401 031242 000401 031303 000000 000000 000704 177777 082333:.....
082347(000754): 000011 000401 000128 000000 000000 000000 000000 000103 000000 000000 000000 000103 082347:.....
082363(000740): 001013 000000 000404 018000 000000 000000 000000 000000 000000 000000 000000 000000 082363:.....
082374(000751): 020040 000004 100000 000004 082374:.....

082400: 100000 000001 000000 110001 018444 100000 005012 000000 082410: 000000 000001 100000 000000 000000 000000 000000 000400
082420: 018274 000000 000000
#####

##### DST 05 [SYSTEM JIT] #####

```



HP3000 III MEMORY DUMPC.06.02 OF SYS VER C UPDATE 01 FIX 08 DUMP TIME 10/18/81. 4:01PM  
(C) HEWLETT-PACKARD CO. 1980

BANK 3

##### DST 117 #####  
165223(000000): 000074 000117 140004 000000 000000 000014 000000 000000 000000 000003 037417 004300 165223: .X.O  
165237(000014): 100060 000003 041517 048515 040518 042040 000255 001400 000170 000050 000000 000000 165237: .O.COMMAND  
165253(000030): 000000 000000 000000 177788 000000 004520 000000 000000 000000 004520 000000 004520 165253: .P  
165267(000044): 177777 177777 000000 000000 002410 010001 030001 000000 000000 000000 000000 000004 165267: .P  
165303(000060): 000016 001400 037024 000000 000000 000000 000000 000000 000000 000000 000000 000000 165303: .O  
165317(000074): 047408 037777 081074 141503 000004 004543 102033 000011 000020 000001 177134 000000 165317: .O.T.B.C.E  
165333(000110): 002250 000000 000001 022057 100433 000010 000000 020220 000000 002404 000203 001480 165333: .P  
165347(000124): 000000 177134 020220 000000 000013 017700 000014 000000 000203 000000 002404 000303 165347: .P  
165363(000140): 001600 080413 003645 000003 000578 102033 000031 000014 000000  
165374(000151): 000303 000001 100000 000001 165374: .P

165400: 100000 000041 000000 110001 020244 100000 021443 000000 165410: 000000 000041 100000 000000 040040 000000 000000 000400  
165420: 010076 000000 000000

##### CST 40  
\*\*\*\* (165223 TO 175573 NOT PRINTED) \*\*\*\* #####  
175574: 141202 000041 100000 000041

175800: 100000 000004 000000 110001 015704 100000 020054 000000 175810: 000000 000004 100000 000000 000134 000000 000000 000400  
175820: 020714 000000 000000

##### DST 134 #####  
175823(000000): 000004 000134 140004 000000 000000 000015 000000 000000 000000 000000 000000 000000 175823: .w.MSG2 4.0  
175837(000014): 000000 100587 000011 048523 043453 020040 020040 022001 002100 000400 000205 000000 175837: .P  
175853(000030): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 175853: .P  
175867(000044): 000000 000388 000388 020112 000000 000401 000000 000001 000001 000182 000215 000000 175867: .P  
175703(000060): 000020 000001 001000 037401 004703 000000 002007 000000 000192 000000 000000 000001 175703: .P  
175717(000074): 000000 177777 177777 177777 177777 177777 177777 177777 177777 177777 177777 177777 175717: .P  
175733(000110): 177777 177777 177777 177777 177777 000000 000000 000000 000000 000000 000000 000000 175733: .P  
175747(000124): 000000 000000 000014 000014 000000 000000 000000 000000 000203 000000 000000 000000 175747: .P  
175763(000140): 000000 000000 000000 000000 000000 000000 000000 000000 000162 000377 000000 000000 175763: .P  
175777(000154): 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 175777: .P  
LINES 175013 - 175422 SAME AS ABOVE  
175423(000000): 000000 000000 000000 000000 008415 008415 005415 008415 000820 000000 000000 000000 175423: .P  
175437(000014): 000000 000000 000000 000000 000030 000000 000000 000000 000000 000000 000000 000000 175437: .P  
175453(000030): 000000 000000 000000 000000 000000 000000 000000 000000 177777 177777 177777 177777 175453: .P  
175467(000044): 000000 000000 000000 000000 000280 000000 000000 000000 000000 000000 000000 000000 175467: .P  
175503(000060): 006270 000000 000000 000000 000000 000000 000000 000000 000300 000000 000000 000000 175503: .P  
175517(000074): 000000 000000 000000 000000 000310 000000 000000 000000 000000 000000 000000 000000 175517: .P  
175533(000110): 000320 000000 000000 000000 000000 000000 000000 000000 000330 000000 000000 000000 175533: .P  
175547(000124): 000000 000000 000000 000000 000340 000000 000000 000000 000000 000000 000000 000000 175547: .P  
175563(000140): 000350 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 175563: .P  
175574(000151): 000000 000004 100000 000004 175574: .P

175600: 020000 000003 000000 000000 000000 000000 000000 000000 174623 175610: 000000 000003 000000 000000 000000 000000 000000 000400  
175620: 021080 000000 000000

##### AVAILABLE AREA #####  
\*\*\*\* (175823 TO 177373 NOT PRINTED) \*\*\*\*

# MMSTAT DEFINITIONS

APPENDIX

C

If problems arise using soft interrupts, the monitor facility can be used to analyze the specific steps that have occurred when using soft interrupts. By enabling MMSTAT (soft interrupts) before the execution of the application, specific information can be logged to tape. Once the application has executed the logging tape can be analyzed to assist in determining the existing problem. The MPE Tables Manual contains the definitions of the entries in the printed log file.

Octal Value	Event Type	Parameter 0	Parameter 1	Parameter 2
62	Open	Port Number	Port DST No.	Flags Parameter
63	Receive Completion	Port Number	MQE Address 15:1 Waitspc	Return Port
64	Send	Port Number	MQE Address 15:1 Q Type	Return Port
65	Change Status	Port Number	0 = Enable 1 = Disable	Head MQE Address
66	Abort	Port Number	Parameter Zero	Return Port
67	Close	Port Number	Port DST	# Open Ports Left
70	Expand	Port DST No.	# Expanded Blocks	Total # of Blocks
71	Timeout Expired	Port Number	MQE Address	Return Port

## MMSTAT Definitions

Octal Value	Event Type	Parameter 1	Parameter 2
72/0	Read Init	# free rec	
72/1	Read Compl	(0:8) error, (8:8) ID	No of Records
72/2	Write Init	(0:8) # rec, (8:8) ID	No of Free Records
72/3	Write Compl	(0:8) error, (8:8) ID	No of Free Records
72/4	Control	(0:8) error, (8:8) ID	(0:4) func, (4:12) parameters
72/5	EOF	(0:8) error, (8:8) ID	No of Records
72/6	Open	(0:8) error, (8:8) ID	No of Records
72/7	Close	(8:8) #Free, (8:8) ID	No of Records
72/10	Initiation	0	(0:8) fix, (8:8) update
73/0	Put Record	(0:8) error, (8:8) ID	(0:3) record type, (3:13) no of records
73/1	Delete Rec	(0:8) error, (8:8) ID	(0:3) record type, (3:13) no of records
73/2	Delete Block	Start of file block #	End of file block #

### Notes:

1. The aa/bb notation in the "octal value" column denotes type/subtype. Type is the actual MMSTAT event number. Subtype is (0:4) of parameter 0.
2. Several items can possibly exceed their fields, in that case the bits beyond the field are lost. These items are number of records, number of free records, start of file, and end of file.
3. Parameter word zero has a common format for all the MMSTAT events.

Field	Description
(0:4)	Event's subtype
(4:2)	File's state 0 = empty 1 = partially full

2 = only a fraction of a free record is left  
3 = completely full

- (6:1) Nonzero indicates that there is one or more waiting readers
- (7:1) Nonzero indicates that there is one or more waiting writers
- (11:1) Nonzero indicates that the write has a carriage control character
- (12:4) Flags local to the accessor
  - (12:1) - the accessor has done no FREADs/FWRITES
  - (13:1) - extended wait
  - (14:1) - nondestructive read
  - (15:1) - writer has not written any records

## Soft Interrupt MMSTAT Entry Forms

MMSTAT tracing for the following types is enabled if the "monitoring" has been turned on by the MON command.

Type	Number	Word 0	Word 1	Word 2
Cause Soft Interrupt	240	(0:4) level (4:2) type (6:2) stype (8:8) pin	Msg Word One	Msg Word two
Pseudoint	241	(0:8) type (8:8) stype	Msg Word One	Msg Word Two
Build Stack	242	Plabel	P-Reg Word of Prev Marker	Status Word of Prev Marker
Change State	243	(14:1) old state (15:1) new state	User's P Register	User's Status Register
Timeout	244	completion type 0 timeout 1 no trlx 2 control y soft int	16 msb of timeout (ms)	16 lsb of timeout (ms)

## Where:

Level	The current status of the interrupt/process
0	process is dying (hard kill or soft kill)
1	other soft interrupts are pending
2	user interrupts are disabled
3	process is impeded, critical, and/or with SIR
4	interrupt was against own process
5	process was either waiting or ready to execute



# MPE MEMORY RESIDENT MESSAGE FACILITY

APPENDIX

D

The memory resident message facility of MPE IV addresses the need for an efficient, simple, and uniform message method for system code to send short status-type messages to processes.

Each process is created with a message harbor which supports a set of message ports which are private to that process. There is a maximum of four ports per harbor in the initial implementation. This limit can be easily extended when new ports are required.

Any system code, even code running on the ICS, can send a message to any port of any process. The destination process's PIN must be known, and a priority conventions on port number and message formats must be established. The caller of SENDMSG may optionally specify that the destination process be awakened from a message wait.

The caller of SENDMSG specifies whether the message is to be buffered in the primary message table or the secondary message table. When the secondary table is specified, if the pool of secondary entries is exhausted, the calling process is queued for a message table entry and blocked until one becomes available. Use of the primary message table is reserved for code running on the ICS or during critical sections (Pdisabled or Disabled intervals) in which it is not possible to release control of the processor to queue for a free message table entry. If the primary table is specified and no free entries are available, the SENDMSG crashes the system.

Messages can be of any length up to the configured maximum. Message length is specified in the call to SENDMSG and RECEIVMSG. In the initial implementation, messages are limited to four words in length. This maximum can be easily increased if the need arises.

By calling PORTSTATUS, a process may at any time determine whether a specified port is non-empty or obtain the port number of the most urgent non-empty port (lowest numerical port number = most urgent port).

By calling RECEIVMSG, a process may received the message at the head of the specified message port. This receive is optionally nondestructive.

A process can wait on a message wait, or on a combination of a message wait and other wait types.

### Message Intrinsic

- A. Procedure SENDMSG (Destpin, Destport, Msglength, Flags);  
Value Destpin, Destport, Msglength,Flags;  
Integer Destpin, Destport, Msglength;  
Option Priviledged, Uncallable;  
Logical Flags;

Destpin, Destport and Msglength has better be within range and reasonable (process and port exist), since SENDMSG checks and will crash if the parameters are bad.

The caller of SENDMSG stacks the message contents before calling the procedure. SENDMSG expects the first message word to be at Q-7-Msglength, and the last message word at Q-8. The message contents at Q-8 to Q-7-Msglength are deleted from top of stack by the exit from SENDMSG to the caller.

Flags.(1:1) = 1 ==> Wake-up destintation process from a message wait  
(0:1) = 1 ==> Place message in secondary message table

Return CC=CCG if process was already awake else CC=CCE.

- B. Logical Procedure PORTSTATUS (Portnumber);  
Value port number;  
Integer port number;  
Option Priviledged, Uncallable;

When supplied a valid port number, PORTSTATUS returns a true value if the port is non-empty and a false value if the port is empty.

When passed a -1 as port number parameter, PORTSTATUS returns the portnumber of the process's most urgent non-empty port (the smaller the number, the more urgent the port).

If all ports are empty, PORTSTATUS returns CC=CCE. If at least one port is non-empty, PORTSTATUS returns CC=CCG.

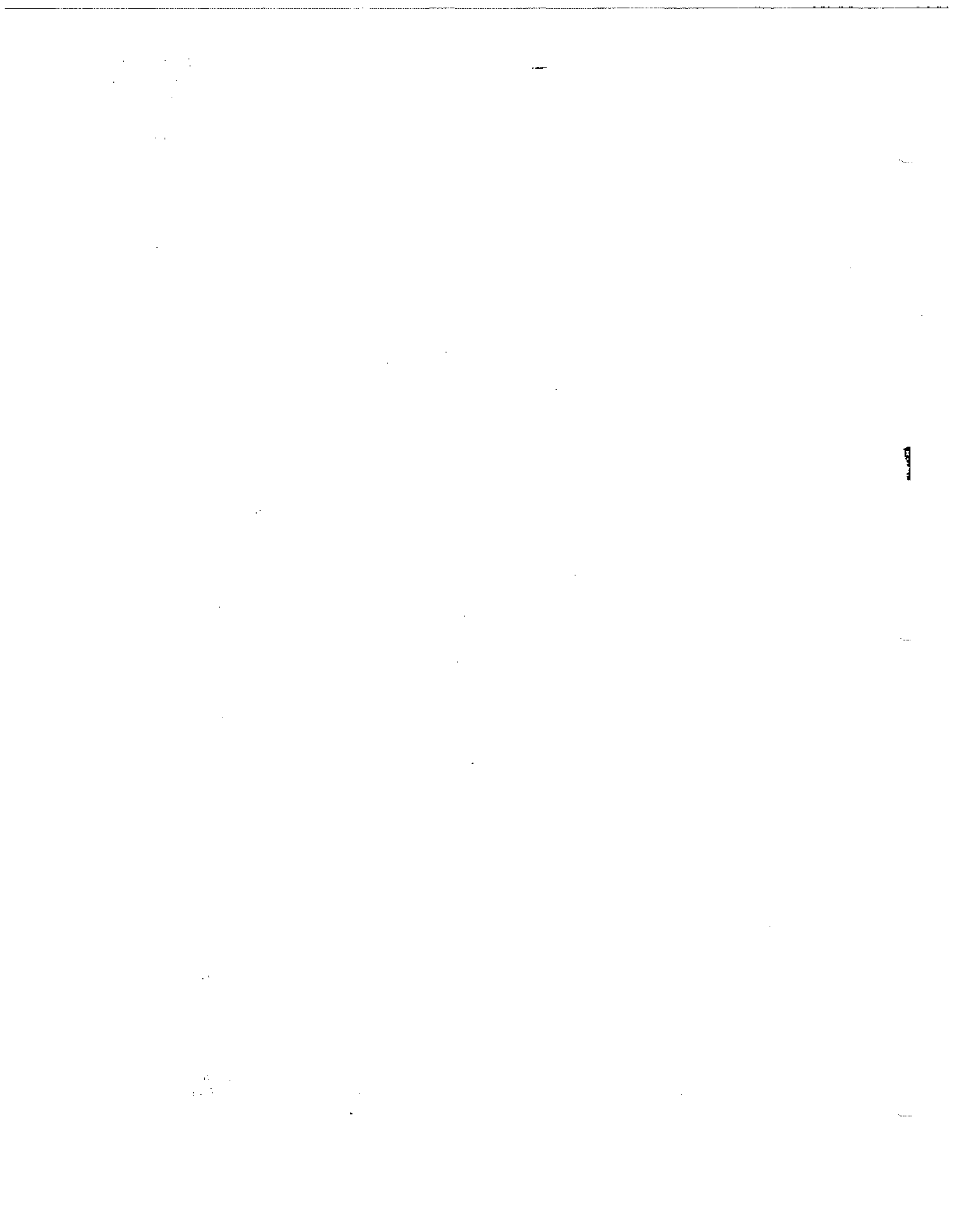
- C. Procedure RECEIVMSG (Portnum, Msglength,Flags);  
Value Portnum,Msglength, Flags;  
Integer Portnum, Msglength;  
Option Priviledged, Uncallable;

Portnum and Msglength had better be within range or else its suddendeath time.

The caller of RECEIVMSG does an ADD S Msglength to make space for the message contents. RECEIVMSG stores the message contents into Q-8, Q-9,.....Q-7-Msglength.

Q-7-Msglength contains the first word of the message.  
Flags.(0:1) = 1 ==> do not release message from head of  
port's message queue (non-destructive  
read)

Return CC=CCG if port was empty, else CC:=CCE.



# FILE SYSTEM BASIC IPC DEFINITIONS

APPENDIX

E

The objective of this set of uncallable procedures is to provide a simple IPC mechanism to support the IPC file access procedures. It enables one process to send short, control messages to another processes.

## General Behavior

### FCPORTOPEN Procedure

The heart of this mechanism is the port. A process desiring to receive messages would first open (create) a port. This process is termed the "port manager". When the port is created, a port number is returned to the opener. Since the port number value cannot be known in advance, potential senders need some method of obtaining the port number from the port manager.

Both the ports and the messages are contained in a single disc resident data segment. There can be a total of over 3,500 open ports and outstanding messages. Thus neither ports nor message blocks are scarce resources.

### FCPORTSEND Procedure

This procedure sends a 0-to 5-word message to a port. Optionally a timeout value may be specified which will limit the duration the message will remain attached to the port. Expiration of the timeout causes the message to be deleted from the target port's queue and placed on the sender's reply port (specified by the sender in the FCPORTSEND procedure call).

### FCPORTRECEIVE

Reads and deletes the head message from a port. The sender's return port number is also given to the receiver, enabling him to send a reply message.

### FCPORTCLOSE

Demolishes the port.

## IPC File's Use of This Mechanism

All open message files have two ports open for the file (read wait queue and write wait queue), plus one port per accessor (reply port). Their use is described in the following.

### Reader and Writer Wait Queues

When an empty message file is accessed by more than one reader (share), then there must be a way of having the readers' FREADs satisfied in the same order that they were issued. That is, there must be a queue of waiting readers. The IPC procedures accomplish this by dedicating a basic IPC port as a "read wait queue". Whenever a reader's request is stalled because the file is

## Basic IPC Procedures

empty, a message is sent to the read wait queue. Subsequent FREADs by other processes will queue up behind the first reader in an FIFO manner. An FWRITE will take the first entry from the wait queue and send a "read may be done" message to the reader's reply port. In a like manner multiple writers will queue on the write wait queue when the file is full.

## Completion Notification for NoWait I/O

The IOWAIT intrinsic waits for a message to be sent to the reply port(s) of the specified user files.

## Timeouts

When an accessor encounters a boundary condition (e.g. a reader accesses an empty file), it may specify that the condition must be satisfied in x seconds (FCONTROL 4). To this end the IPC access procedures merely issue the FCPORSEND to the wait queue with the user's timeout value specified. The timeout will tear the message from the wait queue and place it on the accessor's reply port.

# SYSTEM FAILURES

APPENDIX

F

## BIPC - MODULE 65

### System Failure 690

UGLYPORTPROBLEM is the equate used to describe system failure 690. System failure 690 can occur from 21 different code statements. The main description of this system failure is a problem which occurred during accessing a port. This could be an incorrect Port DST number, incorrect head entry, or an inconsistent number of MQE's between procedures.

## IPC - MODULE 66

### System Failure 495

UGLYMSGACCESS is the equate used to describe system failure 495. Calls to suddendeath using UGLYMSGACCESS can occur from 13 different locations. The causes of this system failure, buffer addresses incorrect, or a wait queue address is incorrect.

## KERNALC - MODULE 92

### System Failure 4

This system failure is common throughout KERNALC. Because this system failure is common in KERNELC it is redefined within the soft interrupt procedures within KERNELC (which means the crossreference separates calls to suddendeath specific to soft interrupts). Causes of System Failure 4 are calling parameters incorrect or problems with messages (don't exist, etc.).

